

| DIVINE TECHNOLOGY VENTURES, an Illinois general partnership, | 02C 7700 |
|--|---|
| Plaintiff, |)) Civil Action No. |
| V. |) ——— « « « « « « « « « « « « « « « « « |
| REDENVELOPE, INC., a Delaware corporation, |) Jury Trial Demanded JOAN H. LEFKOW |
| Defendant. | MAGISTRATE JUDGE KEYS |

COMPLAINT

Plaintiff divine technology ventures ("divine") brings this complaint against defendant RedEnvelope, Inc. ("RedEnvelope") seeking damages, injunctive relief and other relief for patent infringement.

THE PARTIES

- 1. divine is an Illinois general partnership with its principal place of business at 1301 North Elston Avenue, Chicago, Illinois.
- 2. Upon information and belief, RedEnvelope is a Delaware corporation with its principal place of business at 201 Spear Street, San Francisco, California.
- 3. Upon information and belief, RedEnvelope sells gifts for various occasions through catalogs and the Internet. It operates an interactive site on the World Wide Web located at www.redenvelope.com ("RedEnvelope Site") through which customers in this District and elsewhere may order gifts and other products.
- 4. RedEnvelope conducts continuous and systematic business in Illinois and this District. For example, it: (a) routinely offers to contract, via the RedEnvelope Site, with persons

in Illinois; and (b) promotes its goods and services to, and solicits contact from, persons in Illinois via the RedEnvelope Site.

5. Upon information and belief, RedEnvelope makes, uses, sells and/or offers for sale systems and processes to operate its business at the RedEnvelope Site. In general, the RedEnvelope Site includes a network-based sales system that allows customers to purchase products over the Internet and a system to control and monitor customer sessions.

JURISDICTION AND VENUE

- 6. This Court has jurisdiction over divine's patent infringement claims under 28 U.S.C. § 1331 and 1338(a).
- 7. This Court has personal jurisdiction over RedEnvelope because it has transacted and is transacting business in this District, both generally and regarding the allegations in this Complaint, and because RedEnvelope has committed tortious acts within this District.
- 8. Venue properly lies in this Court under 28 U.S.C. §§ 1391(b), (c) and 1400(b) because RedEnvelope is subject to personal jurisdiction and has committed acts of patent infringement in this District.

COUNT I

Patent Infringement of U.S. Patent No. 5,715,314 – 35 U.S.C. § 271

- 9. divine realleges and incorporates by reference the allegations in Paragraphs 1 through 8.
- 10. On February 3, 1998, United States Letters Patent No. 5,715,314 ("314 Patent") was duly and legally issued for an invention entitled "Network Sales System." A true and correct copy of the '314 Patent is attached as PX1.
 - 11. divine owns all rights, title and interest in and to the '314 Patent.

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12. RedEnvelope has directly infringed the claims of the '314 Patent by making, using, selling and/or offering to sell its systems and processes to operate the RedEnvelope Site.

- 13. divine gave written notice by letter dated July 25, 2002 to RedEnvelope of its infringement of the '314 Patent.
- 14. RedÉnvelope will not stop making, using, selling and/or offering for sale its systems and processes to avoid infringing the '314 Patent.
- 15. RedEnvelope's infringement has been deliberate, willful and wanton and with full knowledge of the '314 Patent.
- 16. RedEnvelope's conduct has caused divine to suffer and, unless enjoined by the Court, will cause divine to continue to suffer damage to its operation, reputation and goodwill.
- 17. divine has no adequate remedy at law. RedEnvelope's conduct has caused and, if not enjoined, will continue to cause irreparable damage to divine. As a result of RedEnvelope's wrongful conduct, divine is entitled to injunctive relief.

COUNT II

Patent Infringement of U.S. Patent No. 5,909,492 – 35 U.S.C. § 271

- 18. divine realleges and incorporates by reference the allegations in Paragraphs 1 through 17.
- 19. On June 1, 1999, United States Letters Patent No. 5,909,492 ("'492 Patent") was duly and legally issued for an invention entitled "Network Sales System." A true and correct copy of the '492 Patent is attached as PX2.
 - 20. divine owns all rights, title and interest in and to the '492 Patent.
- 21. RedEnvelope has directly infringed the claims of the '492 Patent by making, using, selling and/or offering to sell its systems and processes to operate the RedEnvelope Site.

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22. divine gave written notice by letter dated August 29, 2002 to RedEnvelope of its infringement of the '492 Patent.

- 23. RedEnvelope will not stop making, using, selling and/or offering for sale its systems and processes to avoid infringing the '492 Patent.
- 24. RedEnvelope's infringement has been deliberate, willful and wanton and with full knowledge of the '492 Patent.
- 25. RedEnvelope's conduct has caused divine to suffer and, unless enjoined by the Court, will cause divine to continue to suffer damage to its operation, reputation and goodwill.
- 26. divine has no adequate remedy at law. RedEnvelope's conduct has caused and, if not enjoined, will continue to cause irreparable damage to divine. As a result of RedEnvelope's wrongful conduct, divine is entitled to injunctive relief.

COUNT III

Patent Infringement of U.S. Patent No. 5,708,780 – 35 U.S.C. § 271

- 27. divine realleges and incorporates by reference the allegations in Paragraphs 1 through 26.
- 28. On January 13, 1998, United States Letters Patent No. 5,708,780 ("780 Patent") was duly and legally issued for an invention entitled "Internet Server Access Control and Monitoring Systems." A true and correct copy of the '780 Patent is attached as PX3.
 - 29. divine owns all rights, title and interest in and to the '780 Patent.
- 30. RedEnvelope has directly infringed the claims of the '780 Patent by making, using, selling and/or offering to sell its systems and processes to operate the RedEnvelope Site.
- 31. divine gave written notice by letter dated July 25, 2002 to RedEnvelope of its infringement of the '780 Patent.

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32. RedEnvelope will not stop making, using, selling and/or offering for sale its systems and processes to avoid infringing the '780 Patent.

- 33. 24. RedEnvelope's infringement has been deliberate, willful and wanton and with full knowledge of the '780 Patent.
- 34. RedEnvelope's conduct has caused divine to suffer and, unless enjoined by the Court, will cause divine to continue to suffer damage to its operation, reputation and goodwill.
- 35. divine has no adequate remedy at law. RedEnvelope's conduct has caused and, if not enjoined, will continue to cause irreparable damage to divine. As a result of RedEnvelope's wrongful conduct, divine is entitled to injunctive relief.

RELIEF REQUESTED

WHEREFORE, divine requests that the Court enter a judgment in divine's favor and against RedEnvelope and provide divine the following relief:

- A. Order, adjudge and decree that RedEnvelope has infringed the '314 Patent in violation of 35 U.S.C. § 271;
- B. Order, adjudge and decree that RedEnvelope willfully and knowingly infringed the '314 Patent;
- C. Order, adjudge and decree that RedEnvelope's infringement of the '314 Patent is exceptional under 35 U.S.C. § 285;
- D. Order, adjudge and decree that RedEnvelope has infringed the '492 Patent in violation of 35 U.S.C. § 271;
- E. Order, adjudge and decree that RedEnvelope willfully and knowingly infringed the '492 Patent;

- F. Order, adjudge and decree that RedEnvelope's infringement of the '492 Patent is exceptional under 35 U.S.C. § 285;
- G. Order, adjudge and decree that RedEnvelope has infringed the '780 Patent in violation of 35 U.S.C. § 271;
- H. Order, adjudge and decree that RedEnvelope willfully and knowingly infringed the '780 Patent;
- I. Order, adjudge and decree that RedEnvelope's infringement of the '780 Patent is exceptional under 35 U.S.C. § 285;
- J. Issue preliminary and permanent injunctive relief prohibiting RedEnvelope and its respective parents, principals, officers, agents, affiliates, servants, attorneys, employees and all others in privity with it from:
 - (i) infringing the '314 Patent;
 - (ii) infringing the '492 Patent; and
 - (iii) infringing the '780 Patent;
- K. Award divine damages for patent infringement including prejudgment interest and costs against RedEnvelope under 35 U.S.C. § 284;
- L. Award divine three times its damages to compensate divine under 35 U.S.C. § 284;
- M. Award divine its reasonable attorneys' fees under 35 U.S.C. § 285; and
- N. Award such other and further relief as the Court may deem just.

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JURY DEMAND

divine demands trial by jury.

Dated this 25th day of October 2002.

Respectfully submitted,

One of the Attorneys for Plaintiff divine technology ventures

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EXHIBIT 1



United States Patent [19]

Payne et al.

[11] Patent Number:

5,715,314

[45] Date of Patent:

Feb. 3, 1998

[54] NETWORK SALES SYSTEM

[75] Inventors: Andrew C. Payne, Lincoln; Lawrence C. Stewart, Burlington; David J.

Mackie. Cambridge, all of Mass.

[73] Assignee: Open Market, Inc., Cambridge, Mass.

[21] Appl. No.: 328,133

[56]

[22] Filed: Oct. 24, 1994

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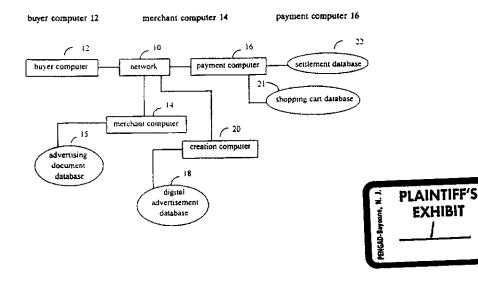
Primary Examiner—Bernarr E. Gregory Attorney, Agent, or Firm—Fish & Richardson P.C.

[57] ABSTRACT

A network-based sales system includes at least one buyer computer for operation by a user desiring to buy a product. at least one merchant computer, and at least one payment computer. The buyer computer, the merchant computer, and the payment computer are interconnected by a computer network. The buyer computer is programmed to receive a user request for purchasing a product, and to cause a payment message to be sent to the payment computer that comprises a product identifier identifying the product. The payment computer is programmed to receive the payment message, to cause an access message to be created that comprises the product identifier and an access message authenticator based on a cryptographic key, and to cause the access message to be sent to the merchant computer. The merchant computer is programmed to receive the access message, to verify the access message authenticator to ensure that the access message authenticator was created using the cryptographic key, and to cause the product to be sent to the user desiring to buy the product.

48 Claims, 25 Drawing Sheets

Microfiche Appendix Included (1 Microfiche, 34 Pages)



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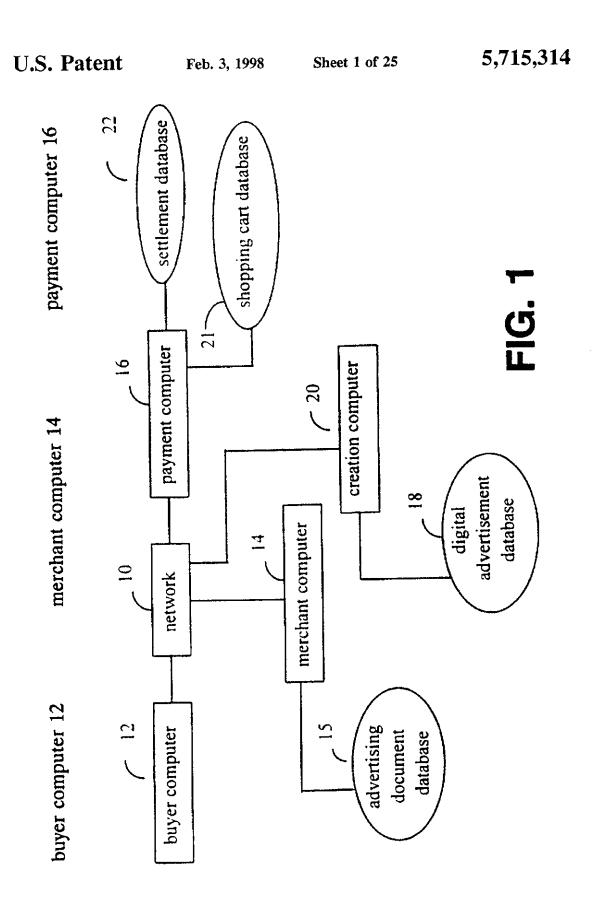
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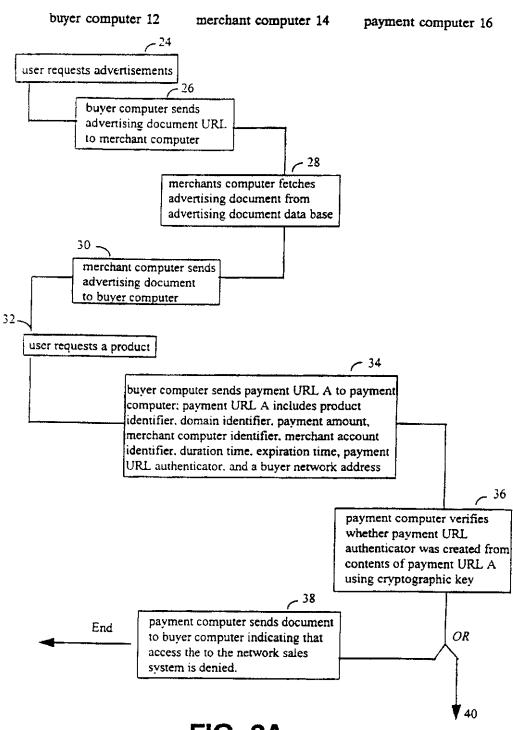


FIG. 2A

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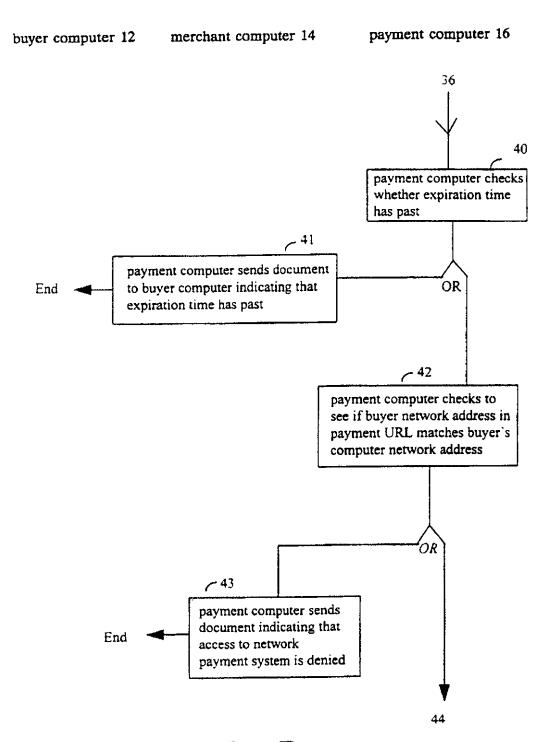
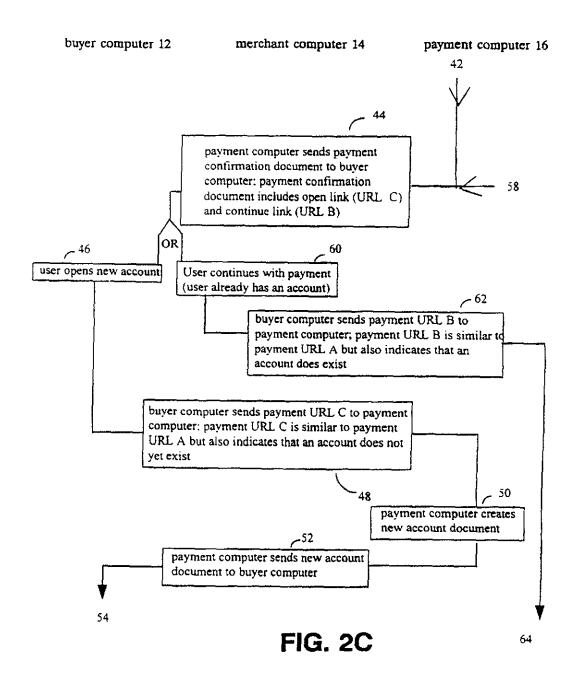


FIG. 2B

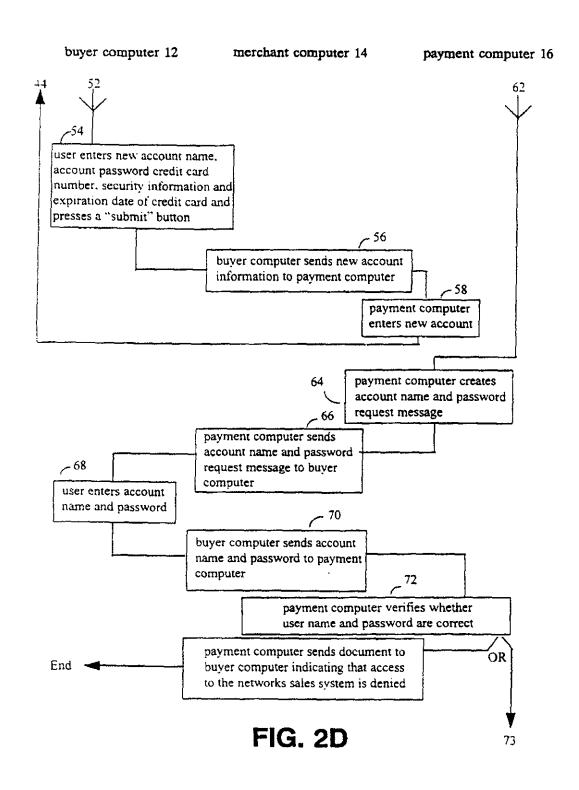
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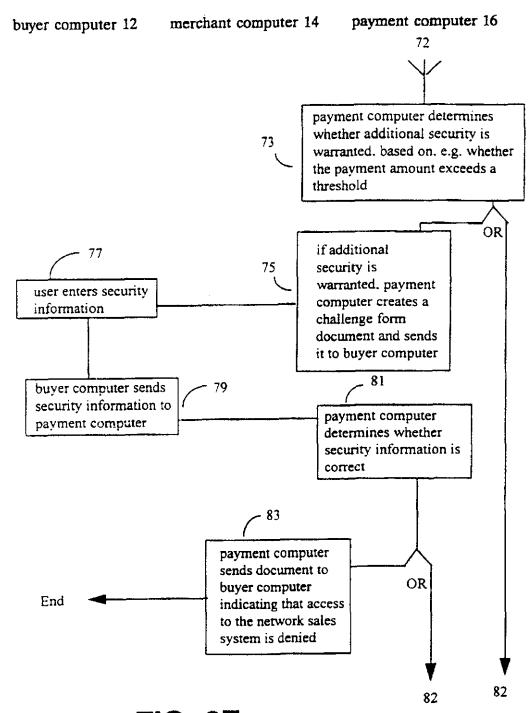


FIG. 2E

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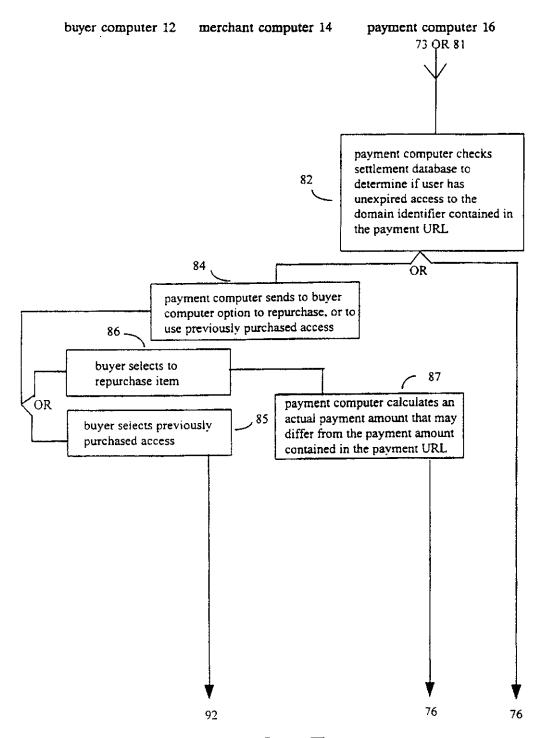


FIG. 2F

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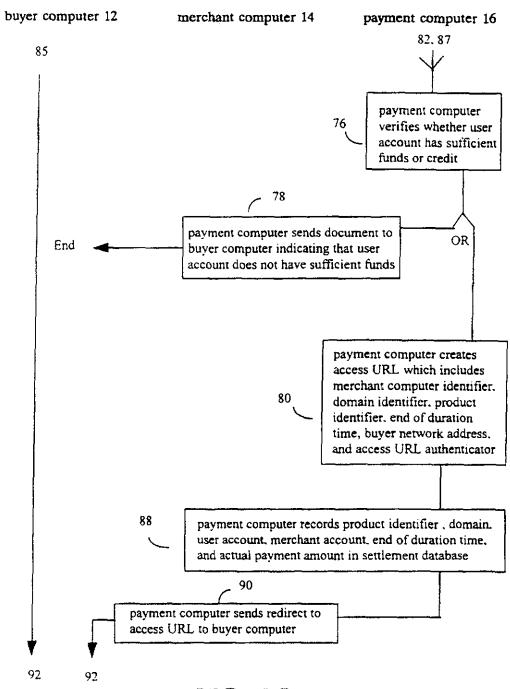
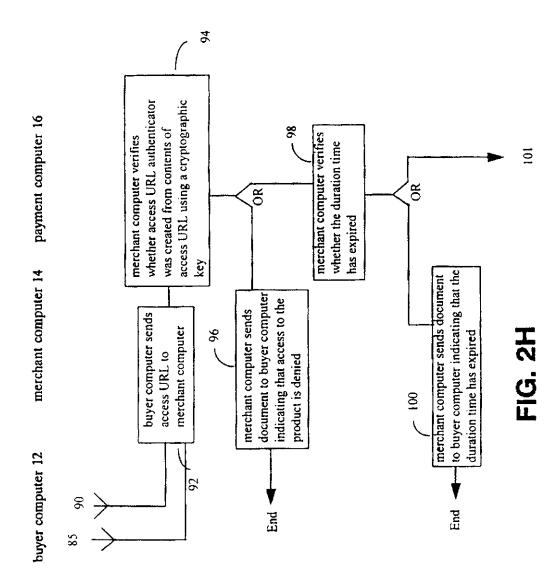


FIG. 2G

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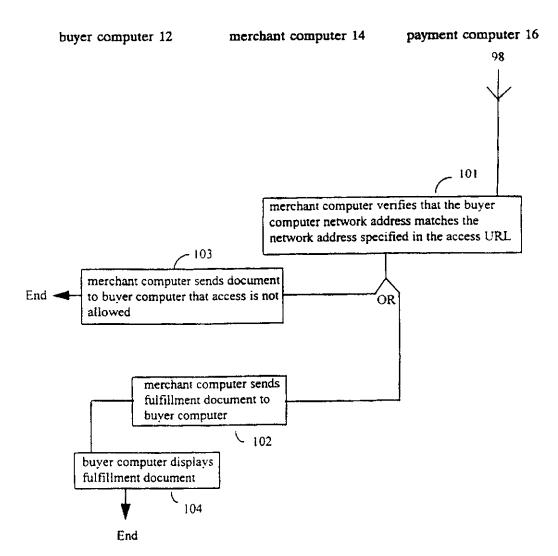
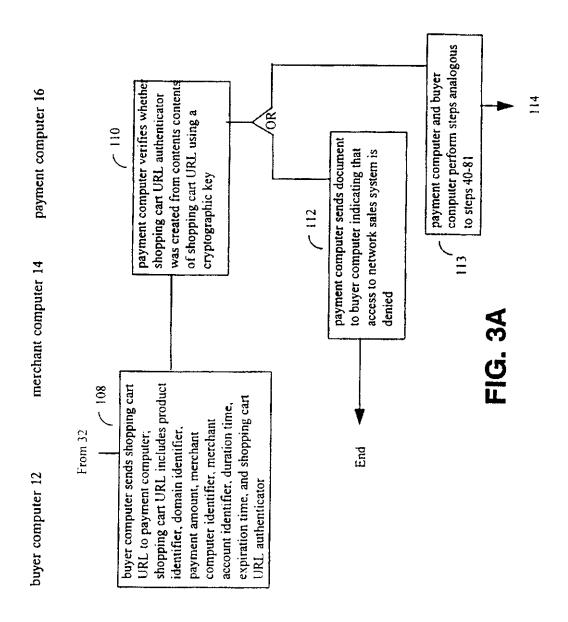


FIG. 21

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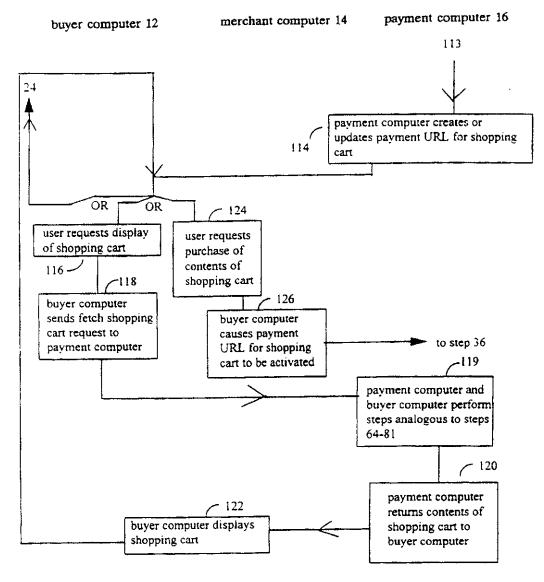
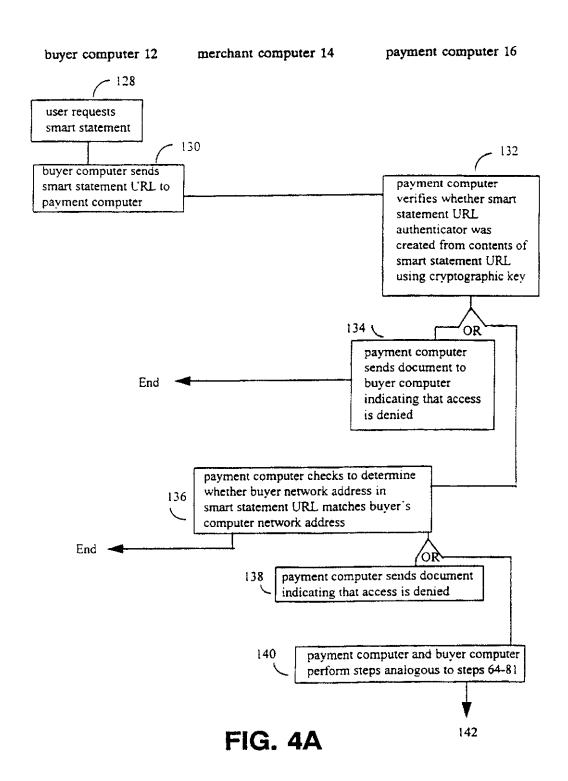


FIG. 3B

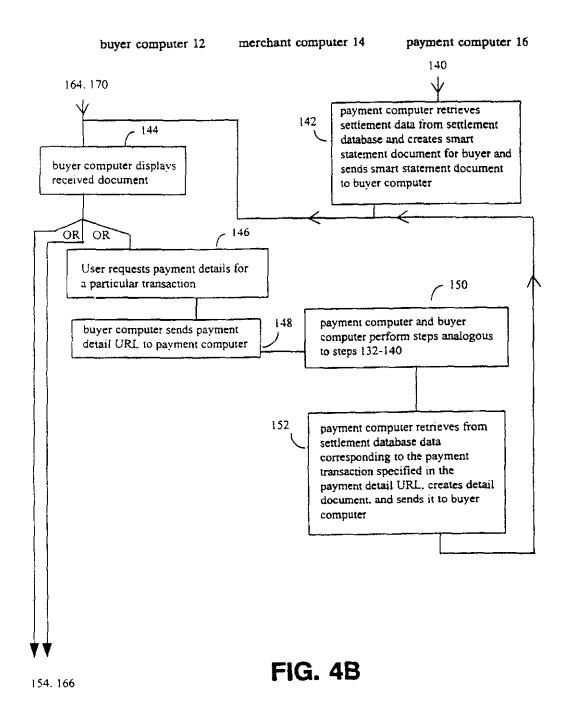
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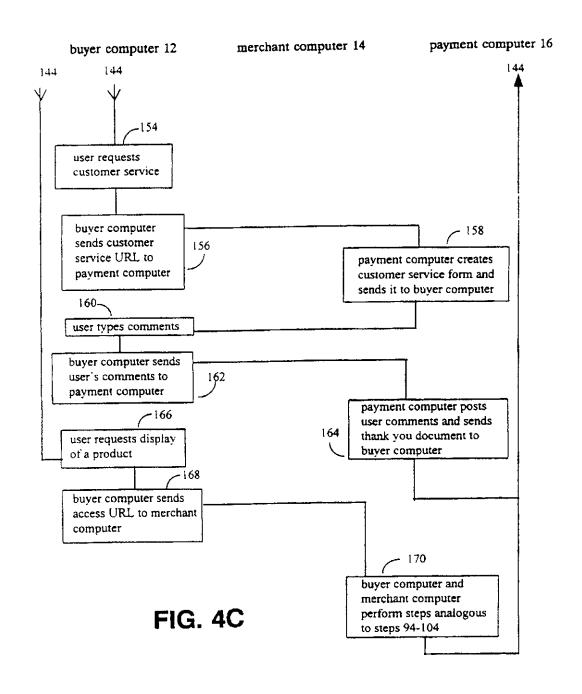
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| Mead Data Central: Internet Information | Δ |
| November 28, 1993 LC's debut on the Internet: Library of Congress catalog On the | |
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FIG. 5

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| Description | Test Merchant ps:Mead Data Central Article 55(US currency) | |
| your account name and | Market account click on "continue" below and you will be promp I password. If you do not have an account, you can establish o o this page to continue your purchase. | |
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FIG. 6

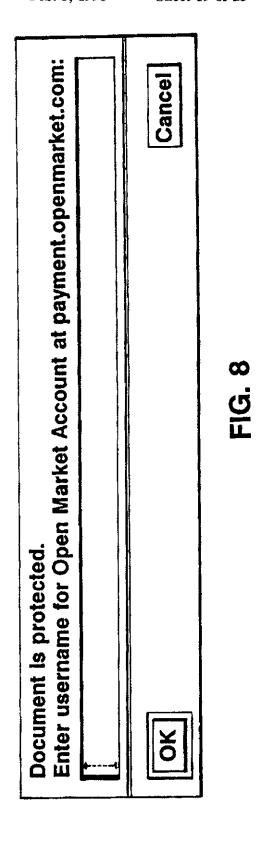
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| Card Number: Expiration Date: (format MM/YY) Check the appropriate boxes: I am the owner of the above credit card. The above address is also the billing address for this credit your OpenMarket account statement is available on-line. At your option you copy of your statement automatically sent to your e-mail address at weekly or intervals. Please choose a statement option. Weekly statements Monthly statements No e-mail statement account name and password for your OpenMarket account. suggest using an account name that is unique and easy to remember such a e-mail address. Your password should be 8 characters or longer. | may a monthly nts |
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FIG. 7

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| Document URL: http://payment.openmarket.com/ben/nph-payment | ے |
| Open Market Payment | Δ |
| You have selected an item that you have purchased recently. Merchant: Test Merchant Description: Mead Data Central Article Amount: 2.85(US currency) | |
| This could happen because you would like to buy the item again or it may have ha by accident. | ppened |
| You can: • Go directly to the previous item • Go ahead and buy the item again | |
| Open Harket, Inc. | |
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FIG. 9

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FIG. 10

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| Document URL: http://payment.openmarket.com/in/nph-stateme | |
| Information about the item. | |
| Transactions in October 1994 | |
| Mon Oct 3 Test Herchant Dilbert subscription 20 seconds amount Tue Oct 4 Test Herchant Mead Data Central Article amount \$2.95 Tue Oct 4 Test Herchant Mead Data Central Article amount \$2.95 Tue Oct 4 Test Merchant Mead Data Central Article amount \$2.95 Tue Oct 4 Test Merchant Mead Data Central Article amount \$0.50 Tue Oct 4 Test Herchant Mead Data Central Article amount \$2.95 Wed Oct 5 Test Merchant Mead Data Central Article A | \$0.10 |
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| August 1994 Return to your Newest Statement | |
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FIG. 11

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| <u>File Options Navigate Annotate</u> | <u>H</u> elp |
|---|--------------|
| Document Title: Smart Statement Detail | 7 |
| Document URL: http://payment.openmarket.com/@c632f154cc8021 | 5 |
| Smart Statement Detail | |
| This is the detailed information about a particular transaction from your Smart Statement | |
| Transaction Information | |
| url: http://www.openmarket.com/demos/aug15/mall/mead-fingerprint/mkarticle.cgo transaction_log_id: 50254.0 currency: US transaction_date: 781377633 initiator: 1.0 expiration: 2592000 description: Mead Data Central Article amount: 2.95 beneficiary: 3.0 ip_address: 199.170.183.13 transaction_type.p domain: mead.internet-1 Merchant Information telephone: 617-621-9501 | |
| address_1: Open Market, Inc. address_2: 215 First Street fax: 617-621-1703 | |
| address_3: Cambridge, MA email: testmerchant@openmarket.com principal_name: Test Merchant | |
| | ▽ |
| Back Forward Home Reload Open Save As Clone New Window Close Window | |

FIG. 12

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| File <u>Options N</u> avigate <u>A</u> nnotate | <u>H</u> elp |
|--|--------------|
| Document Title: Smart Statement Detail | |
| Document URL: http://payment.openmarket.com/@c632f154cc8021 | |
| | A |
| <pre>url: http://www.openmarket.com/demos/aug15/mall/mead-fingerprint/mkartic transaction_log_id: 50254.0 currency: US transaction_data: 781377633</pre> | :le.cgo |
| initiator: 1.0 expiration: 2592000 | |
| description: Mead Data Central Article amount: 2.95 | } |
| beneficiary: 3.0 ip_address: 199.170.183.13 transaction_type.p | |
| domain: mead.internet-1 | 11 |
| Merchant Information | |
| telephone: 617-621-9501 address_1: Open Market, Inc. address_2: 215 First Street | |
| <pre>fax: 617-621-1703 address_3: Cambridge, MA email: testmerchant@openmarket.com principal_name: Test Merchant</pre> | |
| home_url: country: US |]] |
| postal_code: 02142 | |
| Feedback | |
| You can send us comments and suggestions here. | |
| | |
| | |
| | |
| | |
| | |
| Back Forward Home Reload Open Save As Clone New Window Close | Window |

FIG. 13

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| <u>File Options Navigate Annotate</u> | <u>H</u> elp |
|---|--------------|
| Document Title: Open Market Feedback Document URL: http://payment.openmarket.com/ben/fsedback.cg | 5 |
| Or if you prefer, you can send your comments via electronic mail to feedback@openmarket.com or via FAX to +1.617.621.1703.If you would like a please include your e-mail address. | reply |
| Your Open Market accound name (optional): | |
| Your E-mail address (optional): | |
| Subject: | |
| Your comments: | |
| | |
| Submit Feedback | ▽ |
| Back Forward Home Reload Open Save As Clone New Window Close Wi | ndow |

FIG. 14

5.715.314

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NETWORK SALES SYSTEM

REFERENCES TO APPENDICES

Microfiche appendices A-G. 4 sheets of 192 images total. are being submitted with the present application.

A claim of copyright is hereby made by Open Market, Incorporated with respect to the software code contained in the microfiche appendices, as of the date of first issuance of a U.S. patent based on this application. The copyright owner has no objection to the facsimile reproduction by anyone of the microfiche appendices as they appear in the Patent and Trademark office patent file or records, but reserves all other copyright rights whatsoever.

This invention relates to user-interactive network sales systems for implementing an open marketplace for goods or services over computer networks such as the Internet.

U.S. patent application Ser. No. 08/168.519, filed Dec. 16, 1993 by David K. Gifford and entitled "Digital Active Advertising." the entire disclosure of which is hereby incorporated herein in its entirety by reference, now abandoned. describes a network sales system that includes a plurality of buyer computers, a plurality of merchant computers, and a payment computer. A user at a buyer computer asks to have advertisements displayed, and the buyer computer requests 25 advertisements from a merchant computer, which sends the advertisements to the buyer computer. The user then requests purchase of an advertised product, and the buyer computer sends a purchase message to the merchant computer. The merchant computer constructs a payment order 30 that it sends to the payment computer, which authorizes the purchase and sends an authorization message to the merchant computer. When the merchant computer receives the authorization message it sends the product to the buyer

The above-mentioned patent application also describes an alternative implementation of the network sales system in which, when the user requests purchase of an advertised product, the buyer computer sends a payment order directly to the payment computer, which sends an authorization message back to the buyer computer that includes an unforgeable certificate that the payment order is valid. The buyer computer then constructs a purchase message that includes the unforgeable certificate and sends it to the merchant computer. When the merchant computer receives the purchase request it sends the product to the buyer computer, based upon the pre-authorized payment order.

SUMMARY OF THE INVENTION

In one aspect, the invention provides a network-based 50 sales system that includes at least one buyer computer for operation by a user desiring to buy a product, at least one merchant computer, and at least one payment computer. The buyer computer, the merchant computer, and the payment computer are interconnected by a computer network. The 55 buyer computer is programmed to receive a user request for purchasing a product, and to cause a payment message to be sent to the payment computer that comprises a product identifier identifying the product. The payment computer is programmed to receive the payment message, to cause an 60 access message to be created that comprises the product identifier and an access message authenticator based on a cryptographic key, and to cause the access message to be sent to the merchant computer. The merchant computer is programmed to receive the access message, to verify the 65 access message authenticator to ensure that the access message authenticator was created using the cryptographic

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key, and to cause the product to be sent to the user desiring to buy the product.

The invention provides a simple design architecture for the network sales system that allows the merchant computer to respond to payment orders from the buyer computer without the merchant computer having to communicate directly with the payment computer to ensure that the user is authorized to purchase the product and without the merchant computer having to store information in a database regarding which buyers are authorized to purchase which products. Rather, when the merchant computer receives an access message from the buyer computer identifying a product to be purchased, the merchant computer need only check the access message to ensure that it was created by the payment computer (thereby establishing for the merchant computer that the buyer is authorized to purchase the product), and then the merchant computer can cause the product to be sent to the buyer computer who has been authorized to purchase the product.

In another aspect, the invention features a network-based sales system that includes at least one buyer computer for operation by a user desiring to buy products, at least one shopping cart computer, and a shopping cart database connected to the shopping cart computer. The buyer computer and the shopping cart computer are interconnected by a computer network. The buyer computer is programmed to receive a plurality of requests from a user to add a plurality of respective products to a shopping cart in the shopping cart database, and, in response to the requests to add the products, to send a plurality of respective shopping cart messages to the shopping cart computer each of which includes a product identifier identifying one of the plurality of products. The shopping cart computer is programmed to receive the plurality of shopping cart messages, to modify the shopping cart in the shopping cart database to reflect the 35 plurality of requests to add the plurality of products to the shopping cart, and to cause a payment message associated with the shopping cart to be created. The buyer computer is programmed to receive a request from the user to purchase the plurality of products added to the shopping cart and to cause the payment message to be activated to initiate a payment transaction for the plurality of products added to the shopping cart.

In another aspect, the invention features a network-based link message system that includes at least one client computer for operation by a client user and at least one server computer for operation by a server user. The client computer and the server computer are interconnected by a computer network. The client computer is programmed to send an initial link message to the server computer. The server computer is programmed to receive the initial link message from the client computer and to create, based on information contained in the initial link message, a session link message that encodes a state of interaction between the client computer and the server computer. The session link message includes a session link authenticator, computed by a cryptographic function of the session link contents, for authenticating the session link message. The server computer is programmed to cause the session link message to be sent to the client computer. The client computer is programmed to cause the session link message to be sent to a computer in the network that is programmed to authenticate the session link message by examining the session link authenticator and that is programmed to respond to the session link message based on the state of the interaction between the client computer and the server computer.

In another aspect, the invention features a network-based sales system that includes a merchant database having a 3

plurality of digital advertisements and a plurality of respective product fulfillment items, at least one creation computer for creating the merchant database, and at least one merchant computer for causing the digital advertisements to be transmitted to a user and for causing advertised products to be transmitted to the user. The creation computer and the merchant computer are interconnected by a computer network. The creation computer is programmed to create the merchant database, and to transmit the digital advertisements and the product fulfillment items to the merchant 10 computer. The merchant computer is programmed to receive the digital advertisements and product fulfillment items, to receive a request for a digital advertisement from a user, to cause the digital advertisement to be sent to the user, to receive from the user an access message identifying an 15 that the buyer computer creates in FIG. 2. advertised product, and to cause the product to be sent to the user in accordance with a product fulfillment item corresponding to the product.

In another aspect, the invention features a hypertext statement system that includes a client computer for opera- 20 tion by a client user and one or more server computers for operation by a server user. The client computer and the server computers are interconnected by a computer network. At least one of the server computers is programmed to record purchase transaction records in a database. Each of 25 the purchase transaction records includes a product description. The server computer is programmed to transmit a statement document that includes the purchase transaction records to the client computer. The client computer is programmed to display the product descriptions, to receive 30 a request from the client user to display a product corresponding to a product description displayed by the client computer, and to cause a product hypertext link derived from a purchase transaction record to be activated. At least one of the server computers is programmed to respond to activation 35 of the product hypertext link by causing the product to be sent to the client computer.

In another aspect, the invention features a network payment system that includes at least one buyer computer for operation by a user desiring to buy a product and at least one 40 payment computer for processing payment messages from the buyer computer. The buyer computer and the payment computer are interconnected by a computer network. The buyer computer is programmed to cause a payment message to be sent to the payment computer. The payment message 45 the Internet. includes a product identifier identifying the product that the user desires to buy. The payment computer is programmed to receive the payment message, to cause an access message to be created to enable the user to access the product, and to record a purchase transaction record in the settlement data- 50 base. The buyer computer is programmed to cause a request for purchase transaction records to be sent to the payment computer. The payment computer is programmed to receive the request for purchase transaction records and to cause a be sent to the buyer computer.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 is a block diagram of a network sales system in accordance with the present invention.

FIG. 2 (2-A through 2-I) is a flowchart diagram illustrating the operation of a purchase transaction in the network sales system of FIG. 1.

FIG. 3 (3-A through 3-B) is a flowchart diagram illustrating the use of a shopping cart for the purchase of 65 products in connection with the network sales system of FIG. 1.

FIG. 4 (4-A through 4-C) is a flowchart diagram illustrating the operation of a smart statement in the network sales system of FIG. 1.

FIG. 5 is a screen snapshot of an advertising document that the merchant computer sends to the buyer computer in

FIG. 6 is a screen snapshot of a confirmation document that the payment computer sends to the buyer computer in FIG. 2.

FIG. 7 is a screen snapshot of a new account document that the payment computer sends to the buyer computer in

FIG. 8 is a screen snapshot of an account name prompt

FIG. 9 is a screen snapshot of a document that the payment computer sends to the buyer computer in FIG. 2 and that provides an option either to repurchase or to use a previously purchased access.

FIG. 10 is a screen snapshot of a fulfillment document that the merchant computer sends to the buyer computer in FIG.

FIG. 11 is a screen snapshot of a smart statement document that the payment computer sends to the buyer computer in FIG. 4.

FIGS. 12 and 13 are screen snapshots of a transaction detail document that the payment computer sends to the buyer computer in FIG. 4.

FIG. 14 is a screen snapshot of a customer service form that the payment computer sends to the buyer computer in FIG. 4.

DETAILED DESCRIPTION

With reference to FIG. 1. a network sales system in accordance with the present invention includes a buyer computer 12 operated by a user desiring to buy a product, a merchant computer 14, which may be operated by a merchant willing to sell products to the buyer or by a manager of the network sales system, a payment computer 16 typically operated by a manager of the network sales system. and a creation computer 20 typically operated by the merchant. The buyer, merchant, payment, and creation computers are all inter-connected by a computer network 10 such as

Creation computer 20 is programmed to build a "store" of products for the merchant. A printout of a computer program for use in creating such a "store" in accordance with the present invention is provided as Appendix F.

The products advertised by merchant computer 14 may be, for example, newspaper or newsletter articles available for purchase by buyers. Creation computer 20 creates a digital advertisement database 18 that stores advertising documents (which may for example be in the form of document derived from the purchase transaction records to 55 summaries of newspaper or newsletter articles, accompanied by prices) and product fulfillment items (which may be the products themselves if the products can be transmitted over the network, or which may be hard goods identifiers if the products are hard goods, i.e., durable products as opposed to 60 information products). Creation computer 20 transmits contents of the advertising document database 18 to merchant computer 14 to enable the merchant computer to cause advertisements and products to be sent to buyers. Merchant computer 14 maintains advertising documents locally in advertising document database 15. In an alternative embodiment, the creation computer does not have a local digital advertisement database, but instead updates a remote

advertising document database on a merchant computer. These updates can be accomplished using HTML forms or other remote database technologies as is understood by practitioners of the art.

Payment computer 16 has access to a settlement database 22 in which payment computer 16 can record details of purchase transactions. The products may be organized into various "domains" of products, and payment computer 16 can access settlement database 22 to record and retrieve records of purchases of products falling within the various 10 domains. Payment computer 16 also has access to a shopping cart database 21 in which a "shopping cart" of products that a user wishes to purchase can be maintained as the user shops prior to actual purchase of the contents of the shopping cart.

With reference to FIG. 2, a purchase transaction begins when a user at buyer computer 12 requests advertisements (step 24) and buyer computer 12 accordingly sends an advertising document URL (universal resource locator) to merchant computer 14 (step 26). The merchant computer 20 fetches an advertising document from the advertising document database (step 28) and sends it to the buyer computer (step 30). An example of an advertising document is shown in FIG. 5. Details of URLs and how they are used are found in the microfiche Appendix G.

The user browses through the advertising document and eventually requests a product (step 32). This results in the buyer computer sending payment URL A to the payment computer (step 34). Payment URL A includes a product 30 identifier that represents the product the user wishes to buy, a domain identifier that represents a domain of products to which the desired product belongs, a payment amount that represents the price of the product, a merchant computer identifier that represents merchant computer 14, a merchant account identifier that represents the particular merchant account to be credited with the payment amount, a duration time that represents the length of time for which access to the product is to be granted to the user after completion of the purchase transaction, an expiration time that represents 40 a deadline beyond which this particular payment URL cannot be used, a buyer network address, and a payment URL authenticator that is a digital signature based on a cryptographic key. The payment URL authenticator is a hash of other information in the payment URL, the hash being 45 defined by a key shared by the merchant and the operator of the payment computer.

In an alternative embodiment, step 34 consists of the buyer computer sending a purchase product message to the merchant computer, and the merchant computer provides 50 payment VRL A to the buyer computer in response to the purchase product message. In this alternative embodiment, payment URL A contains the same contents as above. The buyer computer then sends the payment URL A it has puter.

When the payment computer receives the payment URL it verifies whether the payment URL authenticator was created from the contents of the payment URL using the cryptographic key (step 36). If not, the payment computer 60 sends a document to the buyer computer indicating that access to the network sales system is denied (step 38). Otherwise, the payment computer determines whether the expiration time has past (step 40). If it has, the payment computer sends a document to the buyer computer indicat- 65 ing that the time has expired (step 41). Otherwise, the payment computer checks the buyer computer network

address to see if it matches the one specified in the payment URL (step 42). If it does not match, the payment computer sends a document to the buyer computer indicating that access to the network payment system is denied (step 43). Otherwise, the payment computer sends a payment confirmation document to the buyer computer, the payment confirmation document including an "open" link and a "continue" link (step 44).

An example of a confirmation document is shown in FIG. 6. The confirmation document asks the user to click on a "continue" button if the user already has an account with the payment computer, or to click on an "open" button if the user does not already have an account and wishes to open one.

If the user clicks on the "open" button (step 46), the buyer computer sends payment URL C to the payment computer (step 48), payment URL C being similar to payment URL A but also indicating that the user does not yet have an account. The payment computer creates a new account document (step 50) and sends it to the buyer computer (step 52). An example of a new account document is shown in FIG. 7. When the user receives the new account document he enters the new account name, an account password, a credit card number, the credit card expiration date, and security information such as the maiden name of the user's mother (step 54), and presses a "submit" button (not shown in FIG. 7). The buyer computer sends the new account information to the payment computer (step 56), which enters the new account in the settlement database (step 58).

If the user clicks on the "continue" button (step 60), the buyer computer sends payment URL B to the payment computer (step 62), payment URL B being similar to payment URL A but also indicating that the user already has an account. The payment computer then instructs the buyer computer to provide the account name and password (steps 64 and 66), and the buyer computer prompts the user for this information by creating an account name prompt (example shown in FIG. 8) and a similar password prompt. The user enters the information (step 68) and the buyer computer sends the account name and password to the payment computer (step 70).

The payment computer verifies whether the user name and password are correct (step 72). If they are not correct, the payment computer sends a document to the buyer computer indicating that access to the network sales system is denied (step 74). Otherwise, the payment computer determines whether additional security is warranted, based on, e.g., whether the payment amount exceeds a threshold (step 73). If additional security is warranted, the payment computer creates a challenge form document and sends it to the buyer computer (step 75). The user enters the security information (step 77), the buyer computer sends the security information to the payment computer (step 79), and the received from the merchant computer to the payment com- 55 payment computer determines whether the security information is correct (step 81). If it is not correct, the payment computer sends a document to the buyer computer indicating that access to the network sales system is denied (step

> If the security information is correct, or if additional security was not warranted, the payment computer checks the settlement database to determine whether the user has unexpired access to the domain identifier contained in the payment URL (step 82). If so, the payment computer sends to the buyer computer a document providing an option either to repurchase or to use the previously purchased access (step 84). An example of such a document is shown in FIG. 9. The

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user can respond to the recent purchase query document by choosing to access the previously purchased document (step 85) or to go ahead and buy the currently selected product (step 86).

If the user chooses to access the previously purchased document, the buyer computer skips to step 92 (see below). If the user chooses to buy the currently selected product, the payment computer calculates an actual payment amount that may differ from the payment amount contained in the payment URL (step 87). For example, the purchase of a product in a certain domain may entitle the user to access other products in the domain for free or for a reduced price for a given period of time.

The payment computer then verifies whether the user account has sufficient funds or credit (step 76). If not, the payment computer sends a document to the buyer computer indicating that the user account has insufficient funds (step 78). Otherwise, the payment computer creates an access URL (step 80) that includes a merchant computer identifier. a domain identifier. a product identifier, an indication of the end of the duration time for which access to the product is to be granted, the buyer network address, and an access URL authenticator that is a digital signature based on a cryptographic key. The access URL authenticator is a hash of other information in the access URL, the hash being defined by a key shared by the merchant and the operator of the payment computer. The payment computer then records the product identifier, the domain, the user account, the merchant account, the end of duration time, and the actual payment amount in the settlement database (step 88).

The payment computer then sends a redirect to access URL to the buyer computer (step 90), which sends the access URL to the merchant computer (step 92). The merchant computer verifies whether the access URL authenticator was created from the contents of the access URL using the cryptographic key (step 94). If not, the merchant computer sends a document to the buyer computer indicating that access to the product is denied (step 96).

Otherwise, the merchant computer verifies whether the duration time for access to the product has expired (step 98). This is done because the buyer computer can request access to a purchased product repeatedly. If the duration time has expired, the merchant computer sends a document to the buyer computer indicating that the time has expired (step 100). Otherwise the merchant computer verifies that the buyer computer network address is the same as the buyer network address in the access URL (step 101), and if so, sends a fulfillment document to the buyer computer (step 102), which is displayed by the buyer computer (step 104). An example of a fulfillment document is shown in FIG. 10. Otherwise, the merchant computer sends a document to the buyer computer indicating that access is not allowed (step 103).

With reference now to FIG. 3, when the merchant computer sends the advertising document to the buyer computer, the user may request that a product be added to a shopping cart in the shopping cart database rather than request that the product be purchased immediately. The buyer computer sends a shopping cart URL to the payment computer (step 108), the shopping cart URL including a product identifier, a domain identifier, a payment amount, a merchant computer identifier, a merchant account identifier, a duration time, an expiration time, and a shopping cart URL authenticator that is a digital signature based on a cryptographic key. The shopping cart URL authenticator is a hash of other information in the shopping cart URL, the hash being defined by

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a key shared by the merchant and the operator of the payment computer.

The payment computer verifies whether the shopping cart URL authenticator was created from the contents of the shopping cart URL using a cryptographic key (step 116). If not, the payment computer sends a document to the buyer computer indicating that access to the network sales system is denied (step 112). Otherwise, before any modification to a user's shopping cart is allowed, user authentication is performed (step 113) in a manner analogous to steps 40-81. Once the user is authenticated, the payment computer creates or updates a payment URL for the shopping cart (step 114).

The user then either requests more advertisements (step 24 in FIG. 2) and possibly adds another product to the shopping cart, requests display of the shopping cart (step 116), or requests purchase of the entire contents of the shopping cart (step 124). If the user requests display of the shopping cart (step 116), the buyer computer sends a fetch shopping cart request to the payment computer (step 118). and the payment computer and buyer computer (step 119) perform steps analogous to steps 64-81. The payment computer returns the contents of the shopping cart to the buyer computer (step 120), which displays the contents of the shopping cart (step 122). If the user requests that the entire contents of the shopping cart be purchased (step 124) the buyer computer causes the payment URL for the shopping cart to be activated (step 126) and the payment URL is processed in a manner analogous to the processing of payment URLs for individual products (beginning with step 36 in FIG. 2).

With reference now to FIG. 4, a user can request display of a "smart statement" that lists purchase transactions for a given month (step 128). When the buyer computer receives such a request, it sends a smart statement URL to the payment computer (step 130).

When the payment computer receives the smart statement URL. it verifies whether the smart statement URL authenticator was created from the contents of the smart statement URL using a cryptographic key (step 132). If not, the payment computer sends a document to the buyer computer indicating that access is denied (step 134). Otherwise, the payment computer checks to determine whether the buyer network address in the smart statement URL matches the buyer computer's actual network address (step 136). If not, the payment computer sends a document to the buyer computer indicating that access is denied (step 138). Otherwise (step 140), the payment computer and buyer computer perform a set of steps analogous to steps 64-81 in FIG. 2 (payment computer requests account name and password. user provides the requested information, and payment computer verifies the information).

In an alternative embodiment steps 132–138 are omitted. After verification of account information is complete, the payment computer retrieves the requested settlement data from the settlement database, creates a smart statement document for the buyer, and sends the smart statement document to the buyer computer (step 142). An example of a smart statement document is shown in FIG. 11. Each purchase transaction record in the smart statement document includes the data of the transaction, the name of the merchant, an identification of the product, and the payment amount for the product. The smart statement document also includes a transaction detail URL for each purchase transaction (these URLs, or hypertext links, are discussed below and are not shown in FIG. 11). The smart statement document

ment also identifies previous statements that the user may wish to have displayed.

The buyer computer displays the retrieved document (step 144), and the user may request transaction details for a particular transaction listed on the smart statement (step 146). If so, the buyer computer sends a transaction detail URL (or "payment detail URL") to the payment computer (step 148). The transaction detail URL includes a transaction identifier, a buyer network address, and a transaction detail URL authenticator. When the payment computer receives the transaction detail URL, it performs (step 150) a set of steps analogous to steps 132-140 (verification of URL authenticator, buyer network address, and account information). The payment computer then retrieves from the settlement database data corresponding to the payment transaction specified in the transaction detail URL, creates a transaction detail document, and sends it to the buyer computer (step 152).

An example of a transaction detail document is shown in FIGS. 12 and 13. The document displays a number of items of information about the transaction, including the transaction date, end of the duration time ("expiration"), a description of the product, the payment amount, the domain corresponding to the product, an identification of the merchant, and the merchant's address.

The smart statement document and the transaction detail document both include customer service URLs (hypertext links) that allow the user to request customer service (i.e., to send comments and suggestions to the payment computer). When the user requests customer service (step 154), the buyer computer sends the customer service URL to the payment computer (step 156), which creates a customer service form and sends it to the buyer computer (step 158). An example of a customer service form is shown in FIG. 14. The user types comments into the customer service form (step 160), and the buyer computer sends the user's comments to the payment computer (step 162). The payment computer then posts the user comments and sends a thank you document to the buyer computer (step 164).

A user may request display of a product included in the smart statement. When the user requests that the product be displayed (step 166), the buyer computer sends the access URL contained in the smart statement document to the merchant computer (step 168), and the buyer computer and 45 merchant computer perform a set of steps analogous to steps 94-104 in FIG. 2 (authentication of access URL, verification whether duration time has expired, verification of buyer network address, and transmission of fulfillment document to buyer computer).

Whenever the present application states that one computer sends a URL to another computer, it should be understood that in preferred embodiments the URL is sent in a standard HTTP request message, unless a URL message is specified as a redirection in the present application. The request 55 message includes components of the URL as described by the standard HTTP protocol definition. These URL components in the request message allow the server to provide a response appropriate to the URL. The term "URL" as used the present application is an example of a "link," which is a 60 each comprises a universal resource locator. pointer to another document or form (including multimedia documents, hypertext documents including other links, or audio/video documents).

When the present application states that one computer sends a document to another computer, it should be under- 65 stood that in preferred embodiments the document is a success HTTP response message with the document in the

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body of the message. When the present application states that a server sends an account name and password request message to the client, it should be understood that in preferred embodiments the account name and password request message is an unauthorized HTTP response. A client computer sends account name and password information to a server as part of a request message with an authorization

The software architecture underlying the particular pre-10 ferred embodiment is based upon the hypertext conventions of the World Wide Web. Appendix A describes the Hypertext Markup Language (HTML) document format used to represent digital advertisements. Appendix B describes the HTML forms fill out support in Mosaic 2.0. Appendix C is a description of the Hypertext Transfer Protocol (HTTP) between buyer and merchant computers, Appendix D describes how documents are named with Uniform Resource Locators (URLs) in the network of computers, and Appendix E describes the authentication of URLs using digital signatures.

A printout of a computer program for use in creating and operating such a "store" in accordance with the present invention is provided as Appendix F. A printout of a computer program for use in operating other aspects of the network sales system in accordance with the present invention is provided in the microfiche appendix G.

There has been described a new and useful network-based sales system. It is apparent that those skilled in the art may make numerous modifications and departures from the specific embodiments described herein without departing from the spirit and scope of the claimed invention.

What is claimed is:

- A network-based sales system, comprising:
- at least one buyer computer for operation by a user desiring to buy a product.
- at least one merchant computer; and
- at least one payment computer;
- said buyer computer, said merchant computer, and said payment computer being interconnected by a computer network:
- said buyer computer being programmed to receive a user request for purchasing a product, and to cause a payment message to be sent to said payment computer that comprises a product identifier identifying said product;
- said payment computer being programmed to receive said payment message, to cause an access message to be created that comprises said product identifier and an access message authenticator based on a cryptographic key, and to cause said access message to be sent to said merchant computer; and
- said merchant computer being programmed to receive said access message, to verify said access message authenticator to ensure that said access message authenticator was created using said cryptographic key. and to cause said product to be sent to said user desiring to buy said product.
- 2. A network-based sales system in accordance with claim 1. wherein said payment message and said access message
- 3. A network-based sales system in accordance with claim 1, wherein said payment computer is programmed to identify said merchant computer upon receipt of said payment message from said buyer computer.
- 4. A network-based sales system in accordance with claim wherein said access message comprises a buyer network. address.

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- 5. A network-based sales system in accordance with claim 4. wherein:
 - said product can be transmitted from one computer to another; and
 - said merchant computer causes said product to be sent to 5 said user by transmitting said product to said buyer network address only.
- 6. A network-based sales system in accordance with claim 4. wherein said merchant computer is programmed to verify whether said buyer network address in said access message 10 matches the actual network address of said buyer computer.
- 7. A network-based sales system in accordance with claim 1. wherein said payment message comprises a buyer network address.
- 8. A network-based sales system in accordance with claim 7, wherein said payment computer is programmed to verify whether said buyer network address in said payment message matches the actual network address of said buyer computer.
- 9. A network-based sales system in accordance with claim 1. wherein said access message authenticator comprises a 20 cryptographic function of contents of said access message based on said cryptographic key.
- 10. A network-based sales system in accordance with claim 1, wherein said payment computer is programmed to verify said payment message authenticator to ensure that 25 said payment message authenticator was created using said cryptographic key.
- 11. A network-based sales system in accordance with claim 10, wherein said payment message authenticator comprises a cryptographic function of contents of said payment 30 message based on said cryptographic key.
- 12. A network-based sales system in accordance with claim 1, wherein said payment message comprises a payment amount.
- 13. A network-based sales system in accordance with claim 1, wherein said payment message comprises a merchant account identifier that identifies a merchant account.
- A network-based sales system in accordance with claim 1, wherein said buyer computer is programmed to transmit a user account identifier to said payment computer that identifies a user account.
- 15. A network-based sales system in accordance with claim 14, wherein:
 - said payment message comprises a payment amount; and said payment computer is programmed to ensure that said user account has sufficient funds or credit to cover said payment amount.
- 16. A network-based sales system in accordance with claim 14, wherein:
 - a merchant account identifier that identifies a merchant
 - said payment computer is programmed to record said payment amount, said user account, and said merchant account in a settlement database.
- 17. A network-based sales system in accordance with claim 16, wherein:
 - said payment message comprises a domain identifier; and said payment computer is programmed to record said domain identifier and said user account in a settlement 60
- 18. A network-based sales system in accordance with claim 17, wherein said payment computer is programmed to check said settlement database, upon receipt of said payment message, to determine whether said user account has pre- 65 claim 1, wherein: viously purchased a product associated with said domain identifier.

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- 19. A network-based sales system in accordance with claim 18, wherein said payment computer is programmed to determine an actual payment amount for said product identified by said product identifier in said payment message based on whether said user account has previously purchased a product associated with said domain identifier.
- 20. A network-based sales system in accordance with claim 1, wherein said buyer computer is programmed to transmit a user authenticator to said payment computer and said payment computer is programmed to verify said user authenticator.
- 21. A network-based sales system in accordance with claim 20, wherein said user authenticator comprises a pass-
- 22. A network-based sales system in accordance with claim 20, wherein:
 - said buyer computer is programmed to transmit security information to said payment computer;
 - said payment computer is programmed to transmit a challenge form to said buyer computer under a predetermined condition, said challenge form asking for said security information previously transmitted by said buyer computer to said payment computer;
 - said payment computer is programmed to respond to said challenge form by querying said user for said security information and transmitting said security information to said payment computer; and
 - said payment computer is programmed to verify authenticity of said security information.
- 23. A network-based sales system in accordance with claim 22, wherein:
 - said payment message comprises a payment amount; and said predetermined condition comprises receipt of a payment amount in said payment message that exceeds a
- 24. A network-based sales system in accordance with claim 1, wherein said payment message comprises a merchant computer identifier that identifies said merchant com-
- 25. A network-based sales system in accordance with claim 24, wherein said access message comprises said merchant computer identifier.
- 26. A network-based sales system in accordance with claim 1, wherein said payment message comprises a dura-45 tion time that specifies a length of time for which access to said product is to be granted.
- 27. A network-based sales system in accordance with claim 26, wherein said payment computer is programmed to use said duration time to compute an end of duration time said payment message comprises a payment amount and 50 and to cause said end of duration time to be included in said access message.
 - 28. A network-based sales system in accordance with claim 27, wherein said merchant computer is programmed to verify, upon receipt of said access message, that said end of 55 duration time has not past.
 - 29. A network-based sales system in accordance with claim 1, wherein said payment message comprises an expiration time after which said payment message can no longer
 - 30. A network-based sales system in accordance with claim 29, wherein said payment computer is programmed to verify, upon receipt of said payment message, that said expiration time has not past.
 - 31. A network-based sales system in accordance with
 - said payment computer is programmed to cause said access message to be sent to said buyer computer; and

said buyer computer is programmed to cause said access message received from said payment computer to be sent to said merchant computer.

32. A network-based sales system, comprising:

- at least one buyer computer for operation by a user 5 desiring to buy a product;
- at least one merchant computer; and
- at least one payment computer;
- said buyer computer, said merchant computer, and said payment computer being interconnected by a computer 10 network:
- said buyer computer being programmed to receive a user request for purchasing a product, and to cause a payment URL to be sent to said payment computer that comprises a product identifier identifying said product, 15 a payment amount, and a payment URL authenticator comprising a cryptographic function of contents of said payment URL based on a cryptographic key;
- said payment computer being programmed to receive said payment URL, to verify said payment URL authenti- 20 cator to ensure that said payment URL authenticator was created using said cryptographic key, to ensure that said user has sufficient funds or credit to cover said payment amount, to identify said merchant computer operated by said merchant willing to sell said product 25 to said buyer, to cause an access URL to be created that comprises said product identifier and an access URL authenticator comprising a cryptographic function of contents of said access URL based on a cryptographic buyer computer;
- said buyer computer being programmed to cause said access URL received from said payment computer to be
- said access URL, to verify said access URL authenticator to ensure that said access URL authenticator was created using said cryptographic key, and to cause said product to be sent to said user desiring to buy said product.
- 33. A method of operating a payment computer in a computer network comprising at least one buyer computer for operation by a user desiring to buy a product, at least one merchant computer, and at least one payment computer, the method comprising the steps of:
- receiving, at said payment computer, a payment message that said buyer computer has caused to be sent to said payment computer in response to a user request for purchasing a product, said payment message comprising a product identifier identifying said product;
- causing an access message to be created that comprises said product identifier and an access message authenticator based on a cryptographic key; and
- causing said access message to be sent to said merchant 55 computer, said merchant computer being programmed to receive said access message, to verify said access message authenticator to ensure that said access message authenticator was created using said cryptographic key, and to cause said product to be sent to said user 60 desiring to buy said product.
- 34. A network-based sales system, comprising:
- at least one buyer computer for operation by a user desiring to buy products;
- at least one shopping cart computer; and
- a shopping cart database connected to said shopping cart computer;

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- said buyer computer and said shopping cart computer being interconnected by a computer network;
- said buyer computer being programmed to receive a plurality of requests from a user to add a plurality of respective products to a shopping cart in said shopping cart database, and, in response to said requests to add said products, to send a plurality of respective shopping cart messages to said shopping cart computer each of which comprises a product identifier identifying one of said plurality of products;
- said shopping cart computer being programmed to receive said plurality of shopping cart messages, to modify said shopping cart in said shopping cart database to reflect said plurality of requests to add said plurality of products to said shopping cart, and to cause a payment message associated with said shopping cart to be cre-
- said buyer computer being programmed to receive a request from said user to purchase said plurality of products added to said shopping cart and to cause said payment message to be activated to initiate a payment transaction for said plurality of products added to said shopping cart;
- said shopping cart being a stored representation of a collection of products, said shopping cart database being a database of stored representations of collections of products, and said shopping cart computer being a computer that modifies said stored representations of collections of products in said database.
- 35. A network-based sales system in accordance with key, and to cause said access URL to be sent to said 30 claim 34, wherein said shopping cart computer is programmed to cause said payment message to be created before said buyer computer causes said payment message to be activated.
- 36. A network-based sales system in accordance with said merchant computer being programmed to receive 35 claim 34, wherein said buyer computer is programmed to products added to said shopping cart.
 - 37. A network-based sales system in accordance with claim 36, wherein said buyer computer is programmed to transmit a fetch shopping cart request to said payment computer in response to receipt of said request from said
 - 38. A network-based sales system in accordance with claim 37, wherein:
 - said payment computer is programmed to respond to said fetch shopping cart request by transmitting a message to said buyer computer indicating said plurality of products added to said shopping cart; and
 - said buyer computer is programmed to display said plurality of products added to said shopping cart.
 - 39. A method of operating a shopping cart computer in a computer network comprising at least one buyer computer for operation by a user desiring to buy products, at least one shopping cart computer, and a shopping cart database connected to said shopping cart computer, said method comprising the steps of:
 - receiving, at said shopping cart computer, a plurality of shopping cart messages sent to said shopping cart computer by said buyer computer in response to receipt of a plurality of requests from a user to add a plurality of respective products to a shopping cart in said shopping cart database, each of said shopping cart messages comprising a product identifier identifying one of said plurality of products;
 - modifying said shopping cart in said shopping cart database to reflect said plurality of requests to add said plurality of products to said shopping cart; and

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causing a payment message associated with said shopping cart to be created;

said buyer computer being programmed to receive a request from said user to purchase said plurality of products added to said shopping cart and to cause said payment message to be activated to initiate a payment transaction for said plurality of products added to said shopping cart;

said shopping cart being a stored representation of a collection of products, said shopping cart database being a database of stored representations of collections of products, and said shopping cart computer being a computer that modifies said stored representations of collections of products in said database.

40. A network-based link message system, comprising: at least one client computer for operation by a client user, and

at least one server computer for operation by a server user; said client computer and said server computer being 20 with claim 40, wherein: interconnected by a computer network; said session link authors.

said client computer being programmed to send an initial link message to said server computer;

said server computer being programmed to receive said initial link message from said client computer, to create, based on information contained in said initial link message, a session link message that encodes a state of interaction between said client computer and said server computer, said session link message comprising a session link authenticator, computed by a cryptographic function of said session link contents, for authenticating said session link message, and to cause said session link message to be sent to said client computer;

said client computer being programmed to cause said session link message to be sent to a computer in said network that is programmed to authenticate said session link message by examining said session link authenticator and that is programmed to respond to said session link message based on said state of said interaction between said client computer and said server computer.

41. A network-based link message system in accordance with claim 40, wherein:

said client computer comprises a buyer computer for operation by a user desiring to buy a product;

said server computer comprises a payment computer for operation by a manager of said network-based link message system; and

said network-based link message system further comprises a merchant computer for operation by a merchant willing to sell said product to said buyer.

42. A network-based link message system in accordance with claim 41, wherein said computer that is programmed to

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authenticate said session link message comprises said merchant computer.

43. A network-based link message system in accordance with claim 41, wherein said initial link message comprises a payment message to said payment computer that comprises a product identifier identifying said product.

44. A network-based link message system in accordance with claim 43, wherein said session link message comprises an access message that comprises said product identifier to be created.

45. A network-based link message system in accordance with claim 44, wherein said merchant computer is programmed to respond to said access message by causing said product to be sent to said user desiring to buy said product.

46. A network-based link message system in accordance with claim 40, wherein said initial link message and said session link message comprise universal resource locators.

47. A network-based link message system in accordance

said session link authenticator comprises a cryptographic function of contents of said session link message based on a cryptographic key; and

said computer to which said client computer is programmed to cause said session link message to be sent is programmed to verify that said session link authenticator was created using said cryptographic key.

48. A method of operating a server computer in a networkbased link message system comprising at least one client computer for operation by a client user and at least one server computer for operation by a server user, said client computer and said server computer being interconnected by a computer network, said method comprising the steps of:

receiving, at said server computer, an initial link message sent to said server computer by said client computer;

creating, based on information contained in said initial link message, a session link message that encodes a state of interaction between said client computer and said server computer, said session link message comprising a session link authenticator, computed by a cryptographic function of said session link contents, for authenticating said session link message; and

causing said session link message to be sent to said client computer:

said client computer being programmed to cause said session link message to be sent to a computer in said network that is programmed to authenticate said session link message by examining said session link authenticator and that is programmed to respond to said session link message based on said state of said interaction between said client computer and said server computer.

* * * * *

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EXHIBIT 2



United States Patent [19]

Payne et al.

[11] Patent Number: 5,909,492

[45] Date of Patent:

Jun. 1, 1999

- [75] Inventors: Andrew C. Payne, Lincoln; Lawrence C. Stewart, Burlington, both of Mass.; David J. Mackie, Brookdale, Calif.
- [73] Assignee: Open Market, Incorporated,

Cambridge, Mass.

[21] Appl. No.: 08/878,396

[22] Filed: Jun. 18, 1997

Related U.S. Application Data

- [63] Continuation of application No. 08/328,133, Oct. 24, 1994, Pat. No. 5,715,314.
- [51] Int. Cl. H04L 9/00 [52] U.S. Cl. 380/24; 380/23; 380/25;

380/49; 380/50; 705/26; 705/27; 705/39;

27, 39, 40, 41, 42, 43, 44, 14, 16

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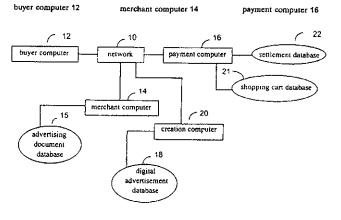
Contents of "Welcome first-time visitors" at www.amazon.com on the Internet as of Jun. 29, 1998.

Primary Examiner—Bernarr E. Gregory Attorney, Agent, or Firm—Fish & Richardson P.C.

[57] ABSTRACT

A network-based sales system includes at least one buyer computer for operation by a user desiring to buy a product, at least one merchant computer, and at least one payment computer. The buyer computer, the merchant computer, and the payment computer are interconnected by a computer network. The buyer computer is programmed to receive a user request for purchasing a product, and to cause a payment message to be sent to the payment computer that comprises a product identifier identifying the product. The payment computer is programmed to receive the payment message, to cause an access message to be created that comprises the product identifier and an access message authenticator based on a cryptographic key, and to cause the access message to be sent to the merchant computer. The merchant computer is programmed to receive the access message, to verify the access message authenticator to ensure that the access message authenticator was created using the cryptographic key, and to cause the product to be sent to the user desiring to buy the product.

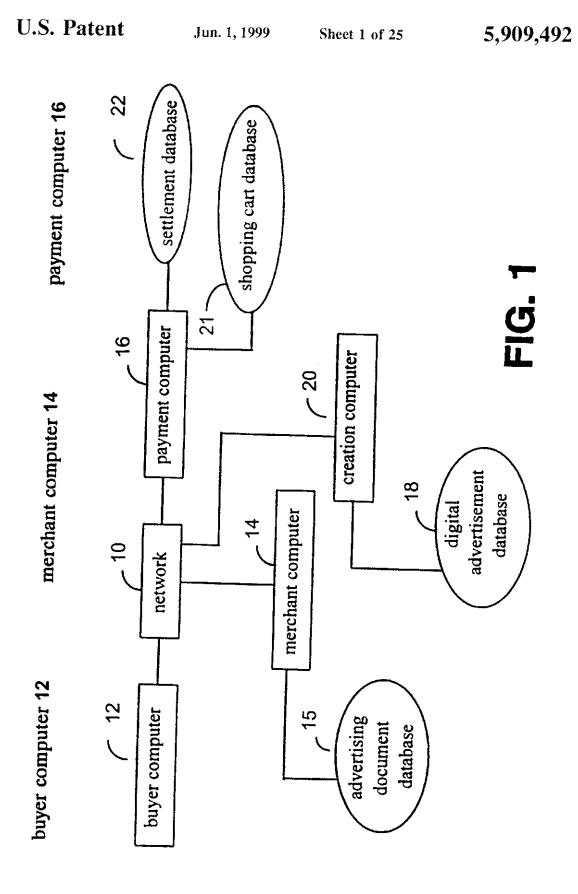
38 Claims, 25 Drawing Sheets





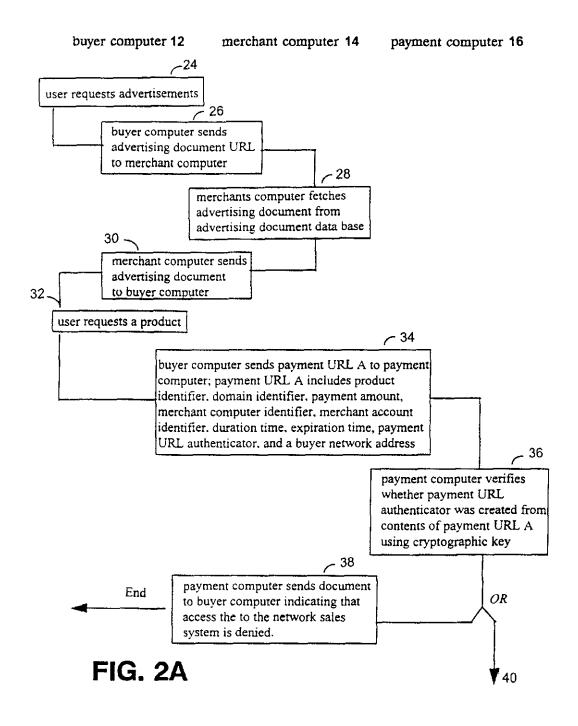
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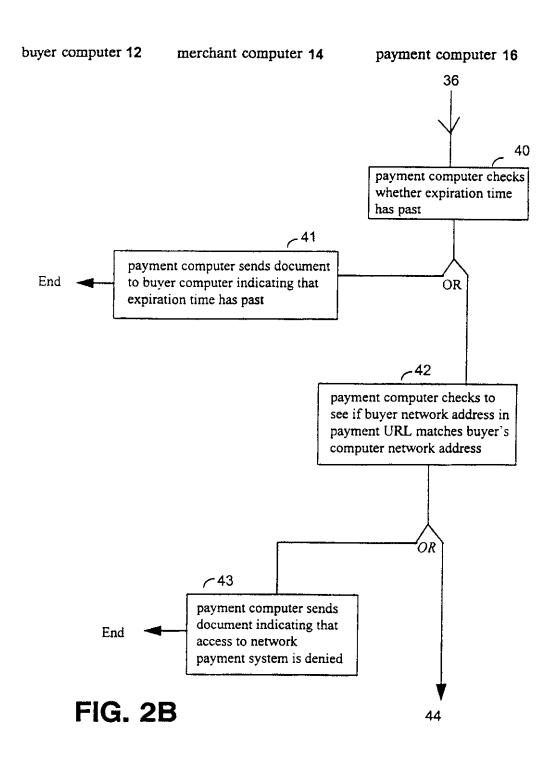
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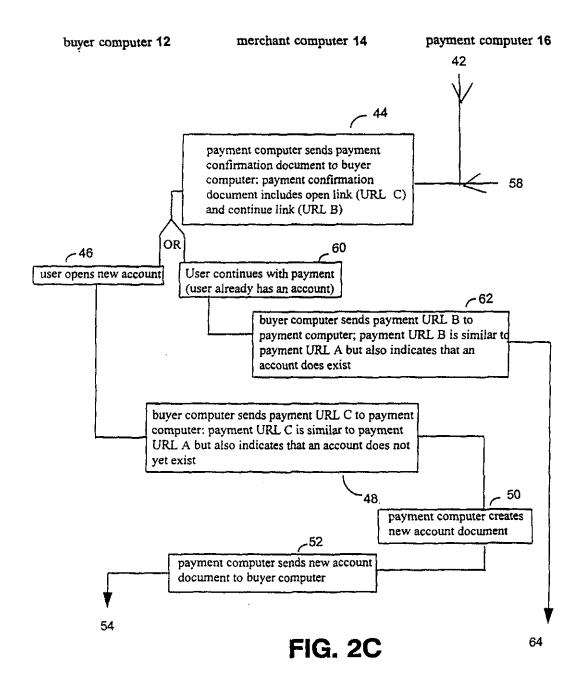
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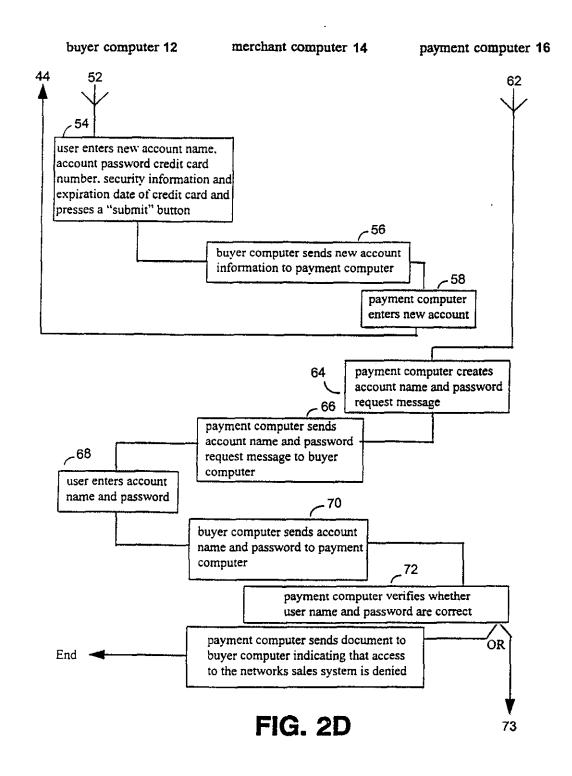
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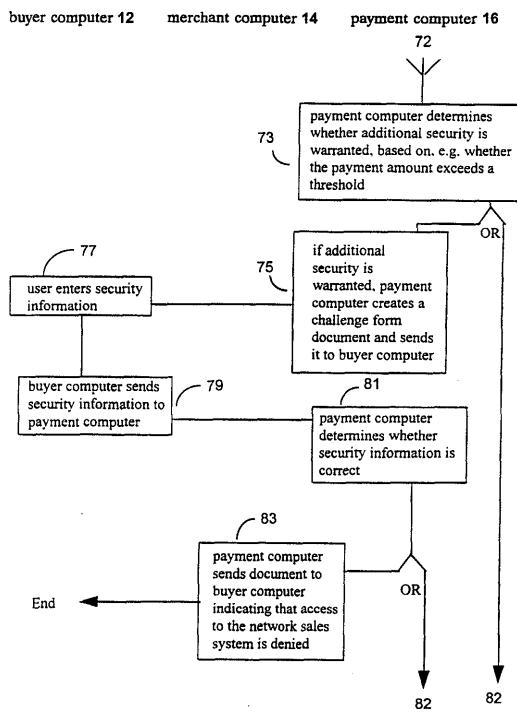


FIG. 2E

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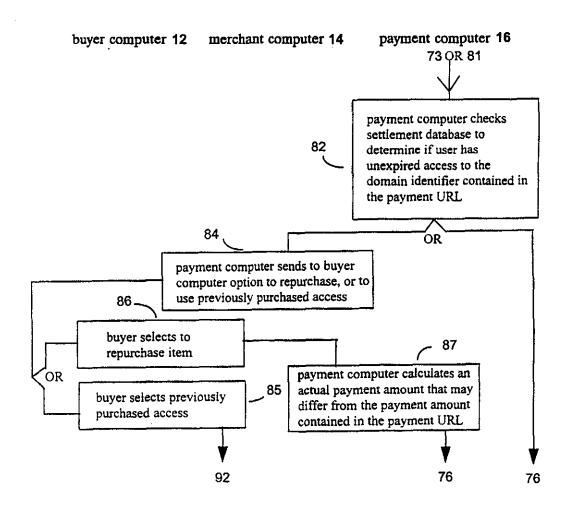
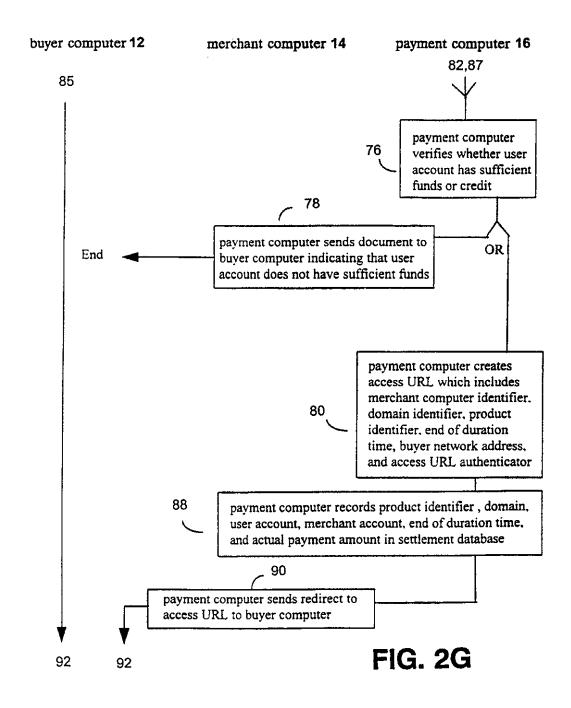


FIG. 2F

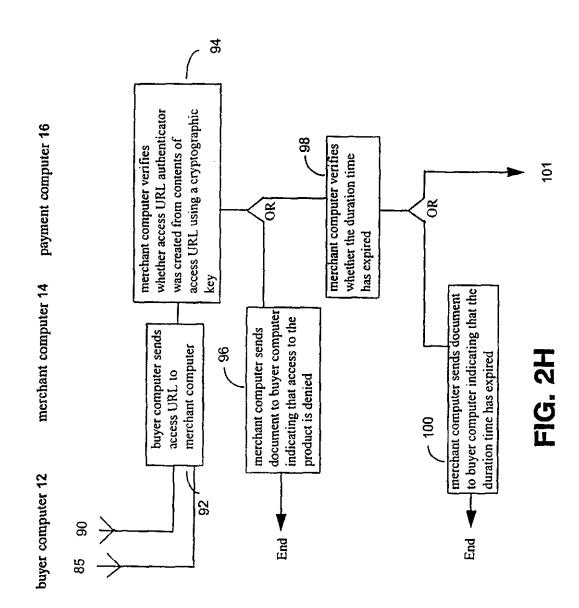
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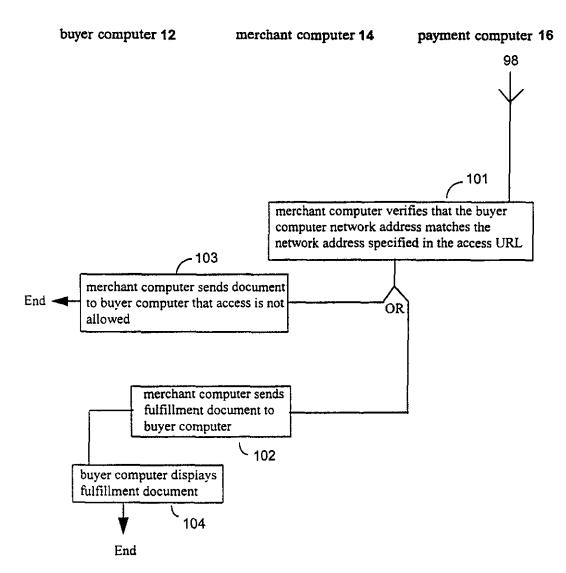
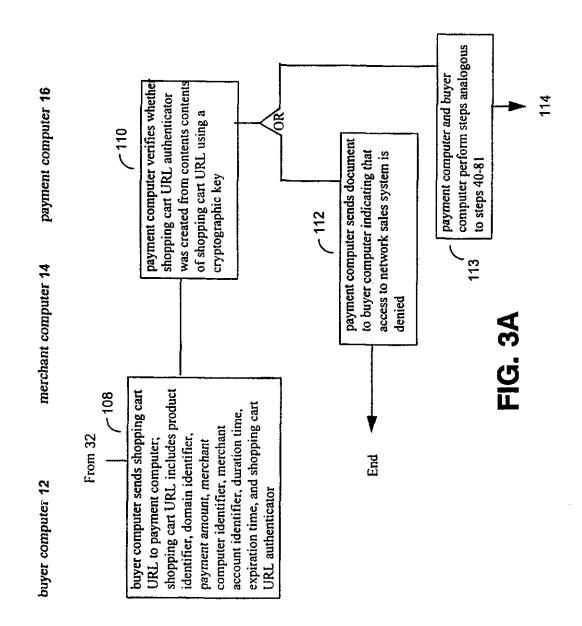


FIG. 21

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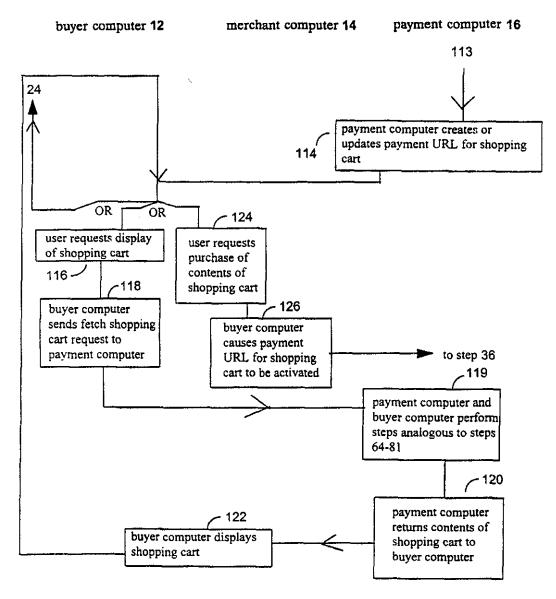
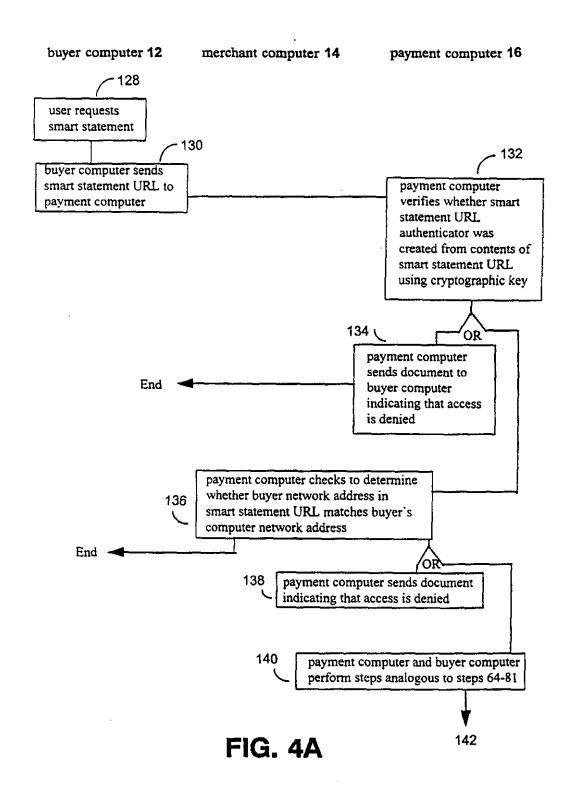


FIG. 3B

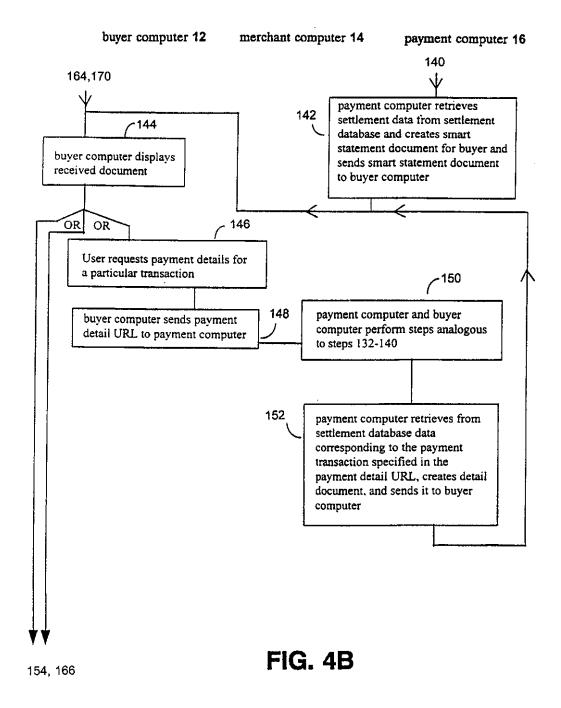
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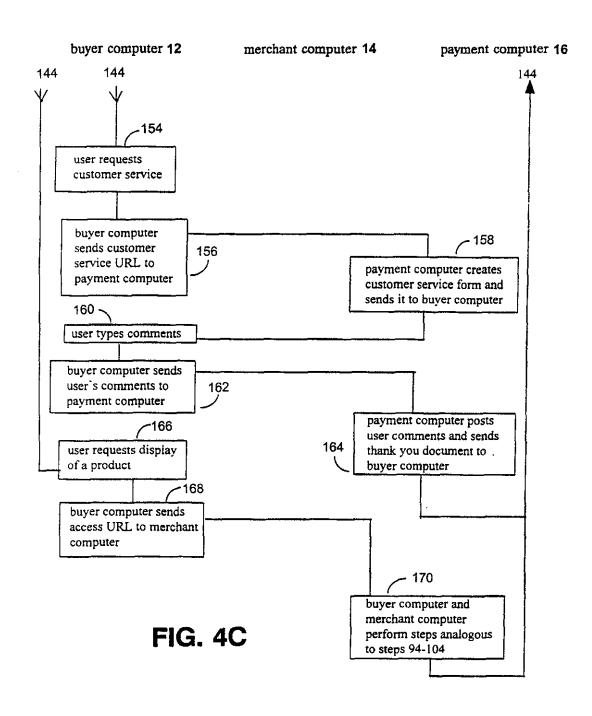
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|--|--------------|
| Document Title: Mead Data Central: Internet Information | |
| Document URL: http://www.openmarketcom/demo/r15/mall/me | |
| Mead Data Central: Internet Information | Δ |
| November 28, 1993 LC's debut on the Internet: Library of Congress catalog On the | |
| Text of Abstract of Article | |
| VERONICA: A GOPHER NAVIGATIONAL TOOL ON THE INTERNET October, 1993 | ⊽ |
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FIG. 5

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| File <u>O</u> ptions | <u>N</u> avigate <u>A</u> nnotate | <u>H</u> el |
|--|---|-------------|
| Document Title: Ope | n Market Payment | |
| Occument URL: htt | p://payment.openmarket.com/ben/nph-payment | |
| Open Market Pay | yment | Δ |
| Merchant:To | ed an item that requires payment est Merchant n:Mead Data Central Article 5(US currency) | |
| you will be prop your account name you can establish | me and password. If you do not have an account, | |
| | account on-line | - |
| NOTE:For demonstrate the password test | ations use the account name testuser@openmarket.comuser. | with |
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FIG. 6

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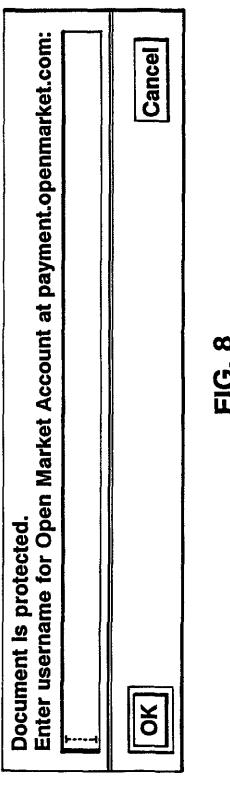
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|--|--------------|
| Document Title: Establish OpenMarket Account Document URL: http://payment.openmarket.com/service/destabli. | 9 |
| Card Number: Expiration Date: (format MM/YY) Check the appropriate boxes: I am the owner of the above credit card. The above address is also the billing address for this credit of the above address is also the billing address for this credit of the copy of your statement is available on-line. At your option you may copy of your statement automatically sent to your e-mail address at weekly or mon intervals. Please choose a statement option. Weekly statements Monthly statements No e-mail statements | y a thly |
| Account name and password Please choose an account name and password for your OpenMarket account. We suggest using an account name that is unique and easy to remember such as ye-mail address. Your password should be 8 characters or longer. Account Name Password | |
| Data transfer complete: Back Forward Home Reload Open Save As Clone New Window Close Window | ow o |

FIG. 7

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|--|----------|
| Document Title: Open Market Payment Document URL: http://payment.openmarket.com/ben/nph-payment | |
| Open Market Payment | Δ |
| You have selected an item that you have purchased recently. Merchant: Test Merchant Description: Mead Data Central Article Amount: 2.85(US currency) This could happen because you would like to buy the item again or it may have happened by accident. | |
| You can: • Go directly to the previous item • Go ahead and buy the item again | |
| Open Market, Inc. | |
| | |
| | ♥ |
| Data transfer complete: Back Forward Home Reload Open Save As Clone New Window Close Window | <u>.</u> |

FIG. 9

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|---|--------------|
| Document Title: LC's debut on the Internet; Library of Congr Document URL: http://www.openmarket.com/@e720f58da6d4ebd268 | |
| LC's debut on the internet Library of Congress catalog | Δ |
| Text of Article | |
| | |
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FIG. 10

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| <u>F</u> ile <u>Options Navigate Annotate</u> | <u>H</u> elp |
|--|---------------|
| Document Title: Smart Statement for Test User | |
| Document URL: http://payment.openmarket.com/in/nph-stateme | |
| Information about the item. Transactions in October 1994 | Δ |
| Mon Oct 3 Test Merchant Dilbert subscription 20 seconds amount Tue Oct 4 Test Merchant Mead Data Central Article amount \$2.95 Tue Oct 4 Test Merchant Mead Data Central Article amount \$2.95 Tue Oct 4 Test Merchant Mead Data Central Article amount \$2.95 Tue Oct 4 Test Merchant Mead Data Central Article amount \$2.95 Tue Oct 4 Test Merchant Mead Data Central Article amount \$2.95 Wed Oct 5 Test Merchant Mead Data Central Article | \$0.10 |
| Previous Statements September 1994 August 1994 | |
| Return to your Newest Statement | |
| Feedback You can send us comments and suggestions here. | |
| Data transfer complete: Back Forward Home Reload Open Save As Clone New Window Close V | Vindow Vindow |

FIG. 11

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| <u>F</u> ile <u>Options Navigate Annotate</u> | <u>H</u> elp |
|---|--------------|
| Document Title: Smart Statement Detail Document URL: http://payment.openmarket.com/@c632f154cc8021 | S |
| Document URL: http://payment.openmarket.com/@c632f154cc8021 | |
| Smart Statement Detail This is the detailed information about a particular transaction from your Smart: | Statement |
| Transaction Information | |
| url: http://www.openmarket.com/demos/aug15/mall/mead-fingerprint/mkarticle transaction_log_id: 50254.0 currency: US transaction_date: 781377633 initiator: 1.0 expiration: 2592000 | . cgo |
| description: Mead Data Central Article amount: 2.95 beneficiary: 3.0 ip_address: 199.170.183.13 transaction_type.p domain: mead.internet-1 | |
| Merchant Information | |
| telephone: 617-621-9501 address_1: Open Market, Inc. address_2: 215 First Street fax: 617-621-1703 address_3: Cambridge, MA email: tastmarchant@openmarket.com | |
| principal_name: Test Merchant | |
| | |
| Data transfer complete: Back Forward Home Reload Open Save As Clone New Window Close | Window |
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FIG. 12

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| <u>F</u> ile | <u>Options</u> | <u>N</u> avigate | <u>A</u> nnotate | | <u>H</u> elp |
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FIG. 14

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NETWORK SALES SYSTEM

CROSS REFERENCE TO RELATED APPLICATION

This is a continuation of U.S. patent application Ser. No. 08/328,133, filed Oct. 24, 1994, now U.S. Pat. No. 5,715, 314.

REFERENCE TO MICROFICHE APPENDICES

Microfiche Appendices A-G are being submitted with the present application, being 4 sheets with 220 total pages.

BACKGROUND OF THE INVENTION

This invention relates to user-interactive network sales 15 systems for implementing an open marketplace for goods or services over computer networks such as the Internet.

U.S. patent application Ser. No. 08/168,519, filed Dec. 16, 1993 by David K. Gifford and entitled "Digital Active Advertising," now abandoned, the entire disclosure of which is hereby incorporated herein in its entirety by reference, describes a network sales system that includes a plurality of buyer computers, a plurality of merchant computers, and a payment computer. A user at a buyer computer asks to have advertisements displayed, and the buyer computer requests advertisements from a merchant computer, which sends the advertisements to the buyer computer. The user then requests purchase of an advertised product, and the buyer computer sends a purchase message to the merchant computer. The merchant computer constructs a payment order that it sends to the payment computer, which authorizes the purchase and sends an authorization message to the merchant computer. When the merchant computer receives the authorization message it sends the product to the buyer 35

The above-mentioned patent application also describes an alternative implementation of the network sales system in which, when the user requests purchase of an advertised product, the buyer computer sends a payment order directly to the payment computer, which sends an authorization message back to the buyer computer that includes an unforgeable certificate that the payment order is valid. The buyer computer then constructs a purchase message that includes the unforgeable certificate and sends it to the merchant computer. When the merchant computer receives the purchase request it sends the product to the buyer computer, based upon the pre-authorized payment order.

SUMMARY OF THE INVENTION

In one aspect, the invention provides a network-based sales system that includes at least one buyer computer for operation by a user desiring to buy a product, at least one merchant computer, and at least one payment computer. The buyer computer, the merchant computer, and the payment 55 computer are interconnected by a computer network. The buyer computer is programmed to receive a user request for purchasing a product, and to cause a payment message to be sent to the payment computer that comprises a product identifier identifying the product. The payment computer is 60 programmed to receive the payment message, to cause an access message to be created that comprises the product identifier and an access message authenticator based on a cryptographic key, and to cause the access message to be sent to the merchant computer. The merchant computer is 65 programmed to receive the access message, to verify the access message authenticator to ensure that the access

message authenticator was created using the cryptographic key, and to cause the product to be sent to the user desiring

to buy the product. The invention provides a simple design architecture for the network sales system that allows the merchant computer to respond to payment orders from the buyer computer without the merchant computer having to communicate directly with the payment computer to ensure that the user is authorized to purchase the product and without the 10 merchant computer having to store information in a database regarding which buyers are authorized to purchase which products. Rather, when the merchant computer receives an access message from the buyer computer identifying a product to be purchased, the merchant computer need only check the access message to ensure that it was created by the payment computer (thereby establishing for the merchant computer that the buyer is authorized to purchase the product), and then the merchant computer can cause the product to be sent to the buyer computer who has been authorized to purchase the product.

In another aspect, the invention features a network-based sales system that includes at least one buyer computer for operation by a user desiring to buy products, at least one shopping cart computer, and a shopping cart database con-25 nected to the shopping cart computer. The buyer computer and the shopping cart computer are interconnected by a computer network. The buyer computer is programmed to receive a plurality of requests from a user to add a plurality of respective products to a shopping cart in the shopping cart database, and, in response to the requests to add the products, to send a plurality of respective shopping cart messages to the shopping cart computer each of which includes a product identifier identifying one of the plurality of products. The shopping cart computer is programmed to receive the plurality of shopping cart messages, to modify the shopping cart in the shopping cart database to reflect the plurality of requests to add the plurality of products to the shopping cart, and to cause a payment message associated with the shopping cart to be created. The buyer computer is programmed to receive a request from the user to purchase the plurality of products added to the shopping cart and to cause the payment message to be activated to initiate a payment transaction for the plurality of products added to the shopping cart.

In another aspect, the invention features a network-based link message system that includes at least one client computer for operation by a client user and at least one server computer for operation by a server user. The client computer and the server computer are interconnected by a computer network. The client computer is programmed to send an initial link message to the server computer. The server computer is programmed to receive the initial link message from the client computer and to create, based on information contained in the initial link message, a session link message that encodes a state of interaction between the client computer and the server computer. The session link message includes a session link authenticator, computed by a cryptographic function of the session link contents, for authenticating the session link message. The server computer is programmed to cause the session link message to be sent to the client computer. The client computer is programmed to cause the session link message to be sent to a computer in the network that is programmed to authenticate the session link message by examining the session link authenticator and that is programmed to respond to the session link message based on the state of the interaction between the client computer and the server computer.

In another aspect, the invention features a network-based sales system that includes a merchant database having a plurality of digital advertisements and a plurality of respective product fulfillment items, at least one creation computer for creating the merchant database, and at least one merchant computer for causing the digital advertisements to be transmitted to a user and for causing advertised products to be transmitted to the user. The creation computer and the merchant computer are interconnected by a computer network. The creation computer is programmed to create the merchant database, and to transmit the digital advertisements and the product fulfillment items to the merchant computer. The merchant computer is programmed to receive the digital advertisements and product fulfillment items, to cause the digital advertisement to be sent to the user, to 15 that the buyer computer creates in FIG. 2. receive a request for a digital advertisement from a user, to receive from the user an access message identifying an advertised product, and to cause the product to be sent to the user in accordance with a product fulfillment item corresponding to the product.

In another aspect, the invention features a hypertext 20 statement system that includes a client computer for operation by a client user and one or more server computers for operation by a server user. The client computer and the server computers are interconnected by a computer network. At least one of the server computers is programmed to 25 record purchase transaction records in a database. Each of the purchase transaction records includes a product description. The server computer is programmed to transmit a statement document that includes the purchase transaction records to the client computer. The client computer is 30 programmed to display the product descriptions, to receive a request from the client user to display a product corresponding to a product description displayed by the client computer, and to cause a product hypertext link derived from a purchase transaction record to be activated. At least one of the server computers is programmed to respond to activation 35 of the product hypertext link by causing the product to be sent to the client computer.

In another aspect, the invention features a network payment system that includes at least one buyer computer for operation by a user desiring to buy a product and at least one 40 payment computer for processing payment messages from the buyer computer. The buyer computer and the payment computer are interconnected by a computer network. The buyer computer is programmed to cause a payment message to be sent to the payment computer. The payment message 45 includes a product identifier identifying the product that the user desires to buy. The payment computer is programmed to receive the payment message, to cause an access message to be created to enable the user to access the product, and to record a purchase transaction record in the settlement data- 50 base. The buyer computer is programmed to cause a request for purchase transaction records to be sent to the payment computer. The payment computer is programmed to receive the request for purchase transaction records and to cause a document derived from the purchase transaction records to 55 be sent to the buyer computer.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 is a block diagram of a network sales system in accordance with the present invention.

FIG. 2 (2-A through 2-I) is a flowchart diagram illustrating the operation of a purchase transaction in the network sales system of FIG. 1.

FIG. 3 (3-A through 3-B) is a flowchart diagram illustrating the use of a shopping cart for the purchase of 65 products in connection with the network sales system of

FIG. 4 (4-A through 4-C) is a flowchart diagram illustrating the operation of a smart statement in the network sales system of FIG. 1.

FIG. 5 is a screen snapshot of an advertising document that the merchant computer sends to the buyer computer in FIG. 2.

FIG. 6 is a screen snapshot of a confirmation document that the payment computer sends to the buyer computer in FIG. 2.

FIG. 7 is a screen snapshot of a new account document that the payment computer sends to the buyer computer in FIG. 2.

FIG. 8 is a screen snapshot of an account name prompt

FIG. 9 is a screen snapshot of a document that the payment computer sends to the buyer computer in FIG. 2 and that provides an option either to repurchase or to use a previously purchased access.

FIG. 10 is a screen snapshot of a fulfillment document that the merchant computer sends to the buyer computer in FIG.

FIG. 11 is a screen snapshot of a smart statement document that the payment computer sends to the buyer computer in FIG. 4.

FIGS. 12 and 13 are screen snapshots of a transaction detail document that the payment computer sends to the buyer computer in FIG. 4.

FIG. 14 is a screen snapshot of a customer service form that the payment computer sends to the buyer computer in FIG. 4.

DETAILED DESCRIPTION

With reference to FIG. 1, a network sales system in accordance with the present invention includes a buyer computer 12 operated by a user desiring to buy a product, a merchant computer 14, which may be operated by a merchant willing to sell products to the buyer or by a manager of the network sales system, a payment computer 16 typically operated by a manager of the network sales system, and a creation computer 20 typically operated by the merchant. The buyer, merchant, payment, and creation computers are all inter-connected by a computer network 10 such as the Internet.

Creation computer 20 is programmed to build a "store" of products for the merchant. A printout of a computer program for use in creating such a "store" in accordance with the present invention is provided as Appendix F.

The products advertised by merchant computer 14 may be, for example, newspaper or newsletter articles available for purchase by buyers. Creation computer 20 creates a digital advertisement database 18 that stores advertising documents (which may for example be in the form of summaries of newspaper or newsletter articles, accompanied by prices) and product fulfillment items (which may be the products themselves if the products can be transmitted over the network, or which may be hard goods identifiers if the products are hard goods, i.e., durable products as opposed to information products). Creation computer 20 transmits contents of the advertising document database 18 to merchant computer 14 to enable the merchant computer to cause advertisements and products to be sent to buyers. Merchant computer 14 maintains advertising documents locally in advertising document database 15. In an alternative embodiment, the creation computer does not have a local digital advertisement database, but instead updates a remote

advertising document database on a merchant computer. These updates can be accomplished using HTML forms or other remote database technologies as is understood by practitioners of the art.

Payment computer 16 has access to a settlement database 22 in which payment computer 16 can record details of purchase transactions. The products may be organized into various "domains" of products, and payment computer 16 can access settlement database 22 to record and retrieve records of purchases of products falling within the various domains. Payment computer 16 also has access to a shopping cart database 21 in which a "shopping cart" of products that a user wishes to purchase can be maintained as the user shops prior to actual purchase of the contents of the shopping cart.

With reference to FIG. 2, a purchase transaction begins when a user at buyer computer 12 requests advertisements (step 24) and buyer computer 12 accordingly sends an advertising document URL (universal resource locator) to merchant computer 14 (step 26). The merchant computer fetches an advertising document from the advertising document database (step 28) and sends it to the buyer computer (step 30). An example of an advertising document is shown in FIG. 5. Details of URLs and how they are used are found in Appendix G.

The user browses through the advertising document and eventually requests a product (step 32). This results in the buyer computer sending payment URL A to the payment computer (step 34). Payment URL A includes a product 30 identifier that represents the product the user wishes to buy, a domain identifier that represents a domain of products to which the desired product belongs, a payment amount that represents the price of the product, a merchant computer identifier that represents merchant computer 14, a merchant account identifier that represents the particular merchant account to be credited with the payment amount, a duration time that represents the length of time for which access to the product is to be granted to the user after completion of the purchase transaction, an expiration time that represents 40 a deadline beyond which this particular payment URL cannot be used, a buyer network address, and a payment URL authenticator that is a digital signature based on a cryptographic key. The payment URL authenticator is a hash of other information in the payment URL, the hash being 45 defined by a key shared by the merchant and the operator of the payment computer.

In an alternative embodiment, step 34 consists of the buyer computer sending a purchase product message to the merchant computer, and the merchant computer provides payment URL A to the buyer computer in response to the purchase product message. In this alternative embodiment, payment URL A contains the same contents as above. The buyer computer then sends the payment URL A it has received from the merchant computer to the payment computer.

When the payment computer receives the payment URL it verifies whether the payment URL authenticator was created from the contents of the payment URL using the cryptographic key (step 36). If not, the payment computer 60 sends a document to the buyer computer indicating that access to the network sales system is denied (step 38). Otherwise, the payment computer determines whether the expiration time has past (step 40). If it has, the payment computer sends a document to the buyer computer indicating that the time has expired (step 41). Otherwise, the payment computer checks the buyer computer network

address to see if it matches the one specified in the payment URL (step 42). If it does not match, the payment computer sends a document to the buyer computer indicating that access to the network payment system is denied (step 43). Otherwise, the payment computer sends a payment confirmation document to the buyer computer, the payment confirmation document including an "open" link and a "continue" link (step 44).

An example of a confirmation document is shown in FIG.

6. The confirmation document asks the user to click on a "continue" button if the user already has an account with the payment computer, or to click on an "open" button if the user does not already have an account and wishes to open one.

If the user clicks on the "open" button (step 46), the buyer computer sends payment URL C to the payment computer (step 48), payment URL C being similar to payment URL A but also indicating that the user does not yet have an account. The payment computer creates a new account document (step 50) and sends it to the buyer computer (step 52). An example of a new account document is shown in FIG. 7. When the user receives the new account document he enters the new account name, an account password, a credit card number, the credit card expiration date, and security information such as the maiden name of the user's mother (step 54), and presses a "submit" button (not shown in FIG. 7). The buyer computer sends the new account information to the payment computer (step 56), which enters the new account in the settlement database (step 58).

If the user clicks on the "continue" button (step 60), the buyer computer sends payment URL B to the payment computer (step 62), payment URL B being similar to payment URLA but also indicating that the user already has an account. The payment computer then instructs the buyer computer to provide the account name and password (steps 64 and 66), and the buyer computer prompts the user for this information by creating an account name prompt (example shown in FIG. 8) and a similar password prompt. The user enters the information (step 68) and the buyer computer sends the account name and password to the payment computer (step 70).

The payment computer verifies whether the user name and password are correct (step 72). If they are not correct, the payment computer sends a document to the buyer computer indicating that access to the network sales system is denied (step 74). Otherwise, the payment computer determines whether additional security is warranted, based on, e.g., whether the payment amount exceeds a threshold (step 73). If additional security is warranted, the payment computer creates a challenge form document and sends it to the buyer computer (step 75). The user enters the security information (step 77), the buyer computer sends the security information to the payment computer (step 79), and the payment computer determines whether the security information is correct (step 81). If it is not correct, the payment computer sends a document to the buyer computer indicating that access to the network sales system is denied (step 83).

If the security information is correct, or if additional security was not warranted, the payment computer checks the settlement database to determine whether the user has unexpired access to the domain identifier contained in the payment URL (step 82). If so, the payment computer sends to the buyer computer a document providing an option either to repurchase or to use the previously purchased access (step 84). An example of such a document is shown in FIG. 9. The

user can respond to the recent purchase query document by choosing to access the previously purchased document (step 85) or to go ahead and buy the currently selected product (step 86).

If the user chooses to access the previously purchased document, the buyer computer skips to step 92 (see below). If the user chooses to buy the currently selected product, the payment computer calculates an actual payment amount that may differ from the payment amount contained in the payment URL (step 87). For example, the purchase of a product in a certain domain may entitle the user to access other products in the domain for free or for a reduced price for a given period of time.

The payment computer then verifies whether the user account has sufficient funds or credit (step 76). If not, the 15 payment computer sends a document to the buyer computer indicating that the user account has insufficient funds (step 78). Otherwise, the payment computer creates an access URL (step 80) that includes a merchant computer identifier, a domain identifier, a product identifier, an indication of the end of the duration time for which access to the product is to be granted, the buyer network address, and an access URL. authenticator that is a digital signature based on a cryptographic key. The access URL authenticator is a hash of other information in the access URL, the hash being defined by a key shared by the merchant and the operator of the payment computer. The payment computer then records the product identifier, the domain, the user account, the merchant account, the end of duration time, and the actual payment amount in the settlement database (step 88).

The payment computer then sends a redirect to access URL to the buyer computer (step 90), which sends the access URL to the merchant computer (step 92). The merchant computer verifies whether the access URL authenticator was created from the contents of the access URL using the cryptographic key (step 94). If not, the merchant computer sends a document to the buyer computer indicating that access to the product is denied (step 96).

Otherwise, the merchant computer verifies whether the duration time for access to the product has expired (step 98). This is done because the buyer computer can request access to a purchased product repeatedly. If the duration time has expired, the merchant computer sends a document to the buyer computer indicating that the time has expired (step 100). Otherwise the merchant computer verifies that the buyer computer network address is the same as the buyer network address in the access URL (step 101), and if so, sends a fulfillment document to the buyer computer (step 102), which is displayed by the buyer computer (step 104). An example of a fulfillment document is shown in FIG. 10. Otherwise, the merchant computer sends a document to the buyer computer indicating that access is not allowed (step 103).

With reference now to FIG. 3, when the merchant computer sends the advertising document to the buyer computer, the user may request that a product be added to a shopping cart in the shopping cart database rather than request that the product be purchased immediately. The buyer computer sends a shopping cart URL to the payment computer (step 60 108), the shopping cart URL including a product identifier, a domain identifier, a payment amount, a merchant computer identifier, a merchant account identifier, a duration time, an expiration time, and a shopping cart URL authenticator that is a digital signature based on a cryptographic key. The shopping cart URL authenticator is a hash of other information in the shopping cart URL, the hash being defined by

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a key shared by the merchant and the operator of the payment computer.

The payment computer verifies whether the shopping cart URL authenticator was created from the contents of the shopping cart URL using a cryptographic key (step 110). If not, the payment computer sends a document to the buyer computer indicating that access to the network sales system is denied (step 112). Otherwise, before any modification to a user's shopping cart is allowed, user authentication is performed (step 113) in a manner analogous to steps 40–81. Once the user is authenticated, the payment computer creates or updates a payment URL for the shopping cart (step 114).

The user then either requests more advertisements (step 24 in FIG. 2) and possibly adds another product to the shopping cart, requests display of the shopping cart (step 116), or requests purchase of the entire contents of the shopping cart (step 124). If the user requests display of the shopping cart (step 116), the buyer computer sends a fetch shopping cart request to the payment computer (step 118), and the payment computer and buyer computer (step 119) perform steps analogous to steps 64-81. The payment computer returns the contents of the shopping cart to the buyer computer (step 120), which displays the contents of the shopping cart (step 122). If the user requests that the entire contents of the shopping cart be purchased (step 124) the buyer computer causes the payment URL for the shopping cart to be activated (step 126) and the payment URL is processed in a manner analogous to the processing of payment URLs for individual products (beginning with step 36 in FIG. 2).

With reference now to FIG. 4, a user can request display of a "smart statement" that lists purchase transactions for a given month (step 128). When the buyer computer receives such a request, it sends a smart statement URL to the payment computer (step 130).

When the payment computer receives the smart statement URL, it verifies whether the smart statement URL authenticator was created from the contents of the smart statement URL using a cryptographic key (step 132). If not, the payment computer sends a document to the buyer computer indicating that access is denied (step 134). Otherwise, the payment computer checks to determine whether the buyer network address in the smart statement URL matches the buyer computer's actual network address (step 136). If not, the payment computer sends a document to the buyer computer indicating that access is denied (step 138). Otherwise (step 140), the payment computer and buyer computer perform a set of steps analogous to steps 64-81 in FIG. 2 (payment computer requests account name and password, user provides the requested information, and payment computer verifies the information).

In an alternative embodiment steps 132-138 are omitted.

After verification of account information is complete, the payment computer retrieves the requested settlement data from the settlement database, creates a smart statement document for the buyer, and sends the smart statement document to the buyer computer (step 142). An example of a smart statement document is shown in FIG. 11. Each purchase transaction record in the smart statement document includes the data of the transaction, the name of the merchant, an identification of the product, and the payment amount for the product. The smart statement document also includes a transaction detail URL for each purchase transaction (these URLs, or hypertext links, are discussed below and are not shown in FIG. 11). The smart statement docu-

ment also identifies previous statements that the user may wish to have displayed.

The buyer computer displays the retrieved document (step 144), and the user may request transaction details for a particular transaction listed on the smart statement (step 146). If so, the buyer computer sends a transaction detail URL (or "payment detail URL") to the payment computer (step 148). The transaction detail URL includes a transaction identifier, a buyer network address, and a transaction detail URL authenticator. When the payment computer receives 10 the transaction detail URL, it performs (step 150) a set of steps analogous to steps 132-140 (verification of URL authenticator, buyer network address, and account information). The payment computer then retrieves from the settlement database data corresponding to the payment 15 transaction specified in the transaction detail URL, creates a transaction detail document, and sends it to the buyer computer (step 152).

An example of a transaction detail document is shown in FIGS. 12 and 13. The document displays a number of items of information about the transaction, including the transaction date, end of the duration time ("expiration"), a description of the product, the payment amount, the domain corresponding to the product, an identification of the merchant, and the merchant's address.

The smart statement document and the transaction detail document both include customer service URLs (hypertext links) that allow the user to request customer service (i.e., to send comments and suggestions to the payment computer). When the user requests customer service (step 154), the buyer computer sends the customer service URL to the payment computer (step 156), which creates a customer service form and sends it to the buyer computer (step 158). An example of a customer service form is shown in FIG. 14. The user types comments into the customer service form (step 160), and the buyer computer sends the user's comments to the payment computer (step 162). The payment computer then posts the user comments and sends a thank you document to the buyer computer (step 164).

A user may request display of a product included in the smart statement. When the user requests that the product be displayed (step 166), the buyer computer sends the access URL contained in the smart statement document to the merchant computer (step 168), and the buyer computer and merchant computer perform a set of steps analogous to steps 94–104 in FIG. 2 (authentication of access URL, verification whether duration time has expired, verification of buyer network address, and transmission of fulfillment document to buyer computer).

Whenever the present application states that one computer sends a URL to another computer, it should be understood that in preferred embodiments the URL is sent in a standard HTTP request message, unless a URL message is specified as a redirection in the present application. The request message includes components of the URL as described by the standard HTTP protocol definition. These URL components in the request message allow the server to provide a response appropriate to the URL. The term "URL" as used the present application is an example of a "link," which is a pointer to another document or form (including multimedia documents, hypertext documents including other links, or audio/video documents).

When the present application states that one computer sends a document to another computer, it should be understood that in preferred embodiments the document is a success HTTP response message with the document in the 10

body of the message. When the present application states that a server sends an account name and password request message to the client, it should be understood that in preferred embodiments the account name and password request message is an unauthorized HTTP response. A client computer sends account name and password information to a server as part of a request message with an authorization field.

The software architecture underlying the particular preferred embodiment is based upon the hypertext conventions of the World Wide Web. Appendix A describes the Hypertext Markup Language (HTML) document format used to represent digital advertisements, Appendix B describes the HTML forms fill out support in Mosaic 2.0, Appendix C is a description of the Hypertext Transfer Protocol (HTTP) between buyer and merchant computers, Appendix D describes how documents are named with Uniform Resource Locators (URLs) in the network of computers, and Appendix E describes the authentication of URLs using digital signatures.

A printout of a computer program for use in creating and operating such a "store" in accordance with the present invention is provided as Appendix F. A printout of a computer program for use in operating other aspects of the network sales system in accordance with the present invention is provided in Appendix G.

There has been described a new and useful network-based sales system. It is apparent that those skilled in the art may make numerous modifications and departures from the specific embodiments described herein without departing from the spirit and scope of the claimed invention.

What is claimed is:

- 1. A network-based sales system, comprising:
- a merchant database comprising a plurality of digital advertisements and a plurality of respective product fulfillment items;
- at least one creation computer for creating said merchant database; and
- at least one merchant computer for causing said digital advertisements to be transmitted to a user and for causing advertised products to be transmitted to said user:
- said creation computers, said merchant computer, and a payment computer being interconnected by a public packet switched computer network;
- said creation computer being programmed to create said merchant database, and to transmit said digital advertisements and said product fulfillment items over said network to said merchant computer;
- said merchant computer being programmed to receive said digital advertisements and product fulfillment items over said network, to receive over said network a request for a digital advertisement from a user, to cause said digital advertisement to be sent to said user over said network, to receive over said network from said user a product request message identifying an advertised product, to receive an access message over said network created by said payment computer, and to cause said product to be sent to said user in accordance with a product fulfillment item corresponding to said product and based upon receipt by the merchant computer of the access message.
- A network-based sales system in accordance with claim
 wherein each of said digital advertisements comprises an abstract of a product and a price.
- 3. A network-based sales system in accordance with claim 2, wherein:

- at least one of said product fulfillment items comprises a product itself; and
- said creation computer is programmed to transmit said product to said merchant computer with said digital advertisements.
- 4. A network-based sales system in accordance with claim 2. wherein:
 - at least one of said product fulfillment items comprises a hard good identifier; and
 - said creation computer is programmed to transmit said 10 hard good identifier to said merchant computer with said digital advertisements.
- 5. A method of operating a merchant computer in a network-based sales system comprising a merchant database that comprises a plurality of digital advertisements and a plurality of respective product fulfillment items, at least one creation computer for creating said merchant database, and at least one merchant computer for causing said digital advertisements to be transmitted to a user and for causing advertised products to be transmitted to said user, and at least one payment computer, said creation computer, said merchant computer, and said payment computer being interconnected by a public packet switched computer network, said method comprising the steps of:
 - receiving, at said merchant computer, said digital advertisements and said product fulfillment items, said digital advertisements and said product fulfillment items having been transmitted over said network to said merchant computer by said creation computer, said merchant database comprising said digital advertisements and said product fulfillment items having been created by said creation computer;

 at least or respon to the form to generally said creation and said product fulfillment items having been created by said creation computer;
 - receiving over said network a request for a digital advertisement from a user;
 - causing said digital advertisement to be sent to said user 35 over said network;
 - receiving over said network from said user a product request message identifying an advertised product;
 - receiving over said network an access message created by said payment computer; and
 - causing said product to be sent to said user in accordance with a product fulfillment item corresponding to said product and based upon receipt by the merchant computer of the access message.
 - 6. A hypertext statement system, comprising:
 - a client computer for operation by a client user; and
 - a plurality of server computers for operation by a server user:
 - said client computer and said server computers being 50 interconnected by a public packet switched computer network;
 - at least one of said server computers being programmed to record information pertaining to purchase transaction records in a database, each of said purchase transaction records comprising a product description, and to cause a statement document comprising said purchase transaction records to be transmitted to said client computer over said network;
 - said client computer being programmed to display said 60 product descriptions, to receive a request from said client user to display a product corresponding to a product description displayed by said client computer, and to cause a product hypertext link derived from a purchase transaction record to be activated;
 - at least one of said server computers, other than a server computer that transmitted said statement document to

- said client computer, being programmed to respond to activation of said product hypertext link by causing said product to be sent to said client computer over said network.
- A hypertext statement system in accordance with claim
 wherein:
 - said client computer is programmed to receive a request from said client user to display transaction details corresponding to a product description displayed by said client computer and to cause a transaction detail hypertext link corresponding to said product description to be activated; and
 - at least one of said server computers is programmed to respond to activation of said transaction detail hypertext link by transmitting said transaction details to said client computer as a transaction detail document.
- 8. A hypertext statement system in accordance with claim 7, wherein:
 - said transaction detail document further comprises a customer service form hypertext link;
 - said client computer is programmed to receive a request from said client user to display a customer service form and to cause said customer service form hypertext link to be activated; and
 - at least one of said server computers is programmed to respond to activation of said customer service form hypertext link by transmitting said customer service form to said client computer.
- 9. A hypertext statement system in accordance with claim 6 wherein
 - said statement document further comprises a customer service form hypertext link;
- said client computer is programmed to receive a request from said client user to display a customer service form and to cause said customer service form hypertext link to be activated; and
- at least one of said server computers is programmed to respond to activation of said customer service form hypertext link by transmitting said customer service form to said client computer.
- 10. A method of operating a server computer in a hypertext statement system comprising a client computer for operation by a client user, and a plurality of server computers for operation by a server user, said client computer and said server computers being interconnected by a public packet switched computer network, said method comprising the steps of:
 - recording, at one of said server computers, information pertaining to purchase transaction records in a database, each of said purchase transaction records comprising a product description; and
 - causing a statement document comprising said purchase transaction records to be transmitted to said client computer over said network;
 - said client computer being programmed to display said product descriptions, to receive a request from said client user to display a product corresponding to a product description displayed by said client computer, and to cause a product hypertext link derived from a purchase transaction record to be activated;
 - at least one of said server computers, other than a server computer that transmitted said statement document to said client computer, being programmed to respond to activation of said product hypertext link by causing said product to be sent to said client computer over said network.

- 11. A network payment system, comprising:
- at least one buyer computer for operation by a user desiring to buy a product; and
- at least one payment computer for processing payment messages from said buyer computer;
- said buyer computer, said payment computer, and a merchant computer being interconnected by a public packet switched computer network;
- said buyer computer being programmed to cause a payment message to be sent to said payment computer over said network;
- said payment computer being programmed to receive said payment message, to cause an access message to be created for transmission over said network to said merchant computer to enable said user to access said product upon verification by said merchant computer that said access message was created by said payment computer, and to record information pertaining to a purchase transaction record in said settlement database; 20
- said buyer computer being programmed to cause a request for a purchase transaction record to be sent to said payment computer over said network; and
- said payment computer being programmed to receive said request for said purchase transaction record and to 25 cause a document derived from said purchase transaction record to be sent to said buyer computer over said network.
- 12. The network payment system of claim 11 wherein the payment message comprises a product identifier identifying 30 the product that the user desires to buy.
- 13. A method of operating a payment computer in a network payment system comprising at least one buyer computer for operation by a user desiring to buy a product, and at least one payment computer for processing payment 35 messages from said buyer computer, and at least one merchant computer, said buyer computer, said payment computer, and said merchant computer being interconnected by a public packet switched computer network, said method comprising the steps of:

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 - receiving, at said payment computer, a payment message that said buyer computer has caused to be sent to said payment computer over said network;
 - causing an access message to be created for transmission to a merchant computer over said network to enable said user to access said product upon verification by said merchant computer that said access message was created by said payment computer;
 - recording information pertaining to a purchase transaction record in said settlement database;
 - receiving over said network a request for a purchase transaction record that said buyer computer has caused to be sent to said payment computer; and
 - causing a document derived from said purchase transaction record to be sent to said buyer computer over said network.
- 14. The method of claim 13 wherein the payment message comprises a product identifier identifying the product that the user desires to buy.
 - 15. A hypertext statement system, comprising: a client computer for operation by a client user; and one or more server computers for operation by a server user:
 - the client computer and the server computers being interconnected by a public packet switched computer network:

- at least one of the server computers being programmed to record information pertaining to purchase transaction records in a database, and to transmit a statement document comprising the purchase transaction records to the client computer over the network;
- the client computer being programmed to display the statement document to receive a request from the client user to display transaction details corresponding to a portion of the statement document displayed by the client computer, and to cause a transaction detail hypertext link corresponding to the portion of the statement document to be activated;
- at least one of the server computers being programmed to respond to activation of the transaction detail hypertext link by transmitting the transaction details to the client computer over the network as a transaction detail document.
- 16. A method of operating a server computer in a hypertext statement system comprising a client computer for operation by a client user, and one or more server computers for operation by a server user, the client computer and the server computers being interconnected by a public packet switched computer network, the method comprising the steps of:
 - recording, at one of the server computers, information pertaining to purchase transaction records in a database; and
 - transmitting a statement document comprising the purchase transaction records to the client computer over the network:
 - the client computer being programmed to display the statement document, to receive a request from the client user to display transaction details corresponding to a portion of the statement document displayed by the client computer, and to cause a transaction detail hypertext link corresponding to the portion of the statement document to be activated;
 - at least one of the server computers being programmed to respond to activation of the transaction detail hypertext link by transmitting the transaction details to the client computer over the network as a transaction detail document.
 - 17. A network-based sales system, comprising:
 - at least one buyer computer for operation by a user desiring to buy products;
 - at least one shopping cart computer; and
 - a shopping cart database connected to the shopping cart computer;
 - the buyer computer and the shopping cart computer being interconnected by a public packet switched computer network;
 - the buyer computer being programmed to receive a plurality of requests from a user to add a plurality of respective products to a shopping cart in the shopping cart database, and, in response to the requests to add the products, to send a plurality of respective shopping cart messages over the network to the shopping cart computer each of which comprises a product identifier identifying one of the plurality of products and at least one of which comprises a universal resource locator;
 - the shopping cart computer being programmed to receive the plurality of shopping cart messages, to modify the shopping cart in the shopping cart database to reflect the plurality of requests to add the plurality of products to the shopping cart, and to cause a payment message

associated with the shopping cart to be created, the payment message comprising a universal resource locator; and

the buyer computer being programmed to receive a request from the user to purchase the plurality of 5 products added to the shopping cart and to cause the payment message to be activated to initiate a payment transaction for the plurality of products added to the shopping cart;

the shopping cart being a stored representation of a 10 collection of products, the shopping cart database being a database of stored representations of collections of products, and the shopping cart computer being a computer that modifies the stored representations of collections of products in the database.

18. A method of operating a shopping cart computer in a public packet switched computer network comprising at least one buyer computer for operation by a user desiring to buy products, at least one shopping cart computer, and a shopping cart database connected to the shopping cart 20 computer, the method comprising the steps of:

receiving, at the shopping cart computer, a plurality of shopping cart messages sent over the network to the shopping cart computer by the buyer computer in response to receipt of a plurality of requests from a user to add a plurality of respective products to a shopping cart in the shopping cart database, each of the shopping cart messages comprising a product identifier identifying one of the plurality of products and at least one of which comprises a universal resource locator;

modifying the shopping cart in the shopping cart database to reflect the plurality of requests to add the plurality of products to the shopping cart; and

causing a payment message associated with the shopping cart to be created, the payment message comprising a 35 universal resource locator;

the buyer computer being programmed to receive a request from the user to purchase the plurality of products added to the shopping cart and to cause the payment message to be activated to initiate a payment transaction for the plurality of products added to the shopping cart;

the shopping cart being a stored representation of a collection of products, the shopping cart database being a database of stored representations of collections of products, and the shopping cart computer being a computer that modifies the stored representations of collections of products in the database.

- 19. A network-based sales system, comprising:
- at least one buyer computer for operation by a user ⁵⁰ desiring to buy a product;
- at least one merchant computer; and
- at least one payment computer;

the buyer computer, the merchant computer, and the 55 payment computer being interconnected by a computer network;

the buyer computer being programmed to receive a user request for purchasing a product, and to cause a payment message to be sent to the payment computer that comprises a product identifier identifying the product;

the payment computer being programmed to receive the payment message, to cause an access message to be created that comprises a product identifier identifying the product and an access message authenticator based on a cryptographic key, and to cause the access message to be sent to the merchant computer; and

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the merchant computer being programmed to receive the access message, to cause the access message authenticator to be verified to ensure that the access message authenticator was created using the cryptographic key, and to cause the product to be received by the user desiring to buy the product.

20. A network-based sales system in accordance with claim 19 wherein the buyer computer is programmed to cause the payment message to be sent to the payment computer by sending a purchase product message to the merchant computer, the merchant computer being programmed to receive the purchase product message, and in response thereto, to send the payment message to the payment computer.

21. A network-based sales system in accordance with claim 19 wherein the merchant computer is programmed itself to verify the access message authenticator.

22. A network-based sales system in accordance with claim 19 wherein the merchant computer is programmed to cause every access message authenticator received by the merchant computer to be verified.

23. A network-based sales system in accordance with claim 19, wherein the payment message comprises a payment amount.

24. A network-based sales system in accordance with claim 19, wherein the payment computer is programmed to record the product identifier and the payment amount.

25. A network-based sales system in accordance with claim 24, wherein the product identifier and the payment amount are recorded in a settlement database.

26. A network-based sales system in accordance with claim 19, wherein the payment message comprises a merchant computer identifier.

27. A network-based sales system in accordance with claim 19, wherein the payment message comprises a payment message authenticator based on a cryptographic key.

28. A network-based sales system in accordance with claim 27, wherein the payment computer is programmed to verify the payment message authenticator to ensure that the payment message authenticator was created using the cryptographic key.

29. A network-based sales system in accordance with claim 19 wherein the computer network is a public packet-switched communications network.

30. A method of operating a payment computer in a computer network comprising at least one buyer computer for operation by a user desiring to buy a product, at least one merchant computer, and at least one payment computer, the method comprising the steps of:

receiving, at the payment computer, a payment message that the buyer computer has caused to be sent to the payment computer in response to a user request for purchasing a product, the payment message comprising a product identifier identifying the product;

causing an access message to be created that comprises a product identifier identifying the product and an access message authenticator based on a cryptographic key; and

causing the access message to be sent to the merchant computer, the merchant computer being programmed to receive the access message, to cause the access message authenticator to be verified to ensure that the access message authenticator was created using the cryptographic key, and to cause the product to be received by the user desiring to buy the product.

31. A network-based sales system, comprising:

at least one buyer computer for operation by a user desiring to buy a product;

at least one merchant computer; and at least one payment computer;

- the buyer computer, the merchant computer, and the payment computer being interconnected by a public packet switched computer network;
- the buyer computer being programmed to receive a request from a user for purchasing a product, and to cause a payment message to be sent over the network to the payment computer;
- the payment computer being programmed to receive the payment message, and, if purchase of the product by the user has not been previously recorded in a settlement database, to cause the user to be charged for the product and to create a new record in the settlement database reflecting purchase of the product by the user, to cause an access message to be created, and to cause the access message to be sent over the network to the merchant computer; and
- the merchant computer being programmed to receive the 20 access message and to cause the user to receive the product.
- 32. The network-based sales system of claim 31 wherein: the payment computer is programmed to cause the access message to be created using a cryptographic key; and 25
- at least one of the computers is programmed to use the access message in a cryptographic process to ensure that the user has paid for the product.
- 33. A method of operating a payment computer in a public packet switched computer network comprising at least one buyer computer for operation by a user desiring to buy a product, at least one merchant computer, and at least one payment computer, the method comprising the steps of:
 - receiving, at the payment computer, a payment message that the buyer computer has caused to be sent over the network to the payment computer in response to a request from a user for purchasing a product, and, if purchase of the product by the user has not been previously recorded in a settlement database, causing the user to be charged for the product and creating a new record in the settlement database reflecting purchase of the product by the user;
 - causing an access message to be created; and
 - causing the access message to be sent over the network to 45 the merchant computer, the merchant computer being programmed to receive the access message and to cause the user to receive the product.
- 34. The method of claim 33 wherein at least one of the computers is programmed to use the access message in a 50 cryptographic process to ensure that the user has paid for the product.
 - 35. A network-based sales system, comprising:
 - at least one buyer computer for operation by a user desiring to buy products;
 - at least one shopping cart computer; and
 - a shopping cart database connected to the shopping cart computer;
 - the buyer computer and the shopping eart computer being 60 interconnected by a public packet switched computer network;
 - the buyer computer being programmed to receive a plurality of requests from a user to add a plurality of respective products to a shopping cart in the shopping 65 cart database, and, in response to the requests to add the products, to send a plurality of respective shopping cart

- messages over the network to the shopping cart computer each of which comprises a product identifier identifying one of the plurality of products;
- the shopping cart computer being programmed to receive the plurality of shopping cart messages, and to modify the shopping cart in the shopping cart database to reflect the plurality of requests to add the plurality of products to the shopping cart; and
- the buyer computer being programmed to receive a request from the user to purchase the plurality of products added to the shopping cart and to cause a payment message to be activated to initiate a payment transaction for the plurality of products added to the shopping cart;
- the shopping cart being a stored representation of a collection of products, the shopping cart database being a database of stored representations of collections of products, and the shopping cart computer being a computer that modifies the stored representations of collections of products in the database.
- 36. A method of operating a shopping cart computer in a public packet switched computer network comprising at least one buyer computer for operation by a user desiring to buy products, at least one shopping cart computer, and a shopping cart database connected to the shopping cart computer, the method comprising the steps of:
 - receiving, at the shopping cart computer, a plurality of shopping cart messages sent over the network to the shopping cart computer by the buyer computer in response to receipt of a plurality of requests from a user to add a plurality of respective products to a shopping cart in the shopping cart database, each of the shopping cart messages comprising a product identifier identifying one of the plurality of products; and
 - modifying the shopping cart in the shopping cart database to reflect the plurality of requests to add the plurality of products to the shopping cart;
 - the buyer computer being programmed to receive a request from the user to purchase the plurality of products added to the shopping cart and to cause a payment message to be activated to initiate a payment transaction for the plurality of products added to the shopping cart;
 - the shopping cart being a stored representation of a collection of products, the shopping cart database being a database of stored representations of collections of products, and the shopping cart computer being a computer that modifies the stored representations of collections of products in the database.
 - 37. A network-based sales system, comprising:
 - a merchant database comprising a plurality of digital advertisements and a plurality of respective product fulfillment items;
 - at least one creation computer for creating the merchant database; and
 - at least one merchant computer for causing the digital advertisements to be transmitted to a user and for causing advertised products to be transmitted to the user:
 - the creation computer and the merchant computer being interconnected by a public packet switched computer network;
 - the creation computer being programmed to create the merchant database, and to transmit the digital advertisements and the product fulfillment items over the network to the merchant computer;

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the merchant computer being programmed to receive the digital advertisements and product fulfillment items over the network, to receive over the network a request for a digital advertisement from a user, to cause the digital advertisement to be sent to the user over the network, to receive over the network from the user a product request message identifying an advertised product, and to cause the product to be sent to the user in accordance with a product fulfillment item corresponding to the product;

at least a portion of the digital advertisements transmitted by the creation computer to the merchant computer over the network being authenticated by at least one digital signature.

38. A method of operating a merchant computer in a 15 network-based sales system comprising a merchant database that comprises a plurality of digital advertisements and a plurality of respective product fulfillment items, at least one creation computer for creating the merchant database, and at least one merchant computer for causing the digital advertisements to be transmitted to a user and for causing advertised products to be transmitted to the user, the creation computer and the merchant computer being interconnected by a public packet switched computer network, the method comprising the steps of:

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receiving, at the merchant computer, the digital advertisements and the product fulfillment items, the digital advertisements and the product fulfillment items having been transmitted over the network to the merchant computer by the creation computer, the merchant database comprising the digital advertisements and the product fulfillment items having been created by the creation computer;

receiving over the network a request for a digital advertisement from a user;

causing the digital advertisement to be sent to the user over the network;

receiving over the network from the user a product request message identifying an advertised product; and causing the product to be sent to the user in accordance with a product fulfillment item corresponding to the product;

at least a portion of the digital advertisements transmitted by the creation computer to the merchant computer over the network being authenticated by at least one digital signature.

* * * * *

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EXHIBIT 3



United States Patent [19]

Levergood et al.

[11] Patent Number:

5,708,780

[45] Date of Patent:

Jan. 13, 1998

[54] INTERNET SERVER ACCESS CONTROL AND MONITORING SYSTEMS

[75] Inventors: Thomas Mark Levergood. Hopkinton; Lawrence C. Stewart. Burlington; Stephen Jeffrey Morris. Westford; Andrew C. Payne. Lincoln; George Winfield Treese, Newton, all of Mass.

[73] Assignee: Open Market, Inc., Cambridge, Mass.

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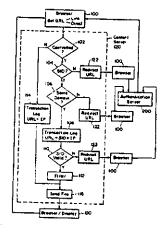
(List continued on next page.)

Primary Examiner—William M. Treat
Assistant Examiner—Patrice L. Winder
Attorney, Agent, or Firm—Hamilton. Brook. Smith &
Reynolds, P.C.

[57] ABSTRACT

This invention relates to methods for controlling and monitoring access to network servers. In particular, the process described in the invention includes client-server sessions over the Internet involving hypertext files. In the hypertext environment, a client views a document transmitted by a content server with a standard program known as the browser. Each hypertext document or page contains links to other hypertext pages which the user may select to traverse. When the user selects a link that is directed to an access-controlled file, the server subjects the request to a secondary server which determines whether the client has an authorization or valid account. Upon such verification, the user is provided with a session identification which allows the user to access to the requested file as well as any other files within the present protection domain.

45 Claims, 7 Drawing Sheets



PLAINTIFF'S EXHIBIT

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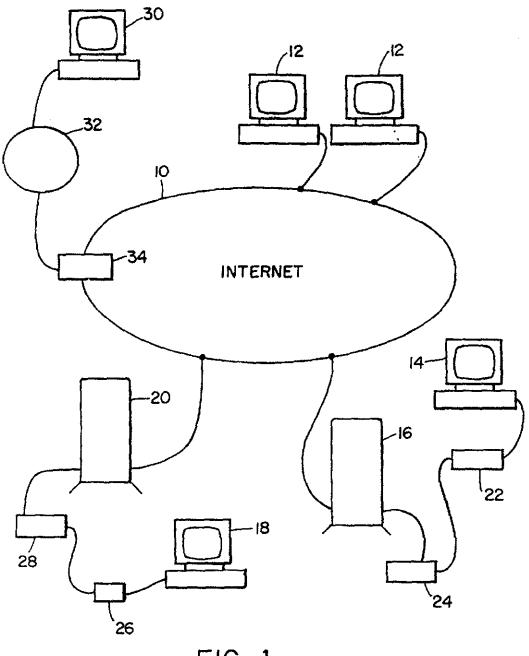


FIG. 1

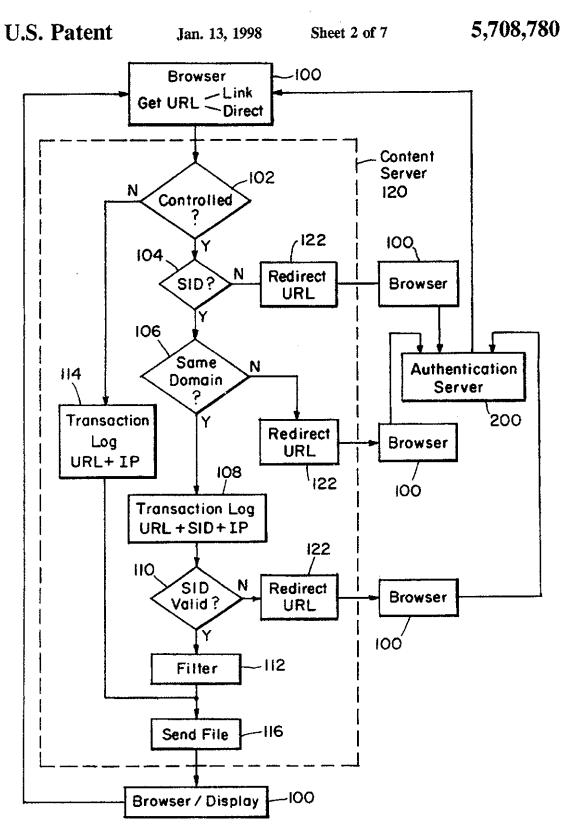
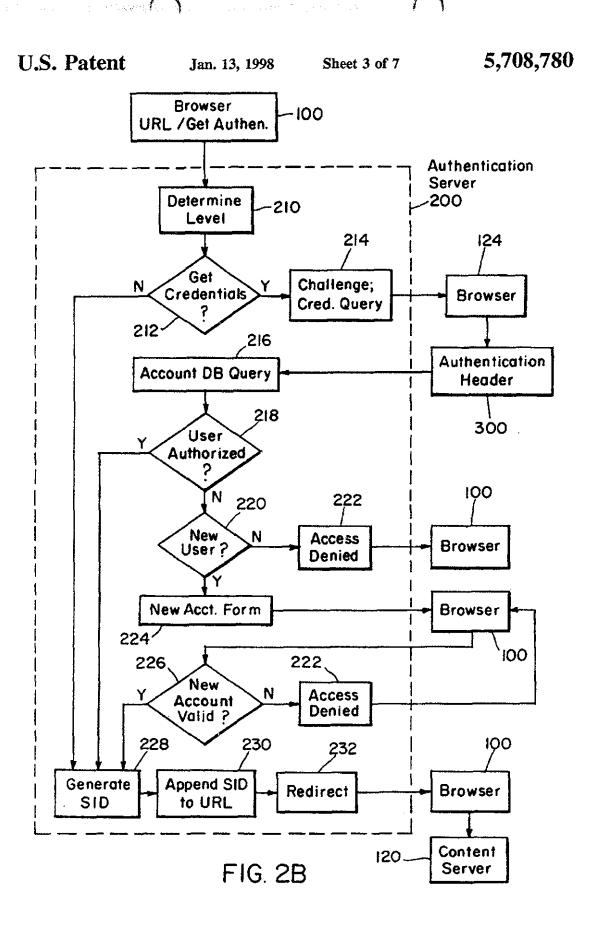
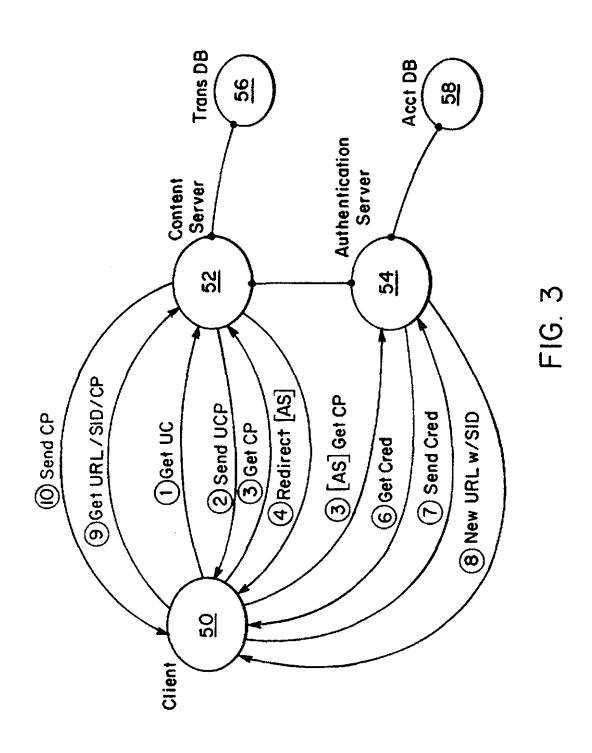


FIG. 2A



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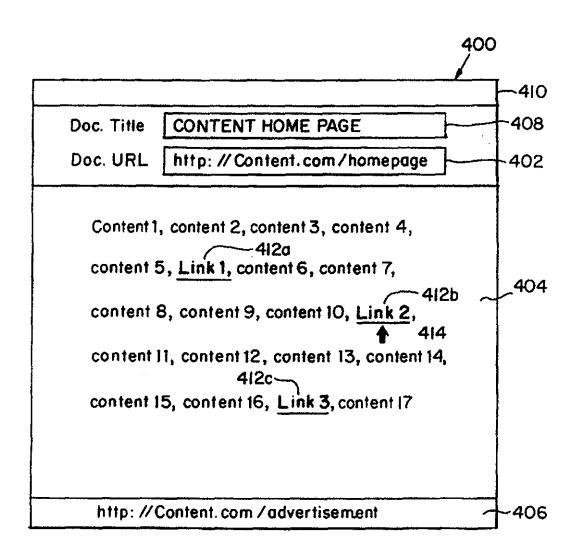


FIG. 4

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| Document View | | | | |
|--|--|--|--|--|
| <u>File Options Navigate Annotate Documents H</u> elp | | | | |
| Title: How to join | | | | |
| URL: http://auth.com/service/nph-createacct.cgi | | | | |
| I. First name | | | | |
| 2. Last name | | | | |
| 3. Choose a screen name (no more than 15 characters) | | | | |
| 4 Channel and the second the seco | | | | |
| 4. Choose a password (no more than 15 characters) Password: | | | | |
| | | | | |
| Re-enter password: | | | | |
| | | | | |
| 5. E-mail address | | | | |
| | | | | |
| 6. Your birthdate (MM/DD/YY | | | | |
| 7. U.S. zip code, or country code | | | | |
| Zip/postal code: | | | | |
| | | | | |
| ISO country code | | | | |
| US | | | | |

FIG. 5

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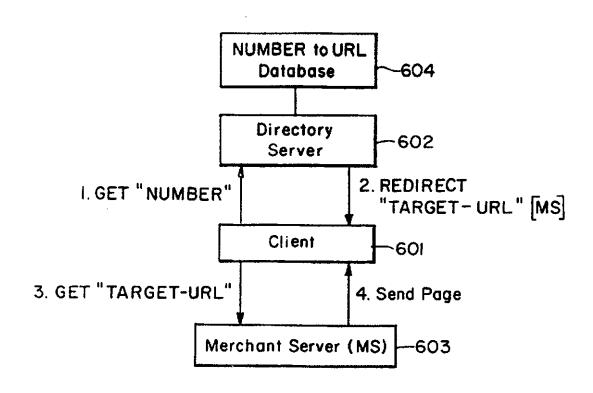


FIG. 6

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INTERNET SERVER ACCESS CONTROL AND MONITORING SYSTEMS

REFERENCE TO APPENDIX

A portion of the disclosure of this patent document contains material which is subject to copyright protection. The copyright owner has no objection to the facsimile reproduction by any one of the patent disclosure, as it appears in the Patent and Trademark Office patent files or records, but otherwise reserves all copyright rights whatsoever.

BACKGROUND OF THE INVENTION

The Internet, which started in the late 1960s, is a vast computer network consisting of many smaller networks that span the entire globe. The Internet has grown exponentially, and millions of users ranging from individuals to corporations now use permanent and dial-up connections to use the Internet on a daily basis worldwide. The computers or networks of computers connected within the Internet, known as "hosts", allow public access to databases featuring information in nearly every field of expertise and are supported by entities ranging from universities and government to many commercial organizations.

The information on the Internet is made available to the public through "servers". A server is a system running on an Internet host for making available files or documents contained within that host. Such files are typically stored on magnetic storage devices, such as tape drives or fixed disks, local to the host. An Internet server may distribute information to any computer that requests the files on a host. The computer making such a request is known as the "client", which may be an Internet-connected workstation, bulletin board system or home personal computer (PC).

TCP/IP (Transmission Control Protocol/Internet Protocol) is one networking protocol that permits full use of the Internet. All computers on a TCP/IP network need unique ID codes. Therefore, each computer or host on the Internet is identified by a unique number code, known as the IP (Internet Protocol) number or address, and corresponding network and computer names. In the past, an Internet user gained access to its resources only by identifying the host computer and a path through directories within the host's storage to locate a requested file. Although various navigating tools have helped users to search resources on the Internet without knowing specific host addresses, these tools still require a substantial technical knowledge of the Internet.

The World-Wide Web (Web) is a method of accessing 50 information on the Internet which allows a user to navigate the Internet resources intuitively, without IP addresses or other technical knowledge. The Web dispenses with command-line utilities which typically require a user to transmit sets of commands to communicate with an Internet 55 server. Instead, the Web is made up of hundreds of thousands of interconnected "pages", or documents, which can be displayed on a computer monitor. The Web pages are provided by hosts running special servers. Software which runs these Web servers is relatively simple and is available 60 on a wide range of computer platforms including PC's. Equally available is a form of client software, known as a Web "browser", which is used to display Web pages as well as traditional non-Web files on the client system. Today, the Internet hosts which provide Web servers are increasing at a 65 rate of more than 300 per month, en route to becoming the preferred method of Internet communication.

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Created in 1991, the Web is based on the concept of "hypertext" and a transfer method known as "HTTP" (Hypertext Transfer Protocol). HTTP is designed to run primarily over TCP/IP and uses the standard Internet setup. 5 where a server issues the data and a client displays or processes it. One format for information transfer is to create documents using Hypertext Markup Language (HTML). HTML pages are made up of standard text as well as formatting codes which indicate how the page should be 10 displayed. The Web client, a browser, reads these codes in order to display the page. The hypertext conventions and related functions of the world wide web are described in the appendices of U.S. patent application Ser. No. 08/328.133, filed on Oct. 24, 1994, by Payne et al. which is incorporated berein by reference.

Each Web page may contain pictures and sounds in addition to text. Hidden behind certain text, pictures or sounds are connections, known as "hypertext links" ("links"), to other pages within the same server or even on other computers within the Internet. For example, links may be visually displayed as words or phrases that may be underlined or displayed in a second color. Each link is directed to a web page by using a special name called a URL (Uniform Resource Locator). URLs enable a Web browser to go directly to any file held on any Web server. A user may also specify a known URL by writing it directly into the command line on a Web page to jump to another Web page.

The URL naming system consists of three parts: the transfer format, the host name of the machine that holds the file, and the path to the file. An example of a URL may be:

http://www.college.univ.edu/Adir/Bdir/Cdir/page.html,

where "http" represents the transfer protocol; a colon and two-forward slashes (://) are used to separate the transfer format from the host name; "www.college.univ.edu" is the host name in which "www" denotes that the file being requested is a Web page; "Adir/Bdir/Cdir" is a set of directory names in a tree structure, or a path, on the host machine; and "page.html" is the file name with an indication that the file is written in HTML.

The Internet maintains an open structure in which exchanges of information are made cost-free without restriction. The free access format inherent to the Internet, however, presents difficulties for those information providers requiring control over their Internet servers. Consider for example, a research organization that may want to make certain technical information available on its Internet server to a large group of colleagues around the globe, but the information must be kept confidential. Without means for identifying each client, the organization would not be able to provide information on the network on a confidential or preferential basis. In another situation, a company may want to provide highly specific service tips over its Internet server only to customers having service contracts or accounts.

Access control by an Internet server is difficult for at least two reasons. First, when a client sends a request for a file on a remote Internet server, that message is routed or relayed by a web of computers connected through the Internet until it reaches its destination host. The client does not necessarily know how its message reaches the server. At the same time, the server makes responses without ever knowing exactly who the client is or what its IP address is. While the server may be programmed to trace its clients, the task of tracing is often difficult, if not impossible. Secondly, to prevent unwanted intrusion into private local area networks (LAN), system administrators implement various data-flow control

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mechanisms, such as the Internet "firewalls", within their networks. An Internet firewall allows a user to reach the Internet anonymously while preventing intruders of the outside world from accessing the user's LAN.

SUMMARY OF THE INVENTION

The present invention relates to methods of processing service requests from a client to a server through a network. In particular the present invention is applicable to processing client requests in an HTIP (Hypertext Transfer Protocol) environment, such as the World-Wide Web (Web). One aspect of the invention involves forwarding a service request from the client to the server and appending a session identification (SID) to the request and to subsequent service requests from the client to the server within a session of requests. In a preferred embodiment, the present method involves returning the SID from the server to the client upon an initial service request made by the client. A valid SID may include an authorization identifier to allow a user to access controlled files.

In a preferred embodiment, a client request is made with a Uniform Resource Locator (URL) from a Web browser. Where a client request is directed to a controlled file without an SID, the Internet server subjects the client to an authorization routine prior to issuing the SID, the SID being protected from forgery. A content server initiates the authorization routine by redirecting the client's request to an authentication server which may be at a different host. Upon receiving a redirected request, the authentication server returns a response to interrogate the client and then issues an SID to a qualified client. For a new client, the authentication server may open a new account and issue an SID thereafter. A valid SID typically comprises a user identifier, an accessible domain, a key identifier, an expiration time such as date, the IP address of the user computer, and an unforgettable digital signature such as a cryptographic hash of all of the other items in the SID encrypted with a secret key. The authentication server then forwards a new request consisting of the original URL appended by the SID to the client in a REDIRECT. The modified request formed by a new URL is automatically forwarded by the client browser to the content server.

When the content server receives a URL request accompanied by an SID, it logs the URL with the SID and the user IP address in a transaction log and proceeds to validate the SID. When the SID is so validated, the content server sends the requested document for display by the client's Web browser.

In the preferred embodiment, a valid SID allows the client 50 to access all controlled files within a protection domain without requiring further authorization. A protection domain is defined by the service provider and is a collection of controlled files of common protection within one or more servers.

When a client accesses a controlled Web page with a valid SID, the user viewing the page may want to traverse a link to view another Web page. There are several possibilities. The user may traverse a link to another page in the same path. This is called a "relative link". A relative link may be made either within the same domain or to a different domain. The browser on the client computer executes a relative link by rewriting the current URL to replace the old controlled page name with a new one. The new URL retains all portions of the old, including the SID, except for the new page name. SI the relative link points to a page in the same protection domain, the SID remains valid, and the request is honored.

However, if the relative link points to a controlled page in a different protection domain, the SID is no longer valid, and the client is automatically redirected to forward the rewritten

URL to the authentication server to update the SID. The updated or new SID provides access to the new domain if the user is qualified.

The user may

The user may also elect to traverse a link to a document in a different path. This is called an "absolute link". In generating a new absolute link, the SID is overwritten by the browser. In the preferred embodiment, the content server, in each serving of a controlled Web page within the domain, filters the page to include the current SID in each absolute URL on the page. Hence, when the user elects to traverse an absolute link, the browser is facilitated with an authenticated URL which is directed with its SID to a page in a different path. In another embodiment, the content server may forego the filtering procedure as above-described and redirect an absolute URL to the authentication server for an update.

An absolute link may also be directed to a controlled file in a different domain. Again, such a request is redirected to the authentication server for processing of a new SID. An absolute link directed to an uncontrolled file is accorded an immediate access.

In another embodiment, a server access control may be maintained by programming the client browser to store an SID or a similar tag for use in each URL call to that particular server. This embodiment, however, requires a special browser which can handle such communications and is generally not suitable for the standard browser format common to the Web.

Another aspect of the invention is to monitor the frequency and duration of access to various pages both controlled and uncontrolled. A transaction log within a content server keeps a history of each client access to a page including the link sequence through which the page was accessed. Additionally, the content server may count the client requests exclusive of repeated requests from a common client. Such records provide important marketing feedback including user demand, access pattern, and relationships between customer demographics and accessed pages and access patterns.

The above and other features of the invention including various novel details of construction and combinations of parts will now be more particularly described with reference to the accompanying drawings and pointed out in the claims. It will be understood that the particular devices and methods embodying the invention are shown by way of illustration only and not as limitations of the invention. The principles and features of this invention may be employed in varied and numerous embodiments without departing from the scope of the invention.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 is a diagram illustrating the Internet operation.

FIG. 2A is a flowchart describing the preferred method of Internet server access control and monitoring.

FIG. 2B is a related flowchart describing the details of the authentication process.

FIG. 3 illustrates an example of a client-server exchange session involving the access control and monitoring method of the present invention.

FIG. 4 is an example of a World Wide Web page.

FIG. 5 is an example of an authorization form page.

FIG. 6 is a diagram describing the details of the translation of telephone numbers to URLs.

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DETAILED DESCRIPTION OF THE INVENTION

Referring now to the drawings, FIG. 1 is a graphical illustration of the Internet. The Internet 10 is a network of millions of interconnected computers 12 including systems owned by Internet providers 16 and information systems (BBS) 20 such as Compuserve or America Online. Individual or corporate users may establish connections to the Internet in several ways. A user on a home PC 14 may purchase an account through the Internet provider 16. Using a modem 22, the PC user can dial up the Internet provide to connect to a high speed modem 24 which, in turn. provides a full service connection to the Internet. A user 18 may also make a somewhat limited connection to the Internet gateway connection to its customers.

FIG. 2A is a flowchart detailing the preferred process of the present invention and FIG. 4 illustrates a sample Web page displayed at a client by a browser. The page includes 20 text 404 which includes underlined link text 412. The title bar 408 and URL bar 402 display the title and URL of the current web page, respectively. As shown in FIG. 4, the title of the page is "Content Home Page" and the corresponding URL is "http://content.com/homepage". When a cursor 414 is positioned over link text 412b, the page which would be retrieved by clicking a mouse is typically identified in a status bar 406 which shows the URL for that link. In this example the status bar 406 shows that the URL for the pointed link 412b is directed to a page called 30 "advertisement", in a commercial content server called "content". By clicking on the link text, the user causes the browser to generate a URL GET request at 100 in FIG. 2A. The browser forwards the request to a content server 120, which processes the request by first determining whether the 35 requested page is a controlled document 102. If the request is directed to an uncontrolled page, as in "advertisement" page in this example, the content server records the URL and the IP address, to the extent it is available, in the transaction log 114. The content server then sends the requested page to 40 the browser 116 for display on the user computer 117.

If the request is directed to a controlled page, the content server determines whether the URL contains an SID 102. For example, a URL may be directed to a controlled page name "report", such as "http://content.com/report", that 45 requires an SID. If no SID is present, as in this example, the content server sends a "REDIRECT" response 122 to the browser 100 to redirect the user's initial request to an authentication server 200 to obtain a valid SID. The details of the authentication process are described in FIG. 2B and 50 will be discussed later, but the result of the process is an SID provided from the authentication server to the client. In the above example, a modified URL appended with an SID may be: "http://content.com/[SID]/report". The preferred SID is a sixteen character ASCII string that encodes 96 bits of SID 55 data, 6 bits per character. It contains a 32-bit digital signature, a 16-bit expiration date with a granularity of one hour, a 2-bit key identifier used for key management, an 8-bit domain comprising a set of information files to which the current SID authorizes access, and a 22-bit user identi- 60 fier. The remaining bits are reserved for expansion. The digital signature is a cryptographic hash of the remaining items in the SID and the authorized IP address which are encrypted with a secret key which is shared by the authentication and content servers.

If the initial GET URL contains a SID, the content server determines whether the request is directed to a page within Ó

the current domain 106. If the request having a SID is directed to a controlled page of a different domain, the SID is no longer valid and, again, the user is redirected to the authentication server 122.

If the request is for a controlled page within the current domain, the content server proceeds to log the request URL, tagged with SID, and the user IP address in the transaction log 108. The content server then validates the SID 110. Such validation includes the following list of checks: (1) the SID's digital signature is compared against the digital signature computed from the remaining items in the SID and the user IP address using the secret key shared by the authentication and content servers; (2) the domain field of the SID is checked to verify that it is within the domain authorized; and (3) the EXP field of the SID is checked to verify that it is later than the current time.

If the validation passes, the content server searches the page to be forwarded for any absolute URL links contained therein 112, that is, any links directed to controlled documents in different content servers. The content server augments each absolute URL with the current SID to facilitate authenticated accesses across multiple content servers. The requested page as processed is then transmitted to the client browser for display 117. The user viewing the requested Web page may elect to traverse any link on that page to trigger the entire sequence again 100.

FIG. 2B describes the details of the authentication process. The content server may redirect the client to an authentication server. The REDIRECT URL might be: "http://auth.com/authenticate?domain=[domain]&URL=http://content.com/report". That URL requests authentication and specifies the domain and the initial URL. In response to the REDIRECT, the client browser automatically sends a GET request with the provided URL.

Whenever the content server redirects the client to the authentication server 200, the authentication server initiates the authorization process by validating that it is for an approved content server and determining the level of authentication required for the access requested 210. Depending on this level, the server may challenge the user 212 for credentials. If the request is for a low level document, the authentication may issue an appropriate SID immediately 228 and forego the credential check procedures. If the document requires credentials, the authentication server sends a "CHALLENGE" response which causes the client browser to prompt the user for credentials 214. A preferred credential query typically consists of a request for user name and password. If the user is unable to provide a password, the access is denied. The browser forms an authorization header 300 from the information provided, and resends a GET request to the authentication server using the last URL along with an authorization header. For example, a URL of such a GET request may be: "http://auth.com/ authenticate?domain=[domain]&URL=http://content.com/ report and the authorization header may be: "AUTHORIZE: authorization ...

Upon receiving the GET request, the authentication server queries an account database 216 to determine whether the user is authorized 218 to access the requested document. A preferred account database may contain a user profile which includes information for identifying purposes, such as client IP address and password, as well as user demographic information, such as user age, home address, hobby, or occupation, for later use by the content server. If the user is authorized, an SID is generated 228 as previously described. If the user is not cleared for authorization, the authentication

server checks to see if the user qualifies for a new account 220. If the user is not qualified to open a new account, a page denying access 222 is transmitted to the client browser 100. If the user is qualified, the new user is sent a form page such as illustrated in FIG. 5 to initiate a real-time on-line registration 224. The form may, for example, require personal information and credit references from the user. The browser is able to transmit the data entered by the user in the blanks 502 as a "POST" message to the authentication server. A POST message causes form contents to be sent to the server in a data body other than as part of the URL. If the registration form filled out by the new user is valid 226, an appropriate SID is generated 228. If the registration is not valid, access is again denied 222.

An SID for an authorized user is appended ("tagged") 230 to the original URL directed to a controlled page on the content server. The authentication server then transmits a REDIRECT response 232 based on the tagged URL to the client browser 100. The modified URL, such as "http:// content.com/[SID]/report" is automatically forwarded to the 20 content server 120.

FIG. 3, illustrates a typical client-server exchange involving the access control and monitoring method of the present invention. In Step 1, the client 50 running a browser transmits a GET request through a network for an uncontrolled 25 page (UCP). For example, the user may request an advertisement page by transmitting a URL "http://content.com/ advertisement", where "content.com" is the server name and "advertisement" is the uncontrolled page name. In Step 2. the content server 52 processes the GET request and transmits the requested page, "advertisement". The content server also logs the GET request in the transaction database 56 by recording the URL, the client IP address, and the current time.

In Step 3, the user on the client machine may elect to 35 traverse a link in the advertisement page directed to a controlled page (CP). For example, the advertisement page may contain a link to a controlled page called "report". Selecting this link causes the client browser 50 to forward a GET request through a URL which is associated with the 40 report file "http://content.com/report". The content server 52 determines that the request is to a controlled page and that the URL does not contain an SID. In Step 4, the content server transmits a REDIRECT response to the client, and, in Step 5, the browser automatically sends the REDIRECT 45 URL to the authentication server 54. The REDIRECT URL sent to the authentication server may contain the following string:

"http://auth.com/authenticate?domain=[domain]&URL= http://content.com/report".

The authentication server processes the REDIRECT and determines whether user credentials (CRED) are needed for authorization. In Step 6, the authentication server transmits a "CHALLENGE" response to the client. As previously described, typical credentials consist of user name and 55 password. An authorization header based on the credential information is then forwarded by the client browser to the authentication server. For example, a GET URL having such an authorization header is:

&URL=http://content.com/report and the authorization header may be: "AUTHORIZE:[authorization]". The authentication server processes the GET request by checking the Account Database 58. If a valid account exists for the user, an SID is issued which authorizes 65 access to the controlled page "report" and all the other pages within the domain.

As previously described, the preferred SID comprises a compact ASCII string that encodes a user identifier, the current domain, a key identifier, an expiration time, the client IP address, and an unforgeable digital signature. In Step 8, the authentication server redirects the client to the tagged URL, "http://content.com/[SID]/report", to the client. In Step 9, the tagged URL is automatically forwarded by the browser as a GET request to the content server. The content server logs the GET request in the Transaction database 56 by recording the tagged URL, the client IP address, and the current time. In Step 10, the content server. upon validating the SID, transmits the requested controlled page "report" for display on the client browser.

According to one aspect of the present invention, the content server periodically evaluates the record contained in the transaction log 56 to determine the frequency and duration of accesses to the associated content server. The server counts requests to particular pages exclusive of repeated requests from a common client in order to determine the merits of the information on different pages for ratings purposes. By excluding repeated calls, the system avoids distortions by users attempting to "stuff the ballot box." In one embodiment, the time intervals between repeated requests by a common client are measured to exclude those requests falling within a defined period of

Additionally, the server may, at any given time, track access history within a client-server session. Such a history profile informs the service provider about link transversal frequencies and link paths followed by users. This profile is produced by filtering transaction logs from one or more servers to select only transactions involving a particular user ID (UID). Two subsequent entries, A and B, corresponding to requests from a given user in these logs represent a link traversal from document A to document B made by the user in question. This information may be used to identify the most popular links to a specific page and to suggest where to insert new links to provide more direct access. In another embodiment, the access history is evaluated to determine traversed links leading to a purchase of a product made within commercial pages. This information may be used, for example, to charge for advertising based on the number of link traversals from an advertising page to a product page or based on the count of purchases resulting from a path including the advertisement. In this embodiment, the server can gauge the effectiveness of advertising by measuring the number of sales that resulted from a particular page, link, or path of links. The system can be configured to charge the merchant for an advertising page based on the number of sales that resulted from that page.

According to another aspect of the present invention, a secondary server, such as the authentication server 200 in FIG. 2B, may access a prearranged user profile from the account database 216 and include information based on such a profile in the user identifier field of the SID. In a preferred embodiment, the content server may use such an SID to customize user requested pages to include personalized content based on the user identifier field of the SID.

In another aspect of the invention, the user may gain "http://autho.com/authenticate?domain=[domain] 60 access to domain of servers containing journals or publications through a subscription. In such a situation, the user may purchase the subscription in advance to gain access to on-line documents through the Internet. The user gains access to a subscribed document over the Internet through the authorization procedure as described above where an authorization indicator is preferably embedded in a session identifier. In another embodiment, rather than relying on a

prepaid subscription, a user may be charged and billed each time he or she accesses a particular document through the. Internet. In that case, authorization may not be required so long as the user is fully identified in order to be charged for the service. The user identification is most appropriately 5 embedded in the session identifier described above.

In another aspect of the invention, facilities are provided to allow users to utilize conventional telephone numbers or other identifiers to access merchant services. These merchant services can optionally be protected using SIDs. In a 10 preferred embodiment, as shown in FIG. 6, a Web browser client 601 provides a "dial" command to accept a telephone number from a user, as by clicking on a "dial" icon and inputting the telephone number through the keyboard. The browser then constructs a URL of the form "http:// 15 directory.net/NUMBER", where NUMBER is the telephone number or other identifier specified by the user. The browser then performs a GET of the document specified by this URL, and contacts directory server 602, sending the NUMBER requested in Message 1.

In another embodiment, implemented with a conventional browser, client 601 uses a form page provided by directory server 602 that prompts for a telephone number or other identifier in place of a "dial" command, and Message 1 is a POST message to a URL specified by this form page.

Once NUMBER is received by directory server 602, the directory server uses database 604 to translate the NUMBER to a target URL that describes the merchant server and document that implements the service corresponding to NUMBER. This translation can ignore the punctuation of 30 the number, therefore embedded parenthesis or dashes are not significant.

In another embodiment an identifier other than a number may be provided. For example, a user may enter a company name or product name without exact spelling. In such a case 35 a "soundex" or other phonetic mapping can be used to permit words that sound alike to map to the same target URL. Multiple identifiers can also be used, such as a telephone number in conjunction with a product name or extension.

In Message 2. Directory server 602 sends a REDIRECT to client 601, specifying the target URL for NUMBER as

computed from database 604. The client browser 601 then automatically sends Message 3 to GET the contents of this URL. Merchant server 603 returns this information in Message 4. The server 602 might have returned a Web page to the client to provide an appropriate link to the required document. However, because server 602 makes a translation to a final URL and sends a REDIRECT rather than a page to

client 601, the document of message 4 is obtained without

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any user action beyond the initial dial input.

The Target URL contained in Message 3 can be an ordinary URL to an uncontrolled page, or it can be a URL that describes a controlled page. If the Target URL describes a controlled page then authentication is performed as previously described. The Target URL can also describe a URL that includes an SID that provides a preauthorized means of accessing a controlled page.

Among benefits of the "dial" command and its implementation is an improved way of accessing the Internet that is compatible with conventional telephone numbers and other identifiers. Merchants do not need to alter their print or television advertising to provide an Internet specific form of contact information, and users do not need to learn about URLs.

In the approach a single merchant server can provide 25 multiple services that correspond to different external "telephone numbers" or other identifiers. For example, if users dial the "flight arrival" number they could be directed to the URL for the arrival page, while, if they dial the "reservations" number, they would be directed to the URL for the reservations page. A "priority gold" number could be directed to a controlled page URL that would first authenticate the user as belonging to the gold users group, and then would provide access to the "priority gold" page. An unpublished "ambassador" number could be directed to a tagged URL that permits access to the "priority gold" page without user authentication.

Equivalents

Those skilled in the art will know, or be able to ascertain using no more than routine experimentation, many equivalents to the specific embodiments or the invention described herein. These and all other equivalents are intended to be

encompassed by the following claims.

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```
-30-
  * tclSidSup.c --
           tcl SID pack/unpack/id routines
     Steve Morris
     morris@openmarket.com
/* #define debug 1 */
#include <stdio.h>
#include <sys/types.h>
#include <stdlib.h>
#include <unistd.h>
#ifdef __alpha
#include <sys/time.h>
#endif
#include <tcl.h>
#include *global.h*
#include *md5.h*
#ifdef debug
#define dbg_out(s) printf("%s\n", s);
#define dbg_outi(s,a) printf("%s %x\n", s,a);
#define dbg_outs(s,a) printf("%s %s\n", s,a);
#else
#define dbg_out(s)
#define dbg_outi(s,a)
#define dbg_outs(s,a)
#endif
/* external routines called
        char * radix64decode_noslash(char *in, int len, int *output_len);
char * radix64encode_noslash(char *in, int len);
int compute_ihash(char *str);
                   sid_rev_zero (0)
#define
#define sig_lw
                          0
#define sig_mask
                          Oxffffffff
#define kid_lw
#define kid_pos
#define kid_mask
                          0x3
#define rev_lw
#define rev_pos
#define rev_mask
#define uid_lw
```

13

```
~31-
#.efine uid_pos
                                                0x3fffff
 #define wid mask
#define rsrv1_lw 1
#define rsrv1_pos 26
#define rsrv1_mask 0x3f
#define dom_lw
#define dom_pos
#define dom_mask
                                                0xff
#define dom_mask 0xf1
#define rsrv2_lw 2
#define rsrv2_pos 8
#define rsrv2_mask 0x3f
#define uctx_lw 2
#define uctx_pos 14
#define uctx_mask 0x3
#define uctx_mask 0x3
#define exp_lw 2
#define exp_pos 16
#define exp_mask 0xffi
                                               0xffff
#define exp_mask
#ifdef __sun_
#define fix_endian(i,o)
                                                                       ecp = (char *) 1; \
eda = (*ecp++) <<0; \
eda |= (*ecp++) <<8; \
eda |= (*ecp++) <<16; \
eda |= (*ecp) <<24; \
(int *) *o = eda; \)
 /* #define fix_endian(i,0)
                                                                         { (int *) *o = (int *) *i;} */
 #define fix_endian(i,o)
#endif
```

5.708.780

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16

-32-

```
/* TclUnpackSid(NoValidate)
        Converts a binary sid to the internal representation, and does minimal validation on the SID, the NoValidate version does even less.
 * Inputs:
                               a sid, looks like "/@@u9GIg4QEAAAABPn"
        the sid
 * Returns:
        ($d $d $d $d)
                            on succes (dom uid keyid exp) on failure
 * or
 */
{
     char temp[512];
     int ret_status = 1;
     int i;
int * bsid = NULL, act_hash;
char ip_address[32]="", secret(200)="", *cp;
     unsigned char *ecp;
unsigned int eda;
     char hash_buffer[512];
     struct sid_info
       unsigned int sig;
unsigned int kid;
unsigned int uid;
unsigned int rev;
unsigned int dom;
unsigned int exp;
        unsigned int uctx;
        } esid;
     if (argc != 2)
       Tcl_AppendResult(interp, *wrong # args: should be \"", argv[0],
    " string\"", (char *) NULL);
return TCL_ERROR;
#ifdef test_ip
     strcpy(ip_address, test_ip);
#ifdef test secret
     strcpy(secret, test_secret);
#endif
       cp = Tcl_GetVar(interp, 'sidip', TCL_GLOBAL_ONLY|TCL_LEAVE_ERR_MSG);
if (cp != NULL) strncpy(ip_address, cp, 31);
       cp = argv[1]+3;
       i = strlen(cp);
if (i != 16) goto sid_bad;
```

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```
/* first convert the SID back to binary*/
       if if to convert the state of the convert the baid = (int *) radix64decode noslash(cp, i, &i);
if (bsid == NULL) goto sid_bad;
if (i != 12) goto sid_bad;
        fix_endian(&bsid(0), &bsid(0));
fix_endian(&bsid(1), &bsid(1));
fix_endian(&bsid(2), &bsid(2));
       dbg_outi("sid[0]=",bsid[0]);
dbg_outi("sid[1]=",bsid[1]);
dbg_outi("sid[2]=",bsid[2]);
      dbg_outi(*sig =*,get_sid(sig_lw,sig_pos,sig_mask));
dbg_outi(*kid =*,get_sid(kid_lw,kid_pos,kid_mask));
dbg_outi(*uid =*,get_sid(uid_lw,uid_pos,uid_mask));
dbg_outi(*rev =*,get_sid(rev_lw,rev_pos,rev_mask));
dbg_outi(*dom =*,get_sid(dom_lw,dom_pos,dom_mask));
dbg_outi(*exp =*,get_sid(exp_lw,exp_pos,exp_mask));
dbg_outi(*uctx =*,get_sid(uctx_lw,uctx_pos,uctx_mask));
dbg_outi(*rs1 =*,get_sid(rsrv1_lw,rsrv1_pos,rsrv1_mask));
dbg_outi(*rs2 =*,get_sid(rsrv2_lw,rsrv2_pos,rsrv2_mask));
      /* check the SID version field */
if (get_sid(rev_lw,rev_pos,rev_mask) != sid_rev_zero) goto sid_bad;
if (get_sid(rsrvl_lw,rsrvl_pos,rsrvl_mask) != 0) goto sid_bad;
if (get_sid(rsrv2_lw,rsrv2_pos,rsrv2_mask) != 0) goto sid_bad;
        /* unpack the sid */
      /* unpack the sid */
esid.sig = get sid(sig lw,sig pos,sig mask);
esid.kid = get_sid(kid lw,kid pos,kid mask);
esid.uid = get_sid(uid lw,uid pos,uid mask);
esid.rev = get_sid(rev lw,rev pos,rev mask);
esid.dom = get_sid(dom lw,dom pos,dom mask);
esid.exp = get_sid(exp lw,exp pos,exp mask);
esid.uctx = get_sid(uctx lw,uctx pos,uctx mask);
      dbg_outi("sig =",esid.sig);
dbg_outi(*kid =",esid.kid);
dbg_outi(*uid =",esid.uid);
dbg_outi("rev =",esid.rev);
dbg_outi("dom =",esid.dom);
       dbg_outi(*exp =",esid.exp);
dbg_outi(*uctx = ",esid.uctx);
       sprintf(temp, *secret (%d) *, esid.kid);
cp = Tcl_GetVar(interp, temp, TCL_GLOBAL_ONLY);
if (cp != NULL) strncpy(secret, cp, 199);
    /* hash the sid and check the signature */
  /* hash the sid and check the signature */
sprintf(hash_buffer, "%s%s%08x%08x", secret,ip_address,bsid{2},bsid{1});
dbg_outs("hashing ->",hash_buffer);
act_hash = compute_ihash(hash_buffer);
fix_endian(&act_hash, &act_hash);
dbg_outi("expected hash=",bsid{0});
dbg_outi("actual hash=",act_hash);
if {act_hash != esid.sig} goto sid_bad;
rtn exit:
```

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}

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-34-

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-35-

```
int TclUnpackSidNoValidate(ClientData dummy, Tcl_Interp *interp,
                              int argc, char **argv)
     char temp[512];
     int ret_status = 1;
     int i;
     int * bsid = NULL, act_hash;
char ip_address[32]="", secret[200]="", *cp;
     unsigned char *ecp;
     unsigned int eda;
char hash_buffer(512);
     struct sid_info
       unsigned int sig;
       unsigned int kid;
       unsigned int uid;
       unsigned int rev;
       unsigned int dom;
       unsigned int exp;
       unsigned int uctx;
       } esid;
     if (argc != 2)
       return TCL_ERROR;
#ifdef test_ip
     strcpy(ip_address, test_ip);
#endif
#ifdef test_secret
    strcpy(secret, test_secret);
       cp = Tcl_GetVar(interp, *sidip*, TCL_GLOBAL_ONLY|TCL_LEAVE_ERR_MSG);
if (cp != NULL) strncpy(ip_address, cp, 31);
       cp = argv[1]+3;
         = strlen(cp);
       if (i != 16) goto sid_bad;
        /* first convert the SID back to binary*/
       bsid = (int *) radix64decode_noslash(cp, i, &i);
if (bsid == NULL) goto sid bad;
if (i != 12) goto sid bad;
       fix_endian(&bsid{0}, &bsid[0]);
fix_endian(&bsid{1}, &bsid{1});
fix_endian(&bsid{2}, &bsid{2});
       dbg_outi("sid[0]=",bsid[0]);
dbg_outi("sid[1]=",bsid[1]);
dbg_outi("sid[2]=",bsid[2]);
       dbg_outi("sig =",get_sid(sig_lw,sig_pos,sig_mask));
```

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```
dbg_outi("kid =",get_sid(kid_lw,kid_pos,kid_mask));
dbg_outi("uid =",get_sid(uid_lw,uid_pos,uid_mask));
dbg_outi("rev =",get_sid(rev_lw,rev_pos,rev_mask));
dbg_outi("dom =",get_sid(dom_lw,dom_pos,dom_mask));
dbg_outi("exp =",get_sid(exp_lw,exp_pos,exp_mask));
dbg_outi("uctx =",get_sid(uctx_lw,uctx_pos,uctx_mask));
dbg_outi("rs1 =",get_sid(rsrv1_lw,rsrv1_pos,rsrv1_mask));
dbg_outi("rs2 =",get_sid(rsrv2_lw,rsrv2_pos,rsrv2_mask));
                 /* check the SID version field */
               if (get_sid(rev_lw,rev_pos,rev_mask) != sid_rev_zero) goto sid_bad; if (get_sid(rsrvl_lw,rsrvl_pos,rsrvl_mask) != 0) goto sid_bad; if (get_sid(rsrv2_lw,rsrv2_pos,rsrv2_mask) != 0) goto sid_bad;
               /* unpack the sid */
esid.sig = get_sid(sig_lw,sig_pos,sig_mask);
esid.kid = get_sid(kid_lw,kid_pos,kid_mask);
esid.uid = get_sid(uid_lw,uid_pos,uid_mask);
esid.rev = get_sid(rev_lw,rev_pos,rev_mask);
esid.dom = get_sid(dom_lw,dom_pos,dom_mask);
esid.exp = get_sid(exp_lw,exp_pos,exp_mask);
esid.uctx = get_sid(uctx_lw.uctx_pos,uctx_mask);
                esid.uctx = get_sid(uctx_lw,uctx_pos,uctx_mask);
                dbg_outi("sig =".esid.sig);
              dbg_outi("sig =".esid.sig);
dbg_outi(*kid =".esid.kid);
dbg_outi("uid =".esid.uid);
dbg_outi("rev =".esid.rev);
dbg_outi("dom =".esid.dom);
dbg_outi("exp =".esid.exp);
dbg_outi("uctx =".esid.uctx);
#if (1 == 0) /* disable validation for the novalidate case */
               sprintf(temp, "secret(%d)", esid.kid);
               cp = Tcl_GetVar(interp, temp, TCL_GLOBAL_ONLY);
               if (cp != NULL) strncpy(secret, cp, 199);
           /* hash the sid and check the signature */
sprintf(hash_buffer, "%s%s%08x*08x*, secret,ip_address,bsid[2],bsid[1]);
dbg_outs("hashing ->",hash_buffer);
act_hash = compute_ihash(hash_buffer);
           act_hash = compute_Inash(hash_buffet,;
fix_endian(&act_hash, &act_hash);
dbg_outi(*expected hash=*,bsid[0]);
dbg_outi(*actual hash=*,act_hash);
if (act_hash != esid.sig) goto sid_bad;
#endif
         rtn_exit:
   if (bsid != NULL) free(bsid);
               if (ret_status) sprintf(interp->result, "{ td td td td td}",
              esid.dom,esid.uid,esid.kid,esid.exp, esid.uctx);
else sprintf(interp->result,"");
              return TCL_OK;
        sid_bad:
              ret_status = 0;
              goto rtn_exit;
      exp_bad:
```

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ret_status = 0;
goto rtn_exit;

}

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```
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```

```
/* TclPackSid
                         Creates the ascii representation of a binary sid
            Inputs:
                          "{ td td td td [td] } Dom uid keyid exp [uctx]
          Returns:
                         ascii bin sid like "/@@u9GIg4QEAAAAABPn"
int TclPackSid(ClientData dummy, Tcl_Interp *interp,
                                                                                                 int argc, char **argv)
              char temp{512};
int bsid[3] = {0,0,0}, i, act_hash;
unsigned char *ecp;
unsigned int eda;
               char ip_address(32), secret(200)=*", *cp;
               int dom, wid, keyid, exp, uctx;
char hash_buffer(512);
               struct sid_info
                        unsigned int sig;
                        unsigned int kid;
                        unsigned int uid;
                        unsigned int rev;
                       unsigned int dom;
                       unsigned int exp;
                       unsigned int uctx;
                        } esid;
               if (argc != 2)
                       Land to the control of the cont
                       return TCL_ERROR;
                 i = sscanf(argv[1], **d *d *d *d *d*d. &dom, &uid, &keyid, &exp, &uctx);
dbg_outs(*scanning...*,argv[1]);
dbg_outi(*scan returned ",i);
if (i == 4) {uctx = 0; i = 5;}
if (i != 5) return TCL_ERROR;
                      esid.kid = keyid;
                      esid.uid = uid;
                      esid.dom = dom;
                      esid.exp = exp;
                   esid.uctx = uctx;
dbg_outi("kid=",esid.kid);
dbg_outi("uid=",esid.uid);
dbg_outi("dom=",esid.dom);
dbg_outi("exp=",esid.exp);
dbg_outi("uctx=",esid.uctx);
```

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```
#ifdef test_ip
       strcpy(ip_address, test_ip);
#endif
#ifdef test_secret
       stropy(secret, test_secret);
#endif
           cp = Tcl_GetVar(interp, "sidip", TCL_GLOBAL_ONLY);
           if (cp != NULL) strncpy(ip_address, cp, 31);
           sprintf(temp,*secret(%d)",esid.kid);
cp = Tcl_GetVar(interp, temp, TCL_GLOBAL_ONLY);
if (cp != NULL) strncpy(secret, cp, 199);
           put_sid(kid_lw,
                                             kid_pos,
                                                                   kid_mask,
                                                                                           esid.kid);
           put_sid(kid_lw, uid_pos, uid_mask, esid.uid);
put_sid(rsrvl_lw, rsrvl_pos, rsrvl_mask, 0);
put_sid(rsrvl_lw, rsrvl_pos, rev_mask, sid_rev_zero);
           put_sid(rev_lw,
           put_sid(dom_1w,
                                                                                           esid.dom);
                                              dom pos.
                                                                   dom_mask,
          put_sid(rsrv2_lw, rsrv2_pos, rsrv2_mask, 0);
put_sid(exp_lw, exp_pos, exp_mask, esi
put_sid(uctx_lw, uctx_pos, uctx_mask, esi
                                                                                           esid.exp);
                                                                                           esid.uctx);
           dbg_outi("sid[0] = ", bsid(0]);
dbg_outi("sid[1] = ", bsid[1]);
dbg_outi("sid[2] = ", bsid[2]);
           sprintf(hash_buffer, "%s%s%08x%08x", secret,ip_address,bsid{2},bsid{1});
dbg_outs("ascii sid built ",hash_buffer);
act_hash = compute_ihash(hash_buffer);
          fix endian(&act hash, &act hash); */
put_sid(sig_lw, sig_pos, sig_mask, act_hash);
           dbg_outi("sid{0}=",bsid{0});
dbg_outi("sid{1}=",bsid{1});
dbg_outi("sid{2}=",bsid{2});
1.
           fix_endian(&bsid[0], &bsid[0]); */
fix_endian(&bsid[1], &bsid[1]);
fix_endian(&bsid[2], &bsid[2]);
           dbg_outi("sid{0}=",bsid{0});
dbg_outi("sid{1}=",bsid{1});
dbg_outi("sid{2}=",bsid{2});
           cp = radix64encode_noslash((char *) bsid, 12);
if (cp == NULL) return TCL_ERROR;
sprintf(temp, "/@@%s", cp);
TCl_SetResult(interp, temp, TCL_VOLATILE);
           if (cp != NULL) free(cp);
return TCL_OK;
}
```

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```
/* TclIdSid
           Scans an ascii line and finds an ascii SID. (no validation though)
     Inputs:
           lineoftext
   * Returns:
           ascii bin_sid, if a sid is found it is returned.
int TclIdSid(ClientData dummy, Tcl_Interp *interp,
                                           int argc, char **argv)
       char *sidp, *cp;
       interp->result(0) = 0;
       if (argc != 2)
           interp->result = "wrong # args";
return TCL_ERROR;
      sidp = (char *) strstr(argv[1], */@@*);
if (sidp == NULL) return TCL_OK;
cp = (char *) strstr(sidp+1,"/*);
if ((cp == NULL) && (strlen(sidp) != 19)) return TCL_OK;
if ((cp - sidp) != 19) return TCL_OK;
strncpy(interp->result, sidp,19);
interp->result[19] = 0;
return TCL_OK;
}
 * Register commands with interpreter.
int SidSupInit(Tcl_Interp *interp)
  Tcl_CreateCommand(interp, "packsid", TclPacksid, NULL, NULL);
Tcl_CreateCommand(interp, "unpacksid", TclUnpacksid, NULL, NULL);
Tcl_CreateCommand(interp, "unpacksidnovalidate", TclUnpacksidNovalidate, NULL,
Tcl_CreateCommand(interp, "issid", TclIdSid, NULL, NULL);
                                                                                                 NULL, NULL);
   return TCL_OK;
```

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```
* ticket.c --

* Commands for TICKET.

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* disclosure, or reproduction is prohibited except as permitted by
* express written license agreement with Open Market, Inc.

* Steve Morris

* morris@OpenMarket.com

* Created: Wed Mar 1 1995

* SSource: /omi/proj/master/omhttpd/Attic/ticket.c,v $

* */

#if !defined(lint)
static const char rcsid[]="$Header: /omi/proj/master/omhttpd/Attic/ticket.c,v 2.
#endif /*not lint*/

#include <stys/utsname.h>
#include "md5.h"
#include "md5.h"
#include "ticket.h"

static TICKET_Server TicketServerData;
```

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لشتميد

* This file implements all the ticket/sid related functions for the server.

* The region commands RequireSID and xxxxx can be used to limit
* access to groups of files based on the authentication of the requestor.
* The two commands are very similar, and only differ in the method used to
* present the authentication data (via the URL) and in handling of the
* failing access case. For failing TICKET's, a "not authorized" message is
* generated. For failing (or absent) SID's, a REDIRECT (either local or via
* CGI script) is performed to forward the request to an authentication server.

* RequireSID domain1 [domain2 ... domainn]

This command denies access unless the specified properties are true of the request. Only one RequireSID or xxxxx command can be used for a given region, though it may specify multiple domains.

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*

* TICKET_RequireSidCmd -
* Checks that the requested URL is authorized via SID to access this region. If the access is not authorized and we do not have a "remote authentication server" registered, then an "unauthorized message" is returned. If a "remote authentication server" has been declared, we REDIRECT to that server, passing the requested URL and required domain's as arguments.

* Results:

* Normal Tcl result, or a REDIRECT request.

* Side effects:

* Either an "unauthorized access" message or a REDIRECT in case of error.

*/

* Static int TICKET_RequireSidCmd(ClientData clientData, Tcl_Interp *interp, int argc, char **argv)

{
if (TicketGlobalData(EnableSidEater)) return TCL_OK; return(ProcessRequires(clientData, interp, argc, argv, ticketSid));

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```
ProcessRequires --
         Checks that the requested URL is authorized to access this
         region. The error cases are treated differently for SID v.s. TICKET.
         For Ticket's, an unauthorized access generates a returned error message.
         For SID's, we first look to see if we are operating in "local authentica
         mode", if we are, we generate a new SID, into the URL and re-process the If not in "local" mode, we look for the presence of a remoteauthenticati
         server, if we have one declared (in the conf file) we REDIRECT to it pas
         the FULL url and a list of domains that would have been legal. If the
         authentication server was not found we return an error message.
         Normal Tcl result, a local reprocess command, or a REDIRECT request.
   Side effects:
         Either an "unauthorized access" message or a REDIRECT in case of error.
static int ProcessRequires(ClientData clientData, Tcl_Interp *interp,
                                int argc, char **argv, int flavor)
    HTTP_Request *reqPtr = (HTTP_Request *) clientData;
HTTP_Server *serverPtr;
    TICKET_Request *ticketPtr;
    DString targetUrl;
DString escapeUrl;
    int i, required dom:
    int firstLegalDom = -1;
    char *NewSid, *cp;
    DStringInit(&targetUrl);
    DStringInit (&escapeUrl);
    /* fetch the server private and ticket specific extension data */
    serverPtr = reqPtr->serverPtr;
ticketPtr = (TICKET_Request *) HT_GetReqExtData(reqPtr, TicketServerData.tic
    ASSERT (ticketPtr != NULL);
   /* compare the requesting SID/Ticket<DOM> to authorized list of domains */ /* a match OR any valid domain and a required domain of TicketFreeArea is su for (i = 1; i < argc; i++)
     required_dom = GetDomain(argv[i],-1);
if (required_dom != -1)
       ((DStringLength(&ticketPtr->ticketIP) == 0) ||
```

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```
(strcmp(DStringValue(&ticketPtr->ticketIP), DStringValue(&reqPtr->r
         DStringFree (&targetUrl);
         DStringFree (&escapeUrl);
         return TCL_OK;
 )
/* count the number of domain crossing that caused re-auth */
if ((flavor == ticketSid) && (ticketPtr->sidDom) != -1) IncTicketCounter(Cou
/st authorization failed, if this was a sid url, and local auth is enabled st/
/* or this was an access to the free area */
/* insert a new sid in the url, and REDIRECT back to the client */
if (TicketGlobalData(EnableLocalAuth) | ((firstLegalDom == TicketGlobalData(FreeArea))
    && (flavor == ticketSid) && (firstLegalDom != -1)))
  http Error(reqptr, NOT_FOUND, "access denied due to poorly formed url");
DStringFree(&targetUrl);
    DStringFree (&escapeUrl);
    if (!ticketPtr->valid)
      DStringFree (&ticketPtr->sid);
    return TCL_RETURN ;
  NewSid = CreateSid(reqPtr,
             firstLegalDom, ticketPtr->uid,
             TicketGlobalData(CurrentSecret), TicketGlobalData(LocalAuthExp),
             ticketPtr->uctx);
  DStringFree (&ticketPtr->sid);
  DStringAppend(&ticketPtr->sid, NewSid, -1);
  ComposeURL(reqPtr, DStringValue(&reqPtr->url), &targetUrl);
IncTicketCounter(CountLocalRedirects);
  HTTP_Error(reqPtr, REDIRECT, DStringValue(&targetUrl));
  DStringFree(&targetUrl);
  DStringFree (&escapeUrl);
  if (!ticketPtr->valid)
    DStringFree (&ticketPtr->sid);
  return TCL_RETURN;
/* authorization failed, build the REDIRECT URL arg's. */
/* If present, REDIRECT to authentication server */
if ((DStringLength(&TicketGlobalData(AuthServer)) != 0) && (flavor == ticket
  HTTP Error (reqPtr, NOT FOUND, "access denied due to poorly formed url");
    DStringFree (&targetUrl);
    DStringFree (&escapeUrl);
    if (!ticketPtr->valid)
      DStringFree (&ticketPtr->sid);
```

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```
47
           return TCL_RETURN ;
      DStringAppend(&targetUrl, DStringValue(&TicketGlobalData(AuthServer)), -1)
     DStringAppend(&targetUrl, "?url=", -1);
ComposeURL(reqPtr, DStringValue(&reqPtr->url), &escapeUrl);
      EscapeUrl (&escapeUrl);
     DStringAppend(&targetUrl, DStringValue(&escapeUrl), -1);
     DStringAppend(&targetUrl, "&domain=", -1);
DStringTrunc(&escapeUrl, 0);
DStringAppend(&escapeUrl, "(", -1);
      for (i=1; i < argc; i++)
                  = GetAsciiDomain(argv(i), NULL);
           if (cp != NULL)
               DStringAppend(&escapeUrl, cp, ~1);
DStringAppend(&escapeUrl, " ", -1);
     DStringAppend(&escapeUrl, *}*, -1);
      EscapeUrl (&escapeUrl);
     DStringAppend(&targetUrl,DStringValue(&escapeUrl), -1);
     DStringFree(&escapeUrl);
HTTP_Error(reqPtr, REDIRECT, DStringValue(&targetUrl));
IncTicketCounter(CountRemoteRedirects);
     DStringFree(&targetUrl);
      if (!ticketPtr->valid)
           DStringFree (&ticketPtr->sid);
     return TCL_RETURN;
/* authorization failed, if this is a ticket access, decode the */
/* reason and handle via a redirect to a handler, or punt a
/* no access message */
if ((flavor == ticketTicket) && (firstLegalDom != -1) && (ticketPtr->ticketD
             check For IP address restrictions */
            ((DStringLength(&ticketPtr->ticketIP) 1= 0) &&
                 (DStringLength(&TicketGlobalData(TicketAdrHandler)) != 0) &&
                 (strcmp(DStringValue(&ticketPtr->ticketIP), DStringValue(&reqPtr->remo
          DstringAppend(&targetUrl, DstringValue(&TicketGlobalData(TicketAdrHandle DstringAppend(&targetUrl, DstringValue(&ticketPtr->fields), -1);
DstringAppend(&targetUrl, "&url=", -1);
DstringAppend(&targetUrl, DstringValue(&reqPtr->url), -1);
IncTicketCounter(CountTicketAddr);
           HTTP_Error(reqPtr, REDIRECT, DStringValue(&targetUrl));
          DStringFree(&targetUrl);
          return TCL_RETURN;
   /* check for expired tickets */
  if (time(0) > ticketPtr->ticketExp)
          DStringAppend(&targetUrl, DStringValue(&TicketGlobalData(TicketExpHandle
DStringAppend(&targetUrl, DStringValue(&ticketPtr->fields), -1);
DStringAppend(&targetUrl, "&url=", -1);
DStringAppend(&targetUrl, DStringAppend(&ta
```

IncTicketCounter(CountExpiredTicket);

5,708,780 **49 50**

HTTP_Error(reqPtr, REDIRECT, DStringValue(&targetUrl));
DStringFree(&targetUrl);
return TCL_RETURN;
}

/* no handler, punt a message */
HTTP_Error(reqPtr, FORBIDDEN, "access denied by Require ticket/sid region co
IncTicketCounter(CountNoRedirects);
if (!ticketPtr->valid)
DStringFree(&ticketPtr->sid);
DStringFree(&targetUrl);
DstringFree(&escapeUrl);
return TCL_RETURN;
}

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```
Get [Ascii] Domain --
         These routine performs an ascii to binary domain name lookup, indexed by 'key') from the server's domain name catalog. Name/number pair's are loaded into the catalog at configuration time with the with the "Domain" configuration command. The Ascii version returns
          a pointer to a character string that represents the domain number.
          The non Ascii version returns an integer representing the domain number.
   Results:
          Integer value of domain. If no domain is available, returns deflt.
   Side effects:
         None.
static int GetDomain(char *domname, int deflt)
     HashEntry *entryPtr;
     DString DomName:
    DStringInit(&DomName);
    DStringAppend(&DomName, domname, -1);
    strtolower(DStringValue(&DomName));
    entryPtr = FindHashEntry(&TicketServerData.Domains, DStringValue(&DomName));
    DStringFree(&DomName);
if (entryPtr == NULL) return deflt;
    return (int) GetHashValue(entryPtr);
static char * GetAsciiDomain(char *domname, char *deflt)
    HashEntry *entryPtr;
    static char buffer (64);
    DString DomName;
    DStringInit(&DomName);
    DStringAppend(&DomName, domname, -1);
    strtolower(DStringValue(&DomName));
    entryPtr = FindHashEntry(&TicketServerData.Domains, DStringValue(&DomName));
    DStringFree (&DomName);
    if (entryPtr == NULL) return deflt;
    sprintf(buffer, "*d", (int) GetHashValue(entryPtr));
    return buffer;
```

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```
TICKET_InsertLocalSid --
                      Given a URL, inspect it to see if it refers to the local server/port
                     if it does, and it does not already contain a SID, insert one if the current request included one. Note, for port 80 access we look
                      for a match with and without the port specifier.
       Results:
                      None.
        Side effects:
                      A SID may be inserted into the URL.
                            void TICKET_InsertLocalSid(HTTP_Request *reqPtr, DString *result)
           HTTP_Server *serverPtr;
           TICKET_Request *ticketPtr;
           char tmp[32];
          DString pattern1; DString pattern2;
          DString tmp_url;
DString *hitPattern * NULL;
           ticketPtr = (TICKET_Request *) HT_GetReqExtData(reqPtr, TicketServerData.tic
           if (ticketPtr == NULL) return;
           serverPtr = reqPtr->serverPtr;
           DStringInit(&patternl);
           DStringInit (&pattern2);
           DStringInit(&tmp_url);
          DStringAppend(&pattern1, "http://", -1);
DStringAppend(&pattern1, DStringValue(&serverPtr->serverName), -1);
DStringAppend(&pattern2, DStringValue(&pattern1), -1);
sprintf(tmp, ":%d", serverPtr->server_port);
DStringAppend(*pattern2, DStringAppen
           DStringAppend(&pattern1, tmp, -1);
           if ((DStringLength(result) >= DStringLength(&pattern1)) &&
                      (strncasecmp(DStringValue(&pattern1), DStringValue(result), DStringLengt
                      hitPattern = &pattern1;
           else
           if ((serverPtr->server_port == 80) &&
     (DStringLength(result) >= DStringLength(&pattern2)) &&
                       (strncasecmp(DStringValue(&pattern2), DStringValue(result), DStringLengt
                      hitPattern = &pattern2;
           if (hitPattern != NULL)
               DStringAppend(&tmp_url, DStringValue(hitPattern), -1);
DStringAppend(&tmp_url, DStringValue(&ticketPtr->sid), -1);
DStringAppend(&tmp_url, &DStringValue(result)[DStringLength(hitPattern)],
DStringFree(result);
```

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```
DStringAppend(result, DStringValue(&tmp_url), -1);
DStringFree(&tmp_url);
)StringFree(&pattern1);
DStringFree(&pattern2);
DStringFree(&tmp_url);
)
```

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DStringFree (&hash);

/* fix_endian(&act_hash, ecp, eda); */

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5-5 CreateSid --This routine takes the passed arguments and creates a sid. Results: A sid. Side effects: char * CreateSid(HTTP_Request *reqPtr, int dom, int wid, int kid, int exp, int uctx) int bsid(3) = $\{0,0,0\}$; char temp_str[512]; DString hash; int act hash; static char sid[64]; unsigned int expire_time; char *secret; char *hashP; char *cp; unsigned char *ecp; unsigned int eda; int endian = 1; DStringInit (&hash); expire time = time(0) + exp; put_sid(dom_lw,
put_sid(uid_lw,
put_sid(kid_lw,
put_sid(exp_lw,
put_sid(uctx_lw,
put_sid(rev_lw, dom_pos, dom_mask, dom); uid_pos, uid_mask, uid); kid_pos, kid_mask, kid); (expire_time>>exp_shft_amt)) exp_pos, exp_mask, uctx_pos, uctx_mask, uctx); rev_pos, rev_mask, sid_rev_zero); secret = GetSecret(kid);
ASSERT (secret != NULL); ASSERT (secret 1s North;
DStringAppend(&hash, secret, -1);
DStringAppend(&hash, DStringValue(&reqPtr->remoteAddr), -1);
sprintf(temp str, "*08x*08x", bsid[2],bsid[1]);
DStringAppend(&hash, temp_str, -1);
/* format of the hash string is *s*s*08x*08x", secret,ip_addr,bsid[2],bsid[1] hashP = DStringValue(&hash); act hash = compute ihash(hashP); while (*hashP != 0) *hashP++ = 0;

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```
.put_sid(sig_lw, sig_pos, sig_mask, act_hash)

/* fix_endian(&bsid[0], ecp, eda); */
    fix_endian(&bsid[1], ecp, eda);
    fix_endian(&bsid[2], ecp, eda);

#if (1 == 0)
        DumpSid();
#endif

cp = radix64encode_noslash((char *) bsid, 12);
    strcpy(sid, SID_prefix);
    strcat(sid, cp);
    free(cp);
    return(sid);
}
```

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```
computeHash --
    Compute the MD5 hash for the specified string, returning the hash as
    a 32-character hex string.
 * Results:
          Pointer to static hash string.
   Side effects:
          None.
static char *computeHash(char *str)
    int i;
MD5_CTX md5;
    unsigned char hash[16];
static char hashstr[33];
char *q;
     MD5Init(&md5);
MD5Update(&md5, (unsigned char *) str, strlen(str));
MD5Final(hash, &md5);
     q = hashstr;
     for(i=0; i<16; i++) {
    sprintf(q, "%02x", hash[i]);
          q += 2;
     *q = '\0';
     return hashstr;
}
```

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```
* ParseSid --
           Called by TICKET_ParseTicket, before any region commands or mount handle
          have run. We parse and handle incomeing sid's.
 * Results:
          None.
 * Side effects:
int ParseSid(HTTP_Request *reqPtr)
     TICKET_Request *ticketPtr;
     HTTP Server *serverPtr;
     DString hash;
     int i;
     char *cp, *cpl;
int *bsid=NULL, act_hash;
     unsigned int cur_tim, tdif, exp_tim;
    char *secret;
char temp_str[512];
     char *hashP;
    cnar -nashr;
int sid ok = 0;
unsigned char *ecp;
unsigned int eda;
int endian = 1;
int ipl,ip2,ip3,ip4;
     /* fetch the server private ticket extension data \star//* note that this sets up a default ticket block for both SID's and Ticket a
    serverPtr = reqPtr->serverPtr;
ticketPtr = (TICKET_Request *) HT_GetReqExtData(reqPtr, TicketServerData.tic
    ASSERT (ticketPtr == NULL);
    ticketPtr = (TICKET_Request *) Malloc(sizeof(TICKET_Request));
HT_AddReqExtData(reqPtr, TicketServerData.ticketExtensionId, ticketPtr, free
    DStringInit (&ticketPtr->rawUrl);
    DStringInit (&ticketPtr->sid);
    DStringInit (&ticketPtr->fields);
    DStringInit (&ticketPtr->signature);
    DStringInit (&ticketPtr->ticketIP);
    ticketPtr->valid
                                = 0;
    ticketPtr->sidDom
    ticketPtr->ticketDom = -1;
    ticketPtr->ticketExp = -1;
                                = 0;
    ticketPtr->uid
                                 = 0:
    sscanf(DStringValue(&reqPtr->remoteAddr), "%d.%d.%d.%d.%d", &ip1, &ip2, &ip3, &
ticketPtr->uid = (((ip1+ip2)<<24) | ((ip3+ip4)<<16) | (rand() & 0xFFFF));
ticketPtr->uctx = 1;
    ticketPtr->uctx
```

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```
/* we are done if sids are not enabled, or this url does not have a sid */
 if (!(TicketGlobalData(EnableSid))) return HT OK;
  cpl = DStringValue(&reqPtr->url);
  if (strstr(cpl, SID_prefix) != cpl)
    return HT OK;
 if (strlen(cp1) == sidLength)
   DstringAppend(&reqPtr->url, */", -1);
DstringAppend(&reqPtr->path, */*, -1);
cpl = DstringValue(&reqPtr->url);
 cp'= strchr(cpl+sizeof(SID_prefix),'/');
if ((cp - cpl) != sidLength)
  return HT_OK;
 IncTicketCounter(CountSidUrl);
 DStringInit(&hash);
 /* if sid eater is enabled, rewrite the url without the sid, and reprocess t
 if (TicketGlobalData(EnableSidEater))
   DStringAppend(&hash, DStringValue(&reqPtr->url), -1);
   DStringFree(&reqPtr->url);
   DStringAppend(&reqPtr->url, DStringValue(&hash)+sidLength, -1);
   DStringTrunc(&hash, 0);
   DStringAppend(&hash, DStringValue(&reqPtr->path), -1);
   DStringFree (&reqPtr->path);
   DStringAppend(&reqPtr->path, DStringValue(&hash)+sidLength, -1);
   DStringFree (&hash);
   IncTicketCounter(CountDiscardedSidUrl);
   return HT_OK;
DStringAppend(&ticketPtr->sid, DStringValue(&reqPtr->url), sidLength);
/* first convert the SID back to binary*/
i = DStringLength(&ticketPtr->sid) -3;
bsid = (int *) radix64decode_noslash(DStringValue(&ticketPtr->sid)+3, i, &i)
if ((bsid == NULL) || (i != 12)) goto rtn_exit;
fix_endian(&bsid[0], ecp, eda);
fix_endian(&bsid[1], ecp, eda);
fix_endian(&bsid[2], ecp, eda);
/* check the SID version field */
if (get_sid(rev_lw,rev_pos,rev_mask) != sid_rev_zero) goto sid_bad;
if (get_sid(rsrvl_lw,rsrvl_pos,rsrvl_mask) != 0) goto sid_bad;
if (get_sid(rsrv2_lw,rsrv2_pos,rsrv2_mask) != 0) goto sid_bad;
/* Get a pointer to the secret */
secret = GetSecret(get_sid(kid_lw,kid_pos,kid_mask));
if (secret == NULL) goto sid_bad;
 /* hash the sid and check the signature */
DStringAppend(&hash, secret, -1);
```

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```
.DstringAppend(&hash, DstringValue(&reqPtr->remoteAddr), -1);
sprintf(temp_str, "%08x%08x", bsid[2],bsid[1]);
DstringAppend(&hash, temp_str, -1);
/* format of the hash string is %s%s%08x%08x", secret,ip_addr,bsid[2],bsid[1]
  hashP = DStringValue(&hash);
  act hash = compute ihash(hashP);
while (*hashP != 0) *hashP++ = 0;
   fix_endian(&act_hash, ecp, eda);
   if (act hash != get_sid(sig_lw,sig_pos,sig_mask)) goto sid_bad;
  /* is is ok, may be expired, but good enough to id user */
  ticketPtr->uid = get_sid(uid_lw,uid_pos,uid_mask);
ticketPtr->uctx = get_sid(uctx_lw,uctx_pos,uctx_mask);
   /* do the SID experation processing */
  cur tim = (time(0)>>exp shft amt) & exp mask;
exp tim = get_sid(exp lw,exp pos,exp mask);
tdif = (exp_tim - cur_tim) & 0xffff;
if (tdif > 0x7fff)
     incTicketCounter(CountExpSid);
     goto sid_exp;
  /* sid is fine, save the sid state, update the url's */
  ticketPtr->sidDom = get_sid(dom_lw,dom_pos,dom_mask);
  ticketPtr->valid = 1;
  sid_ok = 1;
IncTicketCounter(CountValidSid);
sid_bad:
  if (!(sid_ok)) IncTicketCounter(CountInvalidSid);
sid_exp:
  DStringAppend(&ticketPtr->rawUrl, DStringValue(&reqPtr->path), -1);
  DStringTrunc(&reqPtr->path, 0);
  DStringAppend(&reqPtr->path, DStringValue(&ticketPtr->rawUrl)+sidLength, -1)
  DStringTrunc(&ticketPtr->rawUrl, 0);
  DStringAppend(&ticketPtr->rawUrl, DStringValue(&reqPtr->url), -1);
  DStringTrunc(&reqPtr->url, 0);
DStringAppend(&reqPtr->url, DStringValue(&ticketPtr->rawUrl)+sidLength, -1);
rtn_exit:
  DStringFree (&hash);
  if (bsid != NULL) free(bsid);
  return HT_OK;
```

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```
/*

* freeTicketReqData

* This routine frees the storage used by ticket specific request
          data.

* Results:
          None.

* Side effects:
          Memory freed.

* TICKET_Request *ticketPtr = dataPtr;
          DStringFree(&ticketPtr->rawUrl);
          DStringFree(&ticketPtr->rields);
          DStringFree(&ticketPtr->sid);
          DStringFree(&ticketPtr->sid);
          DStringFree(&ticketPtr->signature);
          DStringFree(&ticketPtr->signature);
          DStringFree(&ticketPtr->ticketIP);
          free(ticketPtr);
}
```

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/*

GetKidByKeyID -
Given an ascii KeyID return the binary KeyID.

for untranslatable names, return -1.

Results:

"I've got a secret, now you do too"

Side effects:

"

HashEntry *entryPtr;

entryPtr = FindHashEntry(&TicketServerData.KeyID, (void *) keyID);
if(entryPtr == NULL) return -1;
return (int) GetHashValue(entryPtr);
}

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```
* fieldParse --
     Given a string, a separator character, extracts a field up to the separator into the result string. Does substitution on '*XX' sequences, and returns the pointer to the character beyond last character in '*endptr'.
           Returns a malloc'ed string (caller must free), or NULL if an
           error occurred during processing (such as an invalid '%' sequence).
  * Side effects:
           None.
#define SIZE INC 200 static char *fieldParse(char *str, char sep, char **endptr)
      char buf[3];
     char c;
char *end, *data, *p;
int maxlen, len;
     len = 0;
     maxlen = SIZE INC;
p = data = malloc(maxlen);
    * Loop through string, until end of string or sep character.
     while (*str && *str != sep) {
           if(*str == '%') {
                 if(!isxdigit(str[1]) | !isxdigit(str[2])) {
                       free (data);
                       return NULL;
                buf[0] = str[1];
buf[1] = str[2];
buf[2] = '\0';
                 c = strtol(buf, &end, 16);
           str += 3;
} else if(*str == '+') {
  c = ' ';
                 str++;
           ) else
                 c = *str++;
           *p++ = C;
           len++;
           if(len >= maxlen) {
  maxlen += SIZE_INC;
  data = realloc(data, maxlen);
                 p = data + len;
```

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// *p++ = '\0';
// *endptr = str;
// return data;

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```
DomainNameCmd --
        A call to this routine, builds the ascii domain name
        to binary domain name maping structure for a numeric domain. Syntax is Domain number namel name2 name3 name... name_last at least one name is required. The number is decimal and can be any value except -1. -1 is reserved as a marker
        for untranslatable names.
   Results:
        None.
   Side effects:
        Commands are validate, and entries added to the map
static int DomainNameCmd(ClientData clientData, Tcl_Interp *interp,
                              int argc, char **argv)
    int new, 1;
    HashEntry *entryPtr;
    int DomNumber;
    DString DomName;
    if (argc <3)
      return TCL_ERROR;
    DStringInit(&DomName);
    if (((sscanf(argv[1], "*d", &DomNumber) != 1) | (DomNumber == "1)))
      ", value found was ",argv[1],
                 (char *) NULL);
      return TCL_ERROR;
    for (i = 2; i < argc; i++)
      DStringFree (&DomName);
      DStringAppend(&DomName, argv[i], -1);
strtolower(DStringValue(&DomName));
      entryPtr = CreateHashEntry(&TicketServerData.Domains, DStringValue(&DomNam
      if (new == 0)
        Tcl_AppendResult(interp, argv[0], " directive: ".
```

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```
SecretsCmd --
      A call to this routine, builds kid to secrets table
  Results:
      None.
  Side effects:
      Secrets are stored.
int newKid, newKeyID;
   HashEntry *entryPtrKid = NULL , *entryPtrKeyID = NULL;
   int Kid;
   DString *dsptrKid;
   if (argc != 4)
     (char *) NULL);
     return TCL_ERROR;
   if (sscanf(argv[2], "*d", &Kid) != 1)
     (char *) NULL);
     return TCL_ERROR;
   entryPtrKid = CreateHashEntry(&TicketServerData.SecretsKid, (void *) Kid, &n
   if (strlen(argv[1]))
  entryPtrReyID = CreateHashEntry(&TicketServerData.KeyID, (void *) argv[1],
if ((newKid == 0) [] ((newKeyID == 0) && strlen(argv[1])))
    argv[1],
(char *) NULL);
     return TCL_ERROR;
   if (strlen(argv[1]))
    dsptrkid = (Dstring *) malloc(sizeof(Dstring));
DstringInit(dsptrKid);
    DStringAppend(dsptrKid, argv[3], -1);
```

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SetHashValue(entryPtrKid, dsptrKid);
}

SetHashValue(entryPtrKeyID, Kid);
return TCL_OK;
}

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```
TICKET_Initialize --
          Calls all the necessary routines to initialize the ticket subsystem.
   Results:
          None.
   Side effects:
          Commands added to the region interpreter.
          SID "/@@" url catcher declared.
 */
int TICKET_Initialize(HTTP_Server *serverPtr, Tcl_Interp *interp)
    TicketServerData.ticketExtensionId = HT RegisterExtension(serverPtr, *ticket
    InitHashTable(&TicketServerData.SecretsKid, TCL_ONE_WORD_KEYS);
InitHashTable(&TicketServerData.KeyID, TCL_STRING_KEYS);
    InitHashTable(&TicketServerData.Domains, TCL_STRING_KEYS);
    /* initialize Server ticket data */
    DStringInit (&TicketGlobalData(AuthServer));
    DStringInit(&TicketGlobalData(TicketExpHandler));
    DStringInit (&TicketGlobalData (TicketAdrHandler));
    TicketGlobalData(FreeArea)
                                                         = 0
    TicketGlobalData(EnableLocalAuth)
                                                         = 0;
    TicketGlobalData(CurrentSecret)
TicketGlobalData(EnableSid)
                                                         = 0:
                                                         - 0:
    TicketGlobalData(EnableTicket)
TicketGlobalData(EnableSidEater)
                                                         - 0;
                                                           0 :
    TicketGlobalData(LocalAuthExp)
                                                         = 60*30;
    /* ticket event counters *.
    TicketGlobalData(CountTotalUrl)
                                                         = 0;
    TicketGlobalData(CountSidUrl)
                                                           ٥,
    TicketGlobalData(CountValidSid)
                                                         = 0;
    TicketGlobalData(CountExpSid)
                                                           0;
    TicketGlobalData(CountInvalidSid)
                                                         = 0;
    TicketGlobalData(CountCrossDomain)
    TicketGlobalData(CountLocalRedirects)
                                                         = 0;
    TicketGlobalData(CountRemoteRedirects)
    TicketGlobalData(CountNoRedirects)
    TicketGlobalData(CountDiscardedSidUrl)
    /* Ticket related Config commands */
    Tcl_CreateCommand(interp, "Domain",
                                                                    DomainNameCmd.
                           (ClientData) serverPtr, NULL);
    Tcl_CreateCommand(interp, "Secrets",
                                                                    SecretsCmd,
   TCl_CreateCommand(interp, "Secrets",

(ClientData) serverPtr, NULL);

Tcl_CreateCommand(interp, "AuthenticationServer", CmdStringValue,

(ClientData) &TicketGlobalData(AuthServer), NULL);

Tcl_CreateCommand(interp, "TicketExpirationHandler", CmdStringValue,

(ClientData) &TicketGlobalData(TicketExpHandler), NULL);
```

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```
TICKET_Shutdown --
         Calls all the necessary routines to shutdown the ticket subsystem.
   Results:
         None.
   Side effects:
         Memory freed
void TICKET_Shutdown(HTTP_Server *serverPtr)
    HashEntry *entryPtr;
    HashSearch search;
DString *dstring;
    DStringFree(&TicketGlobalData(AuthServer));
DStringFree(&TicketGlobalData(TicketExpHandler));
DStringFree(&TicketGlobalData(TicketAdrHandler));
    entryPtr = FirstHashEntry(&TicketServerData.SecretsKid, &search);
    while (entryPtr != NULL)
       dstring = GetHashValue(entryPtr);
       DStringFree (dstring);
       free (dstring);
       entryPtr = NextHashEntry(&search);
    DeleteHashTable(&TicketServerData.SecretsKid);
    DeleteHashTable(&TicketServerData.KeyID);
    DeleteHashTable(&TicketServerData.Domains);
```

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```
/*

* TICKET_GetCGIVariables --

* Add TICKET CGI variables to the CGI variable table.

* Results:

* None.

* Side effects:

* Extends the CGI variable hash table.

* TICKET_GetCGIVariables (HTTP_Request *req)

* TICKET_Request *ticketPtr = (TICKET_Request *) HT_GetReqExtData(req, TicketS)

* If there's no extension data, then we're not doing a ticket. Just return

* (ticketPtr == NULL)

* return;

if (DstringLength(&ticketPtr->rawUrl) != 0)

* HT AddCGIParameter(req, "TICKET_URL", DstringValue(&ticketPtr->rawUrl), FA

if (DstringLength(&ticketPtr->sid) != 0)

* HT AddCGIParameter(req, "TICKET_SIST, DstringValue(&ticketPtr->sid), FALSE

if (DstringLength(&ticketPtr->fields) != 0)

* HT AddCGIParameter(req, "TICKET_FIELDS", DstringValue(&ticketPtr->fields),

if (DstringLength(&ticketPtr->fields) != 0)

* HT_AddCGIParameter(req, "TICKET_FIELDS", DstringValue(&ticketPtr->fields),

if (DstringLength(&ticketPtr->signature) != 0)

* HT_AddCGIParameter(req, "TICKET_SIGNATURE", DstringValue(&ticketPtr->signa
```

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```
TICKET_ConfigCheck
        Perform late configuration checks
 * Results:
 * Side effects:
        Possible message loged/printed, and program exit'd.
void TICKET_ConfigCheck()
   HashEntry *entryPtr;
    int kid;
    if ({TicketGlobalData(EnableSid) & ~0x1) != 0}
      LogMessage(LOG_ERR, "EnableSid must be 0 or 1");
      exit(0);
    if (!(TicketGlobalData(EnableSid))) return;
    kid = TicketGlobalData(CurrentSecret);
    if ((kid && kid_mask) != kid)
      LogMessage(LOG_ERR, *CurrentSecret %d is invalid*, kid);
      exit(0);
    entryPtr = FindHashEntry(&TicketServerData.SecretsKid, (void *) kid);
    if (entryPtr == NULL)
      LogMessage(LOG_ERR, "No secret defined for CurrentSecret %d", kid);
      exit(0);
    if ((TicketGlobalData(FreeArea) & -0x255) != 0)
      LogMessage(LOG_ERR, "FreeArea must be between 0 and 255");
      exit(0);
    if ((TicketGlobalData(EnableSidEater) & -0x1) != 0)
      LogMessage(LOG_ERR, "EnableSidEater must be 0 or 1");
      exit(0);
    if ((TicketGlobalData(EnableTicket) & -0x1) != 0)
      LogMessage(LOG_ERR, "EnableTicket must be 0 or 1");
      exīt(0);
```

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```
if ((TicketGlobalData(EnableLocalAuth) & ~0x1) != 0)
{
  LogMessage(LOG_ERR, "EnablLocalAuth must be 0 or 1");
  exit(0);
}
```

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```
DumpStatus --
        Dump the server's ticket stat's
  Results:
        None.
  Side effects:
        None.
#define BUFSI2E 1024
static void DumpStatus(HTTP_Request *reqPtr)
   HTTP_Server *serverPtr = reqPtr->serverPtr;
   char tmp(BUFSIZE), timeStr[BUFSIZE];
struct utsname sysinfo;
    time_t uptime;
    int hours;
   HTTP_Send(reqPtr, "<hr><hr><hr><ht>Ticket Log</h2>","\n", NULL);
    sprintf(tmp, "
                        <b>$8:
                                </b> <d\n", "Number of access
                                                                            ", Ticket
   HTTP Send(reqPtr, tmp, NULL);
   sprintf(tmp,
                        <b>*s: </b>
                                      %d\n", "Number of SID URL's
                                                                            ", Ticket
   HTTP Send(reqPtr, tmp, NULL);
sprintf(tmp, " <b>ts: </
                                </b>
                                      %d\n". "Number of Valid SID's
                                                                            ", Ticket
   sprintf(tmp,
   HTTP Send(reqPtr, tmp, NULL);
sprintf(tmp, " <b>ts: </
    sprintf(tmp,
                                      %d\n", "Number of Expired SID's
                                                                            ", Ticket
                                </b>
                      tmp, NULL); <b>ts: </b>
   HTTP_Send(reqPtr,
                                      %d\n", "Number of Invalid SID's
                                                                            ". Ticket
   sprintf(tmp,
   HTTP_Send(reqPtr, tmp, NULL);
sprintf(tmp, " <b>*s: </
                                </b>
                                      %d\n", "Number of XDomain accesses", Ticket
   sprintf(tmp,
   HTTP Send(reqPtr, sprintf(tmp, "
                      tmp, NULL); <b>ts: </b>
                                       %d\n", "Number of Local Redirects ", Ticket
   sprintf(tmp,
                      tmp, NULL); <b>*s: </
   HTTP_Send(reqPtr,
   sprintf(tmp, "
                                </b>
                                      %d\n", "Number of Remote Redirects", Ticket
   HTTP_Send(reqPtr,
                      tmp, NULL);
                                      %d\n", "Number of No Auth servers ", Ticket
   sprintf(tmp,
                        <b>%s: </b>
   HTTP_Send(reqPtr, tmp, "", NULL);
   uptime = time(NULL) - serverPtr->started;
   uname(&sysinfo);
   strftime(timeStr, BUFSIZE, "%A, %d-%b-%y %T",
       localtime(&serverPtr->started));
```

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```
HTTP_Send(reqPtr, tmp, NULL);
    sprintf(tmp, "
                                                    </b> <d\n",
                      <b>Number of connections:
    serverPtr->numConnects);
HTTP_Send(reqPtr, "", tmp, NULL);
    sprintf(tmp, *
                      <br/><b>Number of HTTP requests:
                                                    </b> *d\n",
   serverPtr->numRequests);
HTTP_Send(reqPtr, tmp, "", NULL);
    hours = max(uptime / 3600, 1);
sprintf(tmp, "This server is averaging <b>*d</b> requests per hour.*,
           serverPtr->numRequests/hours);
    HTTP_Send(reqPtr, tmp, NULL);
    DumpRusage (reqPtr);
   DumpConnections(reqPtr); */
    DNS_DumpStats(reqPtr);
   HTTP_Send(reqPtr, "<hr><address>", DStringValue(&ht_serverSoftware),
            "</address>\n", NULL);
    regPtr->done = TRUE;
#undef BUFSIZE
```

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User: morris Host: uprism.openmarket.com Class: uprism.openmarket.com Job: t.t

What is claimed is:

- 1. A method of processing service requests from a client to a server system through a network, said method comprising the steps of:
 - forwarding a service request from the client to the server system, wherein communications between the client and server system are according to hypertext transfer protocol;
 - returning a session identifier from the server system to the client: and
 - appending as part of a path name in a uniform resource locator the session identifier to the request and to subsequent service requests from the client to the server system within a session of requests.
- 2. A method as claimed in claim 1 wherein the session identifier includes a user identifier.
- A method as claimed in claim 1 wherein the session identifier includes an expiration time for the session.
- 4. A method as claimed in claim 1 wherein the server system records information from the session identifier in a transaction log in the server system.
- 5. A method as claimed in claim 4 wherein the server system tracks the access history of sequences of service requests within a session of requests.
- 6. A method as claimed in claim 5 wherein the server system tracks the access history to determine service requests leading to a purchase made within the session of requests.
- 7. A method as claimed in claim 4 wherein the server system counts requests to particular services exclusive of repeated requests from a common client.
- 8. A method as claimed in claim 4 wherein the server system maintains a data base relating customer information to access patterns.
- 9. A method as claimed in claim 8 wherein the information includes customer demographics.
- 10. A method as claimed in claim 1 wherein the server system assigns the session identifier to an initial service request to the server system.
- 11. A method as claimed in claim 1 wherein the server system subjects the client to an authorization routine prior to issuing the session identifier and the session identifier is protected from forgery.
- 12. A method as claimed in claim 1 wherein the server system comprises plural servers including an authentication server which provides session identifiers for service requests to multiple servers.
 - 13. A method as claimed in claim 12 wherein:
 - a client directs a service request to a first server which is to provide the requested service;
 - the first server checks the service request for a session identifier and only services a service request having a valid session identifier, and where the service request has no valid identifier:
 - the first server redirects the service request from the client to the authorization server;
 - the authorization server subjects the client to the authorization routine and issues the session identifier to be appended to the service request to the first server;
 - the client forwards the service request appended with the session identifier to the first server; and
 - the first server recognizes the session identifier and services the service request to the client; and
 - the client appends the session identifier to subsequent 65 service requests to the server system and is serviced without further authorization.

- 14. A method as claimed in claim 13 wherein the session identifier includes a user identifier.
- 15. A method as claimed in claim 13 wherein the session identifier includes an expiration time for the session.
- 16. A method as claimed in claim 13 wherein the session identifier provides access to a protected domain to which the session has access authorization.
- 17. A method as claimed in claim 16 wherein the session identifier is modified for access to a different protected to domain.
 - 18. A method as claimed in claim 13 wherein the session identifier provides a key identifier for key management.
- 19. A method as claimed in claim 13 wherein the server system records information from the session identifier in a 15 transaction log in the server system.
- 20. A method as claimed in claim 13 wherein the client modifies the path name of a current uniform resource locator using relative addressing and retains the session identifier portion of the path name unmodified for successive requests in the session.
 - 21. A method as claimed in claim 1 wherein:
 - the server system subjects the client to an authorization routine prior to issuing the session identifier and the session identifier is protected from forgery, records information from the session identifier in a transaction log in the server system, tracks request paths relative to hypertext pages, and maintains a data base relating customer demographics to access patterns; and
 - the client modifies the path name of a current uniform resource locator using relative addressing and retains the session identifier portion of the path name unmodified for successive requests in a session.
- 22. A method of processing service requests from a client to a server system through a network, said method comprising the steps of:
 - appending as part of a path name in a uniform resource locator a session identifier to the request, wherein communications between the client and server system are according to hypertext transfer protocol;
 - responding to requests for hypertext pages received from a client through the network by returning the requested hypertext pages to the client;
 - responding to further client requests related to links in the hypertext pages; and tracking the further client requests related to a particular hypertext page.
 - 23. A method as claimed in claim 22 wherein the requests include a common session identifier and the server system tracks client requests within a session of requests.
 - 24. A method of processing service requests from a client to a server system through a network, said method comprising the steps of:
 - appending a session identifier to the request as part of a path name in a uniform resource locator, wherein communications between the client and server system are according to hypertext transfer protocol; and
 - responding to requests for documents received from the client through the network by returning the requested documents wherein the documents are customized for a particular user based on a user profile.
 - 25. A method of processing service requests from a client to a server system through a network, said method comprising the steps of:
 - responding to a request for a document received from the client through the network, wherein communications between the client and server system are according to hypertext transfer protocol;

appending a session identifier, which includes a user identification to the request as part of a path name in a uniform resource locator, and

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returning the requested document wherein the document is customized for a particular user based on the user 5 identification of the session identifier.

26. A method of processing service requests from a client to a server system through a network, said method comprising the steps of:

appending a session identifier to the request as part of a path name in a uniform resource locator, wherein communications between the client and server system are according to hypertext transfer protocol;

responding to requests for information received from the client through the network by returning the requested information to the client; and

counting requests to particular information exclusive of repeated requests from a common client.

27. A method as claimed in claim 26 comprising excluding from the counting requests made for information from the client within a defined period of time.

28. A method of processing service requests for a document received from a client through a network in which the document has been purchased by a user, said method comprising a daccess patterns. prising the steps of:

38. An information of processing service requests for a document and information of the document has been purchased by a user, said method comprising a daccess patterns.

responding to a request for a document received from a client through the network in which the document has been purchased by the user wherein communications between the client and server system are according to 30 hypertext transfer protocol;

appending an authorization identifier to the request as part of a path name in a uniform resource locator; and

returning the requested document if the authorization identifier indicates that the user is authorized to access the document.

29. A method as claimed in claim 28, wherein the authorization identifier is encoded within a session identifier which is appended to the request as part of a path name in a uniform resource locator.

30. A method of processing service requests from a client to a server system through a network, said method comprising the steps of:

responding to a request for a document received from a client through the network, wherein communications between the client and server system are according to hypertext transfer protocol;

appending as part of a path name in a uniform resource locator a session identifier to the request;

returning the requested document to the client, and; charging the user identified in the session identifier for access to the document.

31. A method as claimed in claim 30, wherein a user identifier is encoded within a session identifier which is 55 appended to the request.

32. An information system on a network, comprising:

means for receiving service requests from clients and for determining whether a service request includes a session identifier, wherein communications between the client and server system are according to hypertext transfer protocol;

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means for appending the session identifier as part of a path name in a uniform resource locator in response to an initial service request in a session of requests; and

means for servicing service requests from a client which include the session identifier, the subsequent service request being processed in the session.

33. An information system as claimed in claim 32 wherein the means for providing the session identifier is in a server system which services the requests.

34. An information system as claimed in claim 32 further comprising an authorization routine for authorizing the client prior to issuing the session identifier and means for protecting the session identifier from forgery.

35. An information server system as claimed in claim 32 further comprising a transaction log for recording information from the session identifier.

36. An information system as claimed in claim 32 further comprising means for tracking access history of sequences of service requests within the session of requests.

37. An information system as claimed in claim 32 further comprising means for counting requests to particular services exclusive of repeated requests from a common client.

38. An information system as claimed in claim 32 further comprising a data base relating customer information to access patterns.

39. An information system as claimed in claim 38 wherein the information includes customer demographics.

40. An information server on a network, comprising:

means for appending a session identifier as part of a path name in a uniform resource locator, wherein communications between the client and server system are according to hypertext transfer protocol;

means for responding to requests for hypertext pages received from a client through the network by returning the requested hypertext pages to the client;

means for responding to further requests derived from links in the hypertext pages; and means for tracking the further requests derived from a particular hypertext page.

41. A server as claimed in claim 40 wherein the requests include a common session identifier and the server tracks requests within a session of requests.

42. A server as claimed in claim 41 further comprising a data base relating customer demographics to access patterns.

43. An information server on a network, comprising:

means for appending the session identifier as part of a path name in a uniform resource locator, wherein communications between the client and server system are according to hypertext transfer protocol;

means for responding to requests for service received from a client through a network by returning the requested service to the client; and

means for counting requests to particular service exclusive of repeated requests from a common client.

44. A server as claimed in claim 43 wherein the requests include a common session identifier and the server tracks requests within a session of requests.

45. A server as claimed in claim 43 further comprising means for excluding requests made to a service from the client within a defined period of time.

* * * * *

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Civil Cover Sheet

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UNITED STATES DISTRICT COURT NORTHERN DISTRICT OF ILLINOIS

Civil Cover Sheet

This automated JS-44 conforms generally to the manual JS-44 approved by the Judicial Conference of the United States in September 1974. The data is required for the use of the Clerk of Court for the purpose of initiating the civil docket sheet. The information contained herein neither replaces nor supplements the filing and service of pleadings or other papers as required by law. This form is authorized for use onlyin the Northern District of Illinois.

JUDGE JOAN H. LEFKOW

Plaintiff(s): DIVINE TECHNOLOGY VENTURES, an Illinois general partnership

County of Residence: Cook County, Illinois

Plaintiff's Atty:

SEE ATTACHED

Defendant(s):REDENVELOPE, INC., a Delaware corporation

County of Residence: MAGISTRATE JUDGE KEYS

II. Basis of Jurisdiction:

3. Federal Question (U.S. not a party)

III. Citizenship of Principal Parties (Diversity Cases Only)

> Plaintiff:-N/A Defendant:-N/A

IV. Origin:

1. Original Proceeding

V. Nature of Suit:

830 Patent

VI.Cause of Action:

35 U.S.C. 271 - patent infringement

VII. Requested in Complaint

Class Action: No Dollar Demand: Jury Demand: Yes

VIII. This case **IS NOT** a refiling of a previously dismissed case.

Signature:

Date:

If any of this information is incorrect, please go back to the Civil Cover Sheet Input form using the Babbiton in your browser and change it. Once correct, print this form, sign and date it and submit it with your new civil action. Note: You may need to adjust the font size in your browser display to make the form print properly. Revised: 06/28/00

ATTACHMENT

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UNITED STATES DISTRICT COURT NORTHERN DISTRICT OF ILLINOIS

DOCKETED OCT 2 8 2002

In the Matter of DIVINE TECHNOLOGY VENTURES

v. REDENVELOPE, INC.



DIVINE TECHNOLOGY VENTURES

Plaintiff

Plaintiff

Plaintiff APPEARANCES ARE HEREBY FILED BY THE UNDERSIGNED AS ATTORNEY(S) FOR AN IN. LUNKOW (B) SIGNATURE SIGNATURE NAME / NAME Michael J. Abernathy Robert M. Barrett FIRM FIRM Bell, Boyd & Lloyd LLC Bell, Boyd & Lloyd LLC STREET ADDRESS STREET ADDRESS 70 W. Madison Street, Suite 3300 70 W. Madison Street, Suite 3300 CITY/STATE/ZIP CITY/STATE/ZIP Chicago, Illinois 60602 Chicago, Illinois 60602 TELEPHONE NUMBER TELEPHONE NUMBER (312) 372-1121 (312) 372-1121 IDENTIFICATION NUMBER (SEE ITEM 4 ON REVERSE) IDENTIFICATION NUMBER (SEE ITEM 4 ON REVERSE) 6183736 6180039 MEMBER OF TRIAL BAR? YES X MEMBER OF TRIAL BAR? NO [NO I YES [X] NO L NO L TRIAL ATTORNEY? YES X TRIAL ATTORNEY? YES X DESIGNATED AS LOCAL COUNSEL? YES [] NO X DESIGNATED AS LOCAL COUNSEL? YES | NO [X] (C) (D) SIGNATURE Jason W. Engel Lisa A. Carroll FIRM FIRM Bell, Boyd & Lloyd LLC Bell, Boyd & Lloyd LLC STREET ADDRESS STREET ADDRESS 70 W. Madison Street, Suite 3300 70 W. Madison Street, Suite 3300 CITY/STATE/ZIP CITY/STATE/ZIP Chicago, Illinois 60602 Chicago, Illinois 60602 TELEPHONE NUMBER TELEPHONE NUMBER (312) 372-1121 (312) 372-1121 IDENTIFICATION NUMBER (SEE ITEM 4 ON REVERSE) IDENTIFICATION NUMBER (SEE ITEM 4 ON REVERSE) 6274878 6255717 MEMBER OF TRIAL BAR? NO X MEMBER OF TRIAL BAR? YES TRIAL ATTORNEY? NO [] TRIAL ATTORNEY? NO [] YES [X] YES [X] NO X NO [X] DESIGNATED AS LOCAL COUNSEL? DESIGNATED AS LOCAL COUNSEL?